

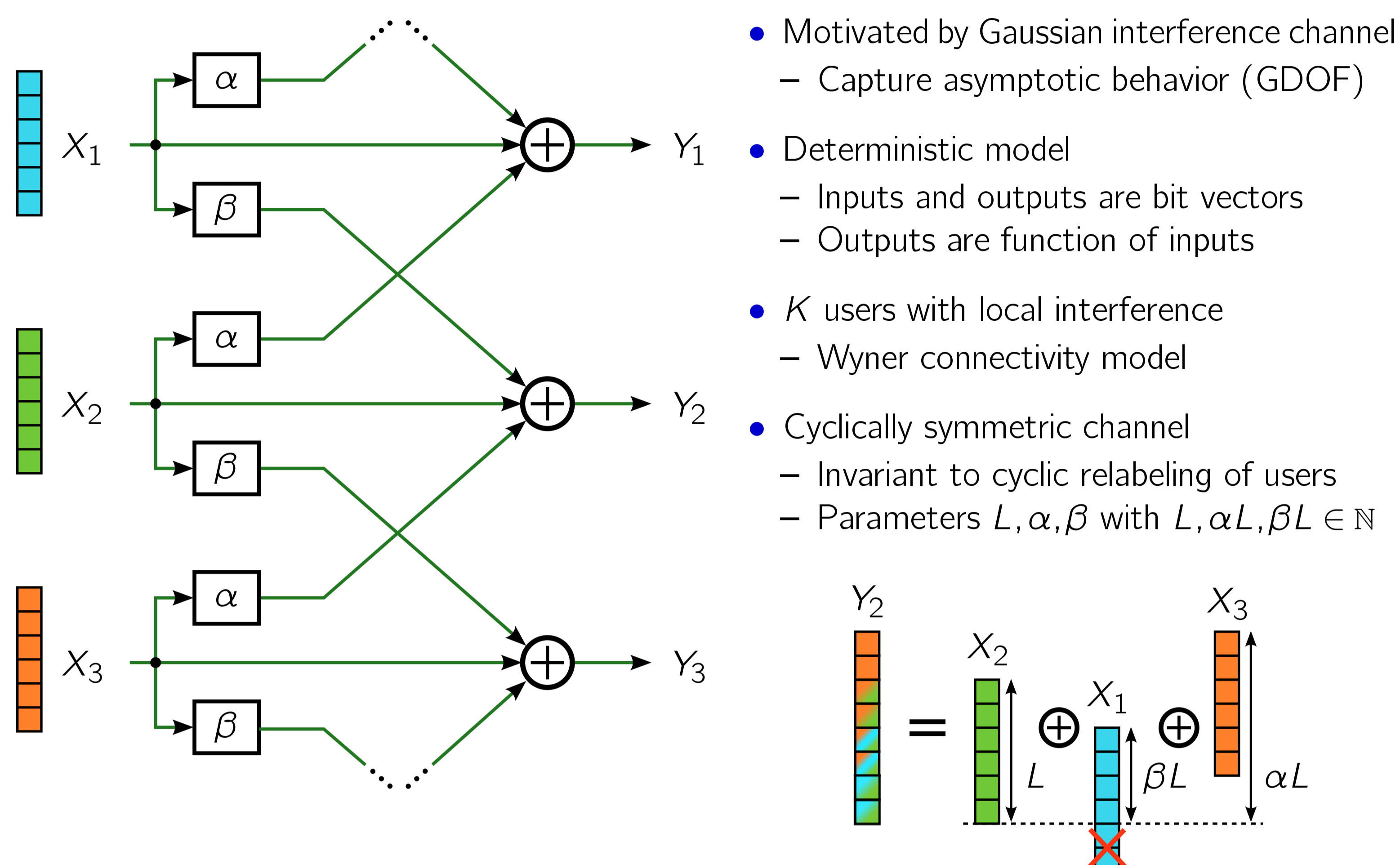
Sum capacity of a Class of Cyclically Symmetric Interference Channels

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Abstract

- Consider a K -user deterministic interference channel
- Find the symmetric capacity (sum capacity) for a large set of parameters
- Capacity is achieved by intricate bit pipe assignment patterns

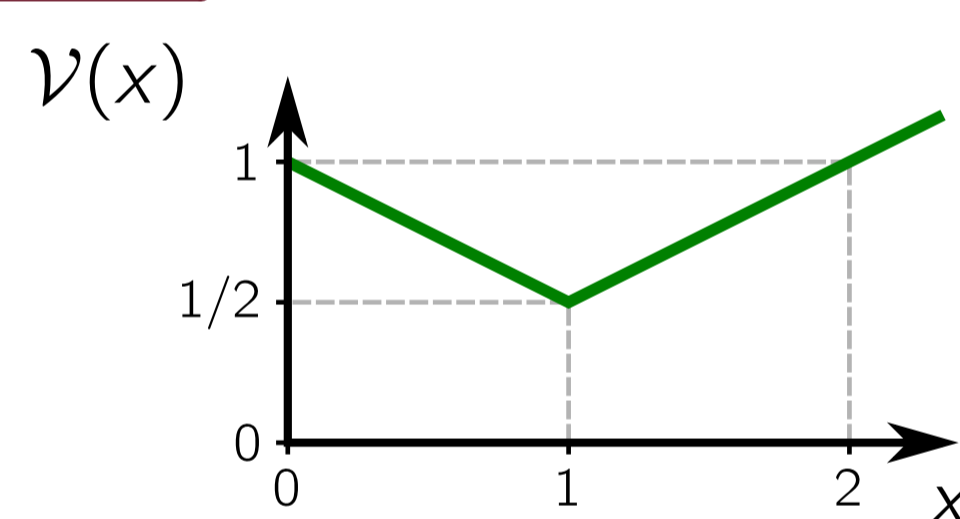
Cyclically symmetric interference channel



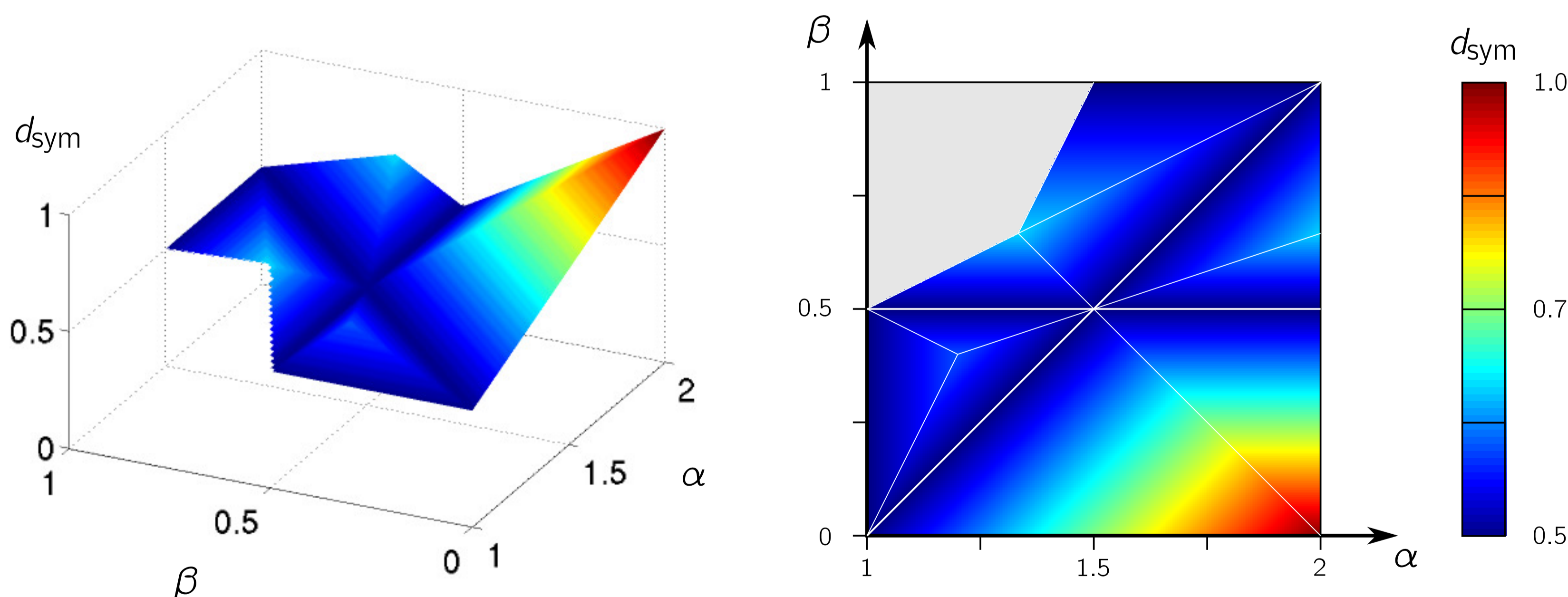
$$\mathbf{Z} = \begin{bmatrix} \mathbf{0}_{L \times L} \\ \mathbf{I}_L \end{bmatrix}, \quad \mathbf{D} = \mathbf{U}^T = \begin{bmatrix} \mathbf{0}^T & \mathbf{0} \\ \mathbf{I}_{2L-1} & \mathbf{0} \end{bmatrix}, \quad \mathbf{Y}_2 = \mathbf{Z}\mathbf{X}_2 \oplus \mathbf{U}^{(\alpha-1)L}\mathbf{Z}\mathbf{X}_3 \oplus \mathbf{D}^{(1-\beta)L}\mathbf{Z}\mathbf{X}_1$$

Main result

- Normalized symmetric capacity $d_{\text{sym}} = \sup R_{\text{sym}}/L$
- Define $\mathcal{V}(x) = \frac{1+|x-1|}{2}$

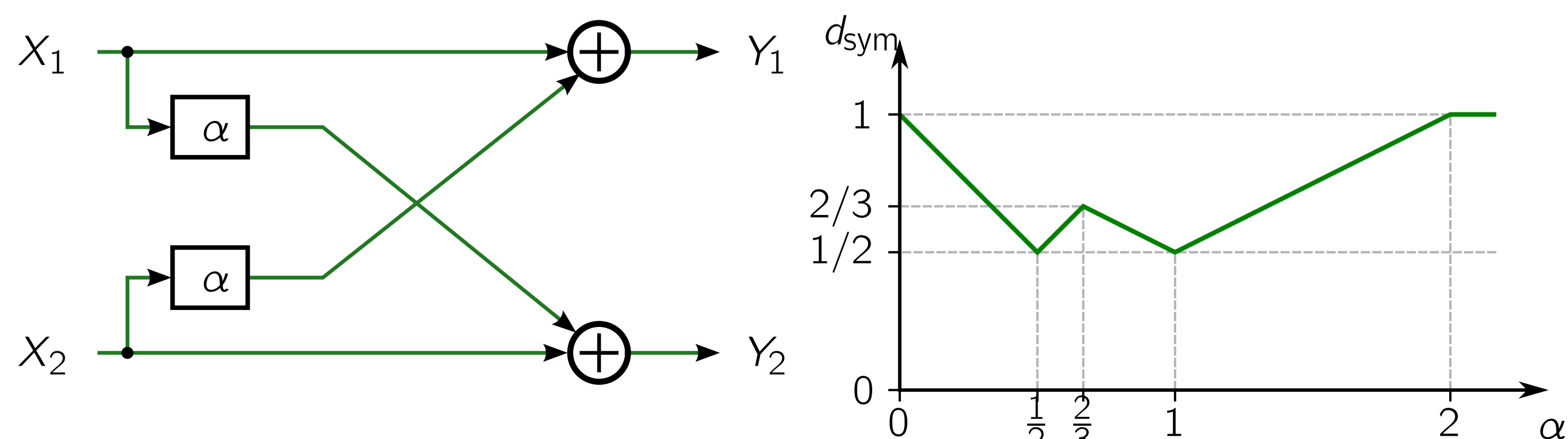


The normalized symmetric capacity of the channel defined above with $(\alpha, \beta) \in [1, 2] \times [0, 1]$ and $\alpha \geq \min\{2\beta, \beta/2 + 1\}$ is

$$d_{\text{sym}} = \min\{1, \mathcal{V}(\alpha), \mathcal{V}(\beta), \mathcal{V}(2\beta), \mathcal{V}(\alpha - \beta)\}$$


Comparison to the two-user case

- Symmetric capacity is known [EC82, ADT07, BT08]
- Achievable by single-letter coding scheme [JV08], which we extend here



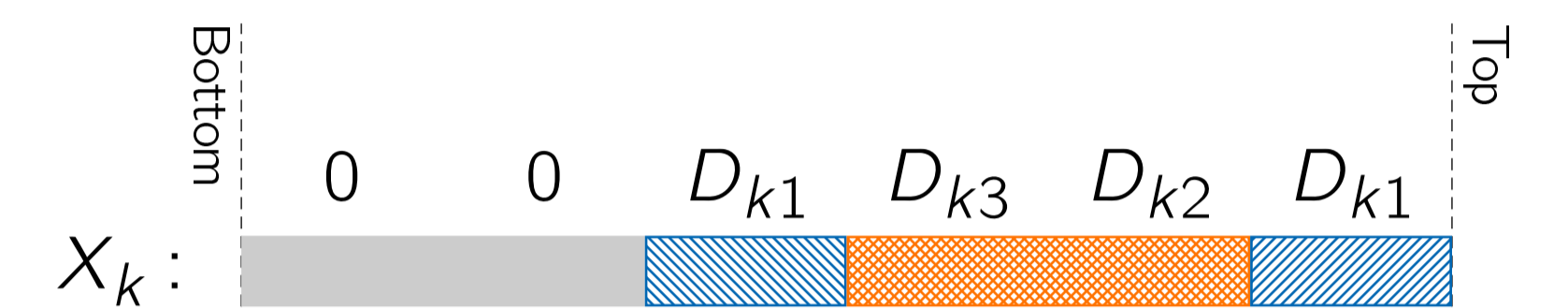
Capacity-achieving codes

- Characteristics of the proposed coding scheme:
 - Single-letter
 - Linear encoding and decoding
 - All transmitters use the same code
- For given (α, β) , encode as follows:

Let D_k : independent message bits
 $\mathbf{G}(\alpha, \beta)$: assignment matrix of size $L \times L$ $d_{\text{sym}}(\alpha, \beta)$
Use $X_k = \mathbf{G}(\alpha, \beta)D_k, \quad k = 1, 2, 3$

- Example:

$$X_k = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} D_{k1} \\ D_{k2} \\ D_{k3} \end{bmatrix}$$

