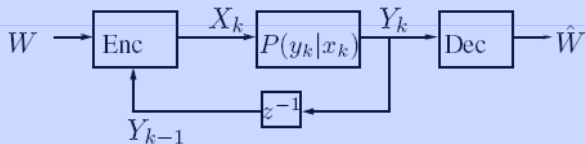


STATUS QUO

- The use of feedback is of the utmost importance in designing **scalable, robust, reliable** communication schemes.
- Posterior Matching Scheme generalizes previous feedback communication schemes for any memoryless channel. [1][2]
- A deep understanding of feedback is still lacking. Are there simple, provably good iterative feedback encoders/decoders?

## THE POSTERIOR MATCHING SCHEME

• A real number  $W_0$  in  $[0,1]$  represents an infinite sequence of bits to be communicated across a channel  $P(y_k|x_k)$  with noiseless feedback.



• The information still missing at the receiver is extracted from the a-posteriori density function.

$$W_n = F_{W|Y^{n-1}}(W_0 | Y_1^{n-1})$$

• The signal is then reshaped to capacity-achieving (optimal) distributions.

$$X_n = F_Q^{-1}(W_n)$$

## RECURSIVE FORM OF THE PM Scheme

$$X_1 = F_Q^{-1}(W_0)$$

$$X_{n+1} = S(X_n, Y_n)$$

- **Converse** to coding theorem: next input should be independent of everything decoder has seen so far
- With an **optimal** decoder, such "posterior matching" schemes achieve capacity
- Does this motivate a **simple iterative decoder**, that **achieves** capacity?
- How do we analyze this easily, exploiting the dynamics?

## MAIN RESULT

For the PM scheme, under appropriate technical assumptions, the following dynamical system decoder achieves capacity:

- Guess the state at time n to be  $u \sim \text{Unif}(0,1)$

$$\hat{W}[n | n] = u$$

- Then recursively solve previous states

$$\hat{W}[i | n] = S_{Y_i}^{-1}(\hat{W}[i+1 | n])$$

- Finally, we arrive at our estimate of  $W_0$

$$\hat{W}_0 = \hat{W}[0 | n]$$

As long as our estimate is within  $2^{-(nC+1)}$  of our actual message, we can decode nC bits.

Since the  $Y$ 's are **independent** and identically distributed, the decoder dynamical system is a **Markov chain** and has a **Lyapunov exponent** of  $-C$  for any initial condition  $u$  on  $(0,1)$ .

## REFERENCES

- [1] O. Shayevitz and M. Feder, "Communication with feedback via posterior matching," in *IEEE International Symposium on Information Theory*, Nice, France, June 2007.
- [2] —, "The posterior matching feedback scheme: Capacity achieving and error analysis," in *IEEE International Symposium on Information Theory*, Toronto, Canada, July 2008.
- [3] H. Ebeid and T. P. Coleman, "Lyapunov Exponents and Reverse Iterated Function System Decoders for Channels with Feedback", in preparation for submission, IEEE International Conference on Communications, September 2009.
- [4] T. P. Coleman "A Stochastic Control Approach to 'Posterior Matching'-style Feedback Communication Systems", submitted to *IEEE International Symposium on Information Theory*, Seoul, Korea, July 2009.

## BINARY SYMETRIC CHANNEL

Consider the binary symmetric channel with error probability  $p$ :

• Figure 1 to the right shows the trajectories of the decoder when started with different values of  $W_{n+1}$ .

• The recursive transmission functions used for encoding and decoding are shown in Figure 1 below.

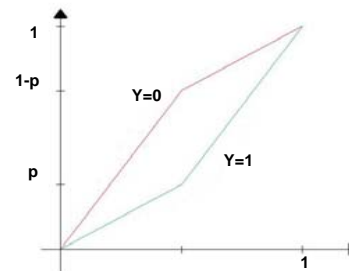


Figure 1. Recursive transmission functions

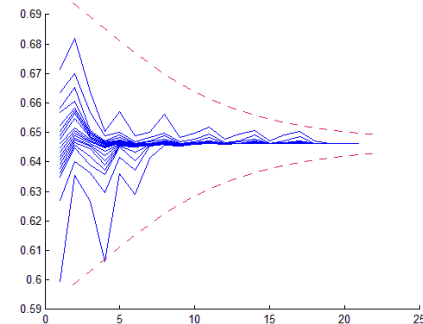


Figure 2. Trajectories of decoder with different initial condition

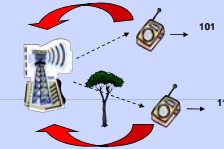
## ASSUMPTIONS AND LIMITATIONS:

- Noiseless feedback
- Memoryless Channels

• **Side benefit:** provides a conceptually simple understanding of how PM scheme uses feedback to achieves capacity

## IMPACT

- Provides **explicit** capacity-achieving **recursive encoders and decoders**
- Can be extended to **networks** with tight converses



## FUTURE GOALS

Use Stochastic Control methodology for a **principled, canonical approach** to address:

- noisy feedback (POMDP)
- Unknown channel (Q-learning)
- Delayed feedback