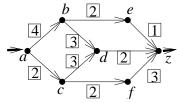
[Homework can be handed in to me or to my mail box in the Math Lounge (opposite the Math main office). Please show your work to receive full credit.]

**A.** Consider a digraph G = (V, A) and two vertices  $s \neq t \in V$ . To hedge against possible arc failures, we wish to find the largest number of paths from s to t that are arc-disjoint (i.e., no two paths share any arc). Formulate this as a maximum flow problem. Explain your answer. [Hint: Use unit capacities.]

**B**. Consider the capacitated digraph shown below:



- (a) Apply the algorithm MAXFLOW to find a maximum flow and minimum cut from a to z.
- (b) For the a-z flow  $x_{ab} = x_{be} = x_{ez} = 1$ ,  $x_{ac} = x_{cd} = x_{dz} = 2$  (with zero flow on remaining arcs), find the residual digraph and an a-z augmenting path. Augment flow along the autmenting path to obtain a new a-z flow.

C. Suppose that  $(x_{uv})_{(u,v)\in A}$  is an s-t flow of value v and [S,T] is an s-t cut of capacity c (with  $T=V\setminus S$ ). Prove that v=c if and only if  $x_{uv}=c_{uv}$  for all  $(u,v)\in [S,T]$  and  $x_{vu}=0$  for all  $(v,u)\in [T,S]$ . [You can use the fact that  $v=\sum_{(u,v)\in [S,T]}x_{uv}-\sum_{(v,u)\in [T,S]}x_{vu}$  for any s-t flow  $(x_{uv})_{(u,v)\in A}$  of value v.]

## $\mathbf{D}$ .

- (a) Give an example of a digraph G = (V, A) with arc capacities of 1 and two vertices s, t such that  $|V| \le 4$  and the s-t flow of maximum value is not unique.
- (b) For your answer to (a). Find an s-t flow of maximum value whose arc flows are not all integer.
- (c) Give an example of a digraph G = (V, A) with arc capacities of 1 and two vertices s, t such that |V| = 3 and the s-t cut of minimum capacity is not unique.
- (d) Consider a digraph G = (V, A) with arc capacities  $c_{uv}$ ,  $(u, v) \in A$ , and  $s \neq t \in V$ . Suppose  $u_1, u_2, u_3, u_4$  is an s-t augmenting path relative to an s-t flow  $(x_{uv})_{(u,v)\in A}$ . Suppose  $(u_2, u_3)$  is a forward arc and  $(u_2, u_1), (u_4, u_3)$  are backward arcs. Write down a formula for the augmentation amount  $\Delta$  in terms of the flow and capacity of these arcs.

**Bonus.** Consider a digraph G=(V,A) with arc capacities  $c_{uv}$ ,  $(u,v) \in A$ , and  $s \neq t \in V$ . Suppose that  $[S_1,T_1]$  and  $[S_2,T_2]$  are two s-t cuts of minimum capacity. Prove that  $[S_1 \cap S_2,T_1 \cup T_2]$  and  $[S_1 \cup S_2,T_1 \cap T_2]$  are also s-t cuts of minimum capacity. [Hint: First show that they are s-t cuts. Then show that their capacities equal that of  $[S_1,T_1]$  and  $[S_2,T_2]$  by considering arcs out of  $S_1 \cap S_2$ ,  $S_1 \setminus S_2$ ,  $S_2 \setminus S_1$ .]