

Distributed Scheduling and Equilibrium Dynamics in Wireless Networks with Correlated Fading Channels

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Introduction

Centralized control schemes for wireless networks (e.g., wireless ad-hoc networks) could be practically difficult to implement:

