# **Near-Optimal Power Control in Wireless Networks: A Potential Game Approach**

# Motivation

- Due to the shared nature of the wireless domain, the performance of each mobile depends on the resources allocated to others
- The power control problem, even as a centralized optimization problem with full information, is a fairly complex problem.
- An additional concern: self-interested behavior of mobiles.
- We consider the power allocation problem from the viewpoint of a central planner who wishes to impose a certain power-dependent objective in the network, using pricing.
- Using a potential-game approach, we provide a general distributed power control scheme that would (approximately) achieve any underlying system objective despite the selfishness of the mobiles.

## **The Model**

- Set of mobiles  $\mathcal{M} = \{1, \ldots, M\}$  share the same wireless spectrum (single channel).
- The power allocation of the mobiles  $\mathbf{p}$  =  $(p_1,\ldots,p_M)$  (where  $0 < \underline{P}_m \leq p_m \leq P_m$ .)
- The rate of user m is given by:

$$r_m(\mathbf{p}) = \log\left(1 + \gamma \mathrm{SINR}_m(\mathbf{p})\right),$$

where,  $\gamma > 0$  is the spreading gain of the CDMA system and

$$\operatorname{SINR}_{m}(\mathbf{p}) = \frac{h_{mm}p_{m}}{N_{0} + \sum_{k \neq m} h_{km}p_{k}},$$

- The payoff of player m:  $u_m(\mathbf{p}) = r_m(\mathbf{p})$  $\lambda_m p_m.$
- $\lambda_m$  is set by the central planner, to maximize a system utility-function  $U_0(\cdot)$ .

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Summary



such that

$$\Phi(x_m, x_{-m}) - \Phi(y_m, x_{-m}) = u_m(x_m, x_{-m})$$

Let  $\tilde{r}_m(\mathbf{p}) = \log(\gamma \text{SINR}_m(\mathbf{p}))$ . Consider a related game  $\tilde{\mathcal{G}}$  with utilities  $\tilde{u}_m(\mathbf{p}) = \tilde{r}_m(\mathbf{p}) - \lambda_m p_m$ . Note that  $\tilde{u}_m(\mathbf{p}) \approx u_m(\mathbf{p})$ . for  $\gamma \gg 1$  or  $h_{mm} \gg h_{km}$  for all  $k \neq m$ .

- $\tilde{\mathcal{G}}$  is a potential game with potential  $\phi(\mathbf{p}) = \sum_{m} \log(p_m) \lambda_m p_m$ . Additionally,  $\tilde{\mathcal{G}}$  has a unique equilibrium.
- A simple linear pricing scheme can be used to set a desired operating point as an equilibrium in  $\hat{\mathcal{G}}$ .

**Theorem 1.** Given a desired operating point  $\mathbf{p}^*$ , the unique equilibrium of  $\mathcal{G}$  is  $\mathbf{p}^*$  if the prices  $\lambda^*$  are chosen as

$$\lambda_m^* = \frac{1}{p_m^*}, \quad \text{for all } m$$

BR dynamics:  $\dot{p}_m = \beta_m(\mathbf{p}_{-m}) - p_m$  for all  $m \in \mathcal{M}$ , where  $\beta_m(\mathbf{p}_{-m}) = \arg \max_{p_m \in \mathcal{P}_m} u_m(p_m, \mathbf{p}_{-m})$ . Our idea: Use perturbed system analysis to study dynamics and equilibria in  $\mathcal{G}$  using properties of  $\mathcal{G}$ .

 $(x_{-m}) - u_m(y_m, x_{-m}),$ 

 $f\in\mathcal{M}^{2}$ 

**Theorem 2.** In the original game, best-response dynamics converges uniformly to the set  $\mathcal{I}_{\epsilon}$ , where

The above characterization is used to bound the deviation from the optimal system utility. For ex-

Percent loss in the system utility with respect to change in  $\gamma$ 

- prices.

# **Analysis of Dynamics**

$$\epsilon \leq \frac{1}{\gamma} \sum_{m \in \mathcal{M}} \frac{1}{\underline{SINR}_m}.$$

### **Future Work**

Multichannel networks.

• Distributed methods for calculating the optimal

• Applications of the potential-game approach to additional wireless network domains.