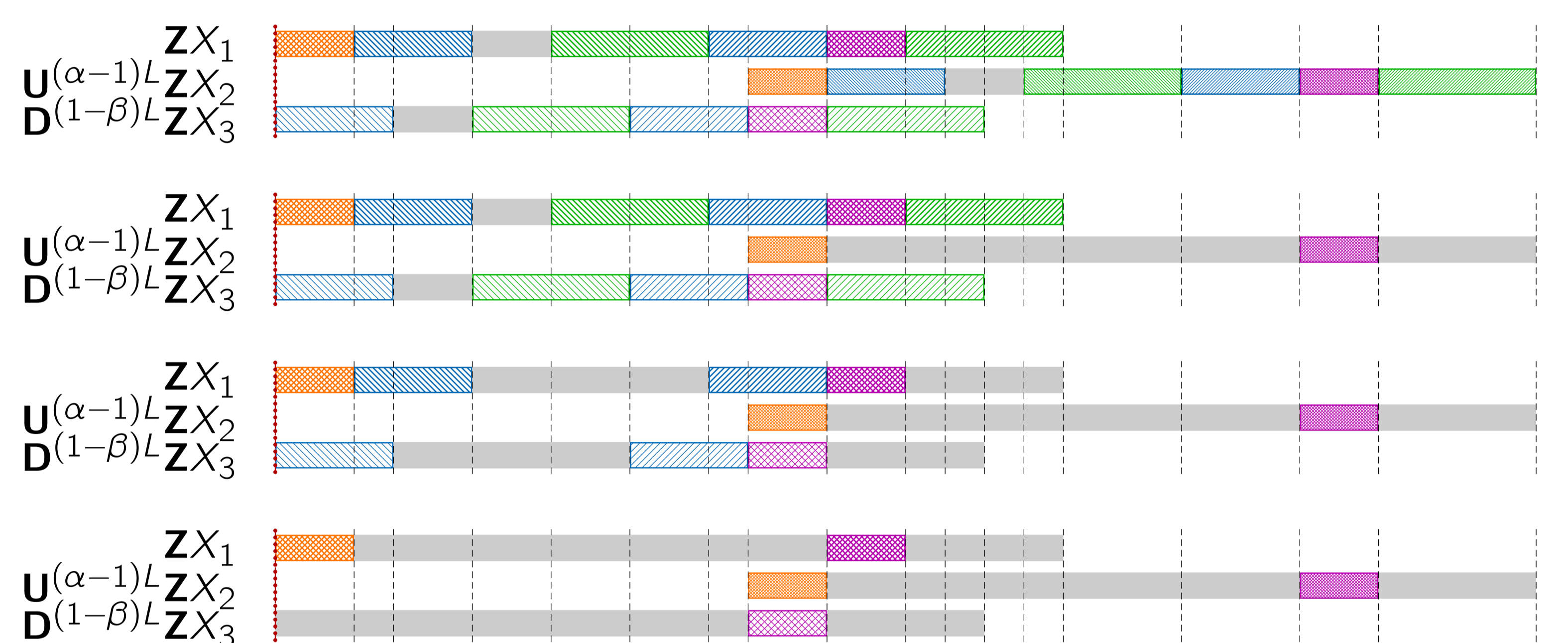


Example

- Consider $\alpha = 1.6, \beta = 0.9$, with optimum $d_{\text{sym}} = 0.55$

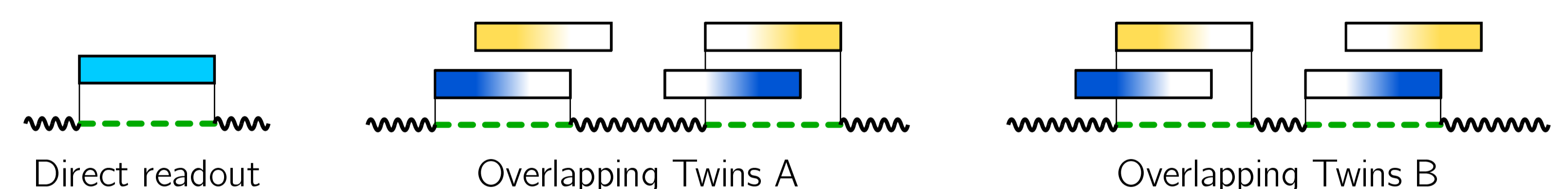
$$X_k: \text{[Bit assignment diagram for } \alpha=1.6, \beta=0.9 \text{]}$$

- Receiver signal $Y_1 = \mathbf{Z}\mathbf{X}_1 \oplus \mathbf{U}^{(\alpha-1)L}\mathbf{Z}\mathbf{X}_2 \oplus \mathbf{D}^{(1-\beta)L}\mathbf{Z}\mathbf{X}_3$:

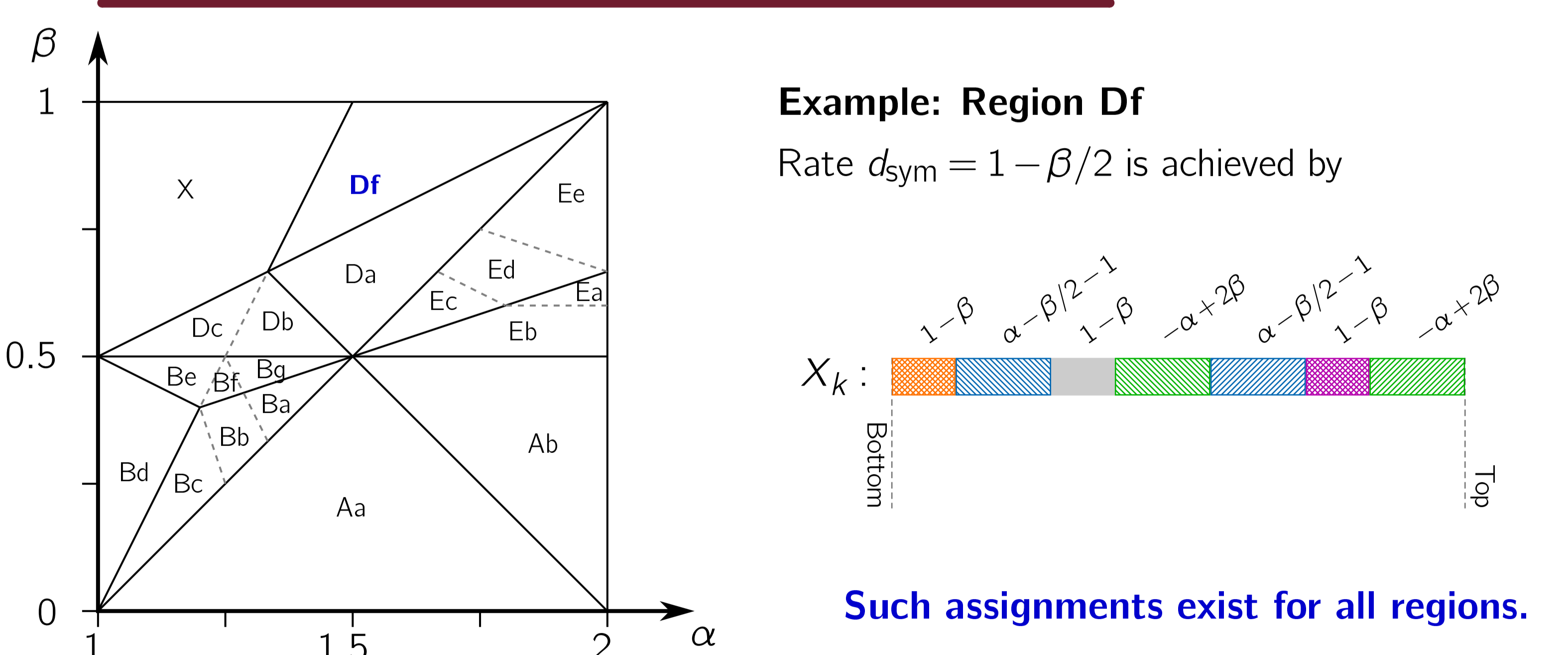


Decoding rules

- Each step above matches one of three basic rules



Different patterns for different regions



References

- [ADT07] A. S. Avestimehr, S. N. Diggavi, and D. N. C. Tse. "A deterministic approach to wireless relay networks". In "Proc. Allerton", Sep. 2007.
- [BT08] G. Bresler and D. Tse. "The two-user Gaussian interference channel: a deterministic view". *Euro. Trans. Telecomm.*, vol. 19, no. 4, pp. 333–354, Jun. 2008. ArXiv:0807.3222.
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- [JV08] S. A. Jafar and S. Vishwanath. "Generalized degrees of freedom of the symmetric K user Gaussian interference channel". 2008. arXiv:0804.4489.