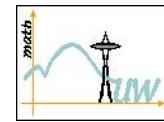


Smoothing methods for Second-Order Cone Programs/Complementarity Problems

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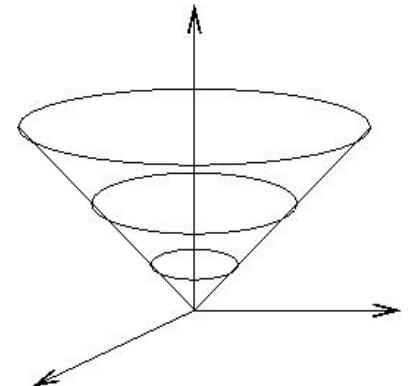
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Talk Outline

- Second-Order Cone (SOC) Program and Complementarity Problem
- Unconstrained Diff. Min. Reformulation
- Numerical Experience
- SOCP from Dist. Geometry Optim



Convex SOCP

$$\begin{array}{ll} \min & g(x) \\ \text{s.t.} & Ax = b \\ & x \in K \end{array}$$

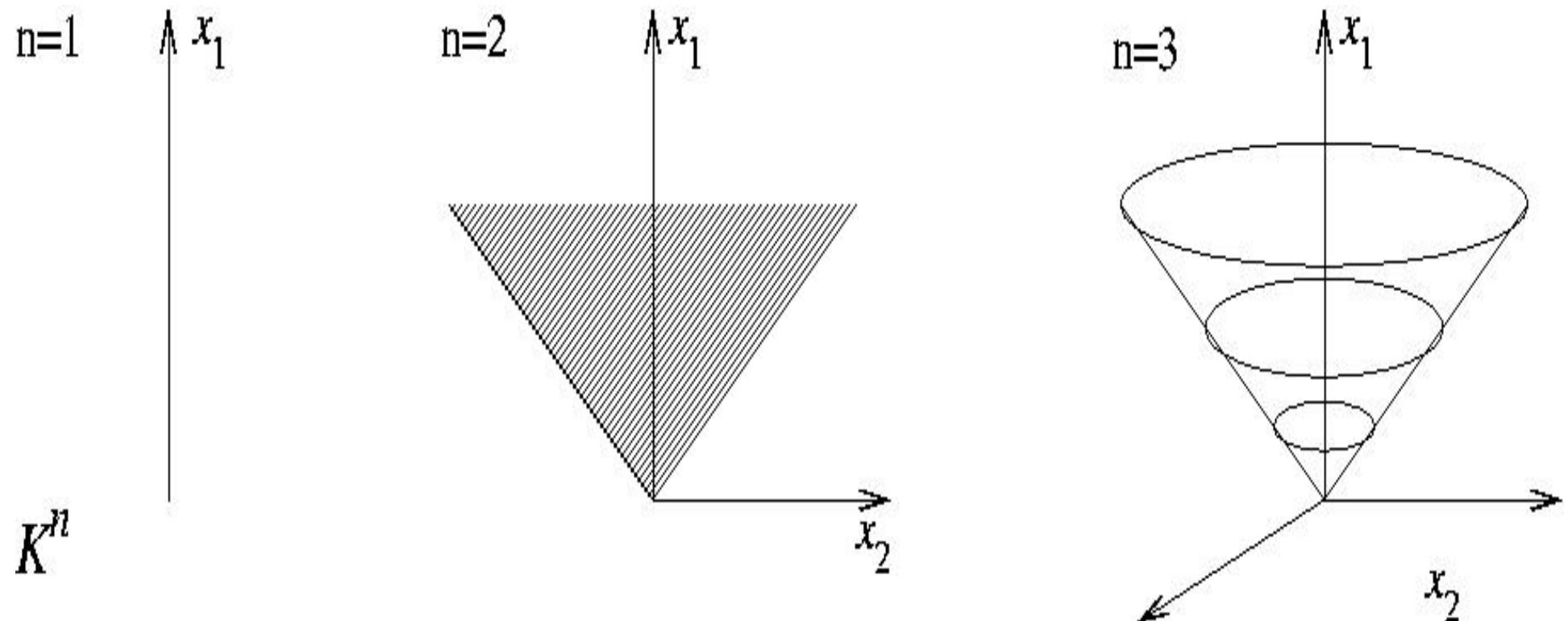
$$A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$$

$g : \mathbb{R}^n \rightarrow \mathbb{R}$, convex, twice cont. diff.

$$K = K^{n_1} \times \cdots \times K^{n_p}$$

$$K^{n_i} \stackrel{\text{def}}{=} \left\{ x_i = \begin{bmatrix} x_{1i} \\ x_{2i} \end{bmatrix} \in \mathbb{R} \times \mathbb{R}^{n_i-1} : \|x_{2i}\|_2 \leq x_{1i} \right\}$$

Special cases? LP, SOCP,...

SOC K^n 

Suff. Optim. Conditions

$$\begin{aligned} x \in K, \quad y \in K, \quad x^T y = 0, \\ Ax = b, \quad y = \nabla g(x) - A^T \zeta_d \end{aligned}$$

\iff

$$\begin{aligned} x \in K, \quad y \in K, \quad x^T y = 0, \\ x = F(\zeta), \quad y = G(\zeta) \end{aligned}$$

with

$$\begin{aligned} F(\zeta) &= d + (I - A^T(AA^T)^{-1}A)\zeta \\ G(\zeta) &= \nabla g(F(\zeta)) - A^T(AA^T)^{-1}A\zeta \quad (Ad = b) \end{aligned}$$

SOCCP

Find $\zeta \in \Re^n$ satisfying

$$x \in K, \quad y \in K, \quad x^T y = 0,$$

$$x = F(\zeta), \quad y = G(\zeta)$$

$F, G : \Re^n \rightarrow \Re^n$ smooth

$\nabla F(\zeta), -\nabla G(\zeta)$ column-monotone $\forall \zeta \in \Re^n$, i.e.,

$$\nabla F(\zeta)u - \nabla G(\zeta)v = 0 \quad \Rightarrow \quad u^T v \geq 0$$

Special cases? convex SOCP, monotone NCP,...

How to solve SOCCP?

For LP, simplex methods and interior-point methods.

For SOCP, interior-point methods.

For convex SOCP and column-monotone SOCCP?

Interior-point methods not amenable to warm start. Non-interior methods?

Nonsmooth Eq. Reformulation

$$x_i \cdot y_i \stackrel{\text{def}}{=} \begin{bmatrix} x_{1i} \\ x_{2i} \end{bmatrix} \cdot \begin{bmatrix} y_{1i} \\ y_{2i} \end{bmatrix} = \begin{bmatrix} x_i^T y_i \\ x_{1i}y_{2i} + y_{1i}x_{2i} \end{bmatrix} \quad (\text{Jordan product assoc. with } K^{n_i})$$

$$\phi_{\text{FB}}(x, y) \stackrel{\text{def}}{=} \left[(x_i^2 + y_i^2)^{1/2} - x_i - y_i \right]_{i=1}^p$$

Fact (Fukushima,Luo,T '02):

$$\phi_{\text{FB}}(x, y) = 0 \iff x \in K, y \in K, x^T y = 0$$

Thus, SOCCP is equivalent to

$$\phi_{\text{FB}}(F(\zeta), G(\zeta)) = 0$$

ϕ_{FB} is strongly semismooth (Sun,Sun '03)

Unconstr. Smooth Min. Reformulation

$$\min f_{\text{FB}}(\zeta) \stackrel{\text{def}}{=} \|\phi_{\text{FB}}(F(\zeta), G(\zeta))\|_2^2$$

F, G smooth and $\nabla F(\zeta), -\nabla G(\zeta)$ column-monotone $\forall \zeta \in \Re^n$ (e.g., LP, SOCP, convex SOCP, monotone NCP)

For monotone NCP ($K = \Re_+^n$),

f_{FB} is smooth, and $\nabla f_{\text{FB}}(\zeta) = 0 \iff \zeta$ is a soln

(Geiger,Kanzow '96)

The same holds for SOCCP.

(J.-S. Chen,T '04)

Advantage? Any gradient method for unconstrained diff. min. (e.g., CG, BFGS, L-BFGS) can be used to find $\nabla f_{\text{FB}}(\zeta) = 0$.

Numerical Experience on Convex SOCP

$$\begin{aligned} x &= F(\zeta) = d + (I - P)\zeta \\ y &= G(\zeta) = \nabla g(F(\zeta)) - P\zeta \end{aligned}$$

with $P = A^T(AA^T)^{-1}A$, $Ad = b$. (Solve $\min \|Ax - b\|$ to find d)

- Implement in Matlab CG-PR, BFGS, L-BFGS (memory=5) to minimize $f_{\text{FB}}(\zeta)$, using Armijo stepsize rule, with $\zeta^{\text{init}} = 0$. Stop when

$$\max\{f_{\text{FB}}(\zeta), |x^T y|\} \leq \text{accr.}$$

- Let $\psi_{\text{FB}}(x, y) \stackrel{\text{def}}{=} \|\phi_{\text{FB}}(x, y)\|_2^2$. Then

$$\begin{aligned} f_{\text{FB}}(\zeta) &= \psi_{\text{FB}}(x, y) \\ \nabla f_{\text{FB}}(\zeta) &= (I - P)\nabla_x \psi_{\text{FB}}(x, y) - P\nabla_y \psi_{\text{FB}}(x, y) \end{aligned}$$

Compute $P\zeta$ using Cholesky factorization of AA^T or using preconditioned CG. Compute $\psi_{\text{FB}}(x, y)$ and $\nabla \psi_{\text{FB}}(x, y)$ within Fortran Mex files.

DIMACS Challenge SOCPs

- Problem names and statistics:

$\text{nb } (m = 123, n = 2383, K = (K^3)^{793} \times \mathbb{R}_+^4)$

$\text{nb-L2 } (m = 123, n = 4195, K = K^{1677} \times (K^3)^{838} \times \mathbb{R}_+^4)$

$\text{nb-L2-bessel } (m = 123, n = 2641, K = K^{123} \times (K^3)^{838} \times \mathbb{R}_+^4)$

Compare iters/cpu(sec)/accuracy with Sedumi 1.05 ([Sturm '01](#)), which implements a predictor-corrector interior-point method.

Problem Name	SeDuMi (pars.eps=1e-5) iter/cpu	L-BFGS-Chol (accur=1e-5) iter/cpu
nb	19/7.6	1042/16.5
nb-L2	11/11.1	330/9.2
nb-L2-bessel	11/5.3	108/1.7

Table 1: (cpu times are in sec on an HP DL360 workstation, running Matlab 6.1)

Regularized Sum-of-Norms Problems

$$\min_{w \geq 0} \sum_{i=1}^M \|A_i w - b_i\|_2 + h(w),$$

$A_i \sim U[-1, 1]^{m_i \times \ell}$, $b_i \sim U[-5, 5]^{m_i}$, $m_i \sim U\{2, 3, \dots, r\}$ ($r \geq 2$).

$$h(w) = 1^T w + \frac{1}{3}\|w\|_3^3 \quad (\text{cubic reg.})$$

Reformulate as a convex SOCP:

$$\begin{aligned} & \text{minimize} && \sum_{i=1}^M z_i + h(w) \\ & \text{subject to} && A_i w + s_i = b_i, \quad (z_i, s_i) \in K^{m_i+1}, \quad i=1, \dots, M, \quad w \in \mathbb{R}_+^\ell. \end{aligned}$$

Problem ℓ, M, r (m, n)	BFGS-Chol iter/cpu	CG-PR-Chol iter/cpu	L-BFGS-Chol iter/cpu
500,10,10 (56,566)	352/24.6	1703/6.6	497/2.4
500,50,10 (283,833)	546/85.1	3173/69.0	700/12.4
500,10,50 (246,756)	272/36.3	1290/23.0	371/5.6

Table 2: (cpu times are in sec on an HP DL360 workstation, running Matlab 6.5.1, with accur=1e-3)

Smoothing Newton Step

$$\phi_{\text{FB}}^{\mu}(x, y) \stackrel{\text{def}}{=} (x^2 + y^2 + \mu^2 e)^{1/2} - x - y$$

with $e = (\underbrace{1, 0, \dots, 0}_{n_1}, \dots, \underbrace{1, 0, \dots, 0}_{n_p})^T, \quad \mu > 0$ (Fukushima,Luo,T '02)

Given ζ , choose $\mu > 0$ and solve

$$\nabla \phi_{\text{FB}}^{\mu}(F(\zeta), G(\zeta))^T \Delta \zeta = -\phi_{\text{FB}}(F(\zeta), G(\zeta))$$

Use $\Delta \zeta$ to accelerate convergence.

This requires more work per iteration. Use it judiciously.

Observations

For our unconstrained smooth merit function approach:

Advantage:

- Less work/iteration, simpler matrix computation than interior-point methods.
- Applicable to convex SOCP and column-monotone SOCCP.
- Useful for warm start?

Drawback:

- Many more iters. than interior-point methods.
- Lower solution accuracy.

SOCP from Dist. Geometry Optim

n pts in \mathbb{R}^d ($d = 2, 3$).

Know x_{m+1}, \dots, x_n and Eucl. dist. estimate for pairs of ‘neighboring’ pts

$$d_{ij} > 0 \quad \forall (i, j) \in \mathcal{A} \subseteq \{1, \dots, n\} \times \{1, \dots, n\}.$$

Estimate x_1, \dots, x_m .

Problem (nonconvex):

$$\min_{x_1, \dots, x_m} \sum_{(i,j) \in \mathcal{A}} |\|x_i - x_j\|_2^2 - d_{ij}^2|$$

- Objective function is nonconvex. 

- Problem is NP-hard (reduction from PARTITION). 
- Use a convex (SDP, SOCP) relaxation. SOCP is much easier to solve than SDP.

SOCP Relaxation

$$\begin{aligned}
 v_{\text{opt}} = & \min_{x_1, \dots, x_m, y_{ij}} \sum_{(i,j) \in \mathcal{A}} |y_{ij} - d_{ij}^2| \\
 \text{s.t. } & y_{ij} = \|x_i - x_j\|_2^2 \quad \forall (i, j) \in \mathcal{A}
 \end{aligned}$$

Relax “=” to “ \geq ” constraint:

$$\begin{aligned}
 v_{\text{socp}} = & \min_{x_1, \dots, x_m, y_{ij}} \sum_{(i,j) \in \mathcal{A}} |y_{ij} - d_{ij}^2| \\
 \text{s.t. } & y_{ij} \geq \|x_i - x_j\|_2^2 \quad \forall (i, j) \in \mathcal{A}
 \end{aligned}$$

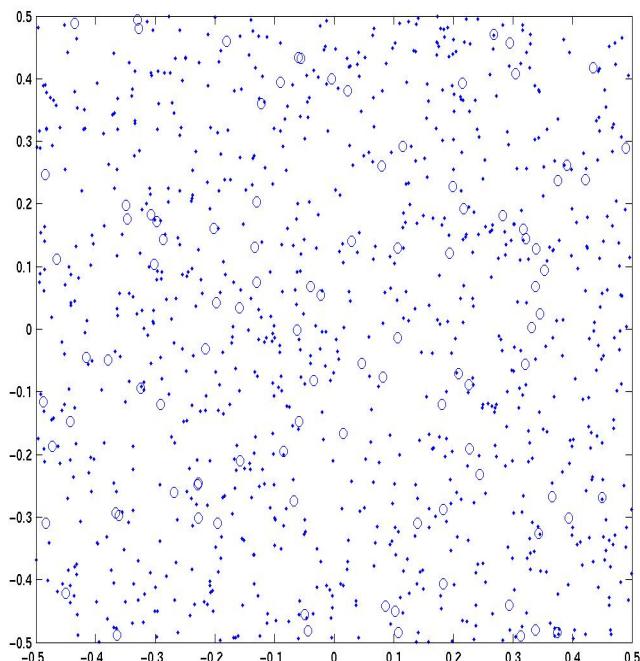
$$y \geq \|x\|_2^2 \iff y + 1 \geq \|(y - 1, 2x)\|_2$$

(also Doherty,Pister,El Ghaoui '03)

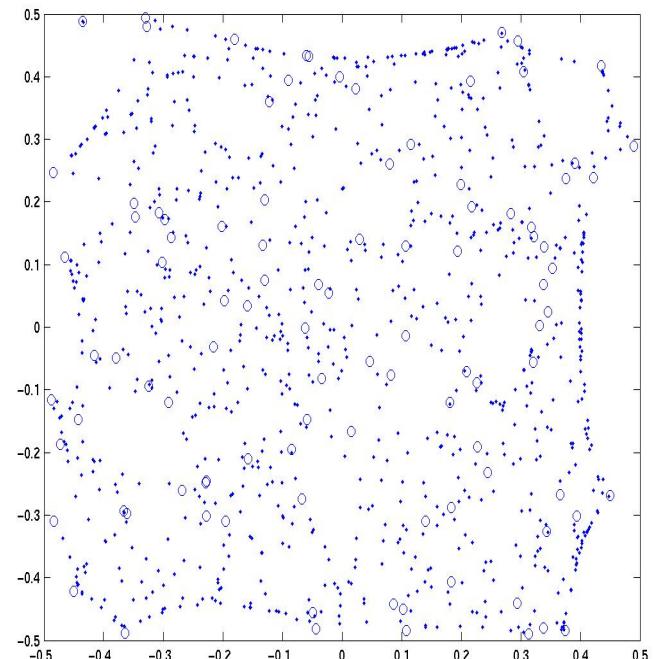
Simulation Results

Uniformly generate $\tilde{x}_1, \dots, \tilde{x}_n$ in $[-.5, .5]^2$, $m = 0.9n$

\tilde{x}_i and \tilde{x}_j are neighbor if $\text{dist} < .06$. Set $d_{ij} = \|\tilde{x}_i - \tilde{x}_j\|_2$ (Biswas, Ye '03)



True soln
($m = 900, n = 1000$)



SOCP soln found by IP method or
smoothing method

Lasty...

Thank you for coming!

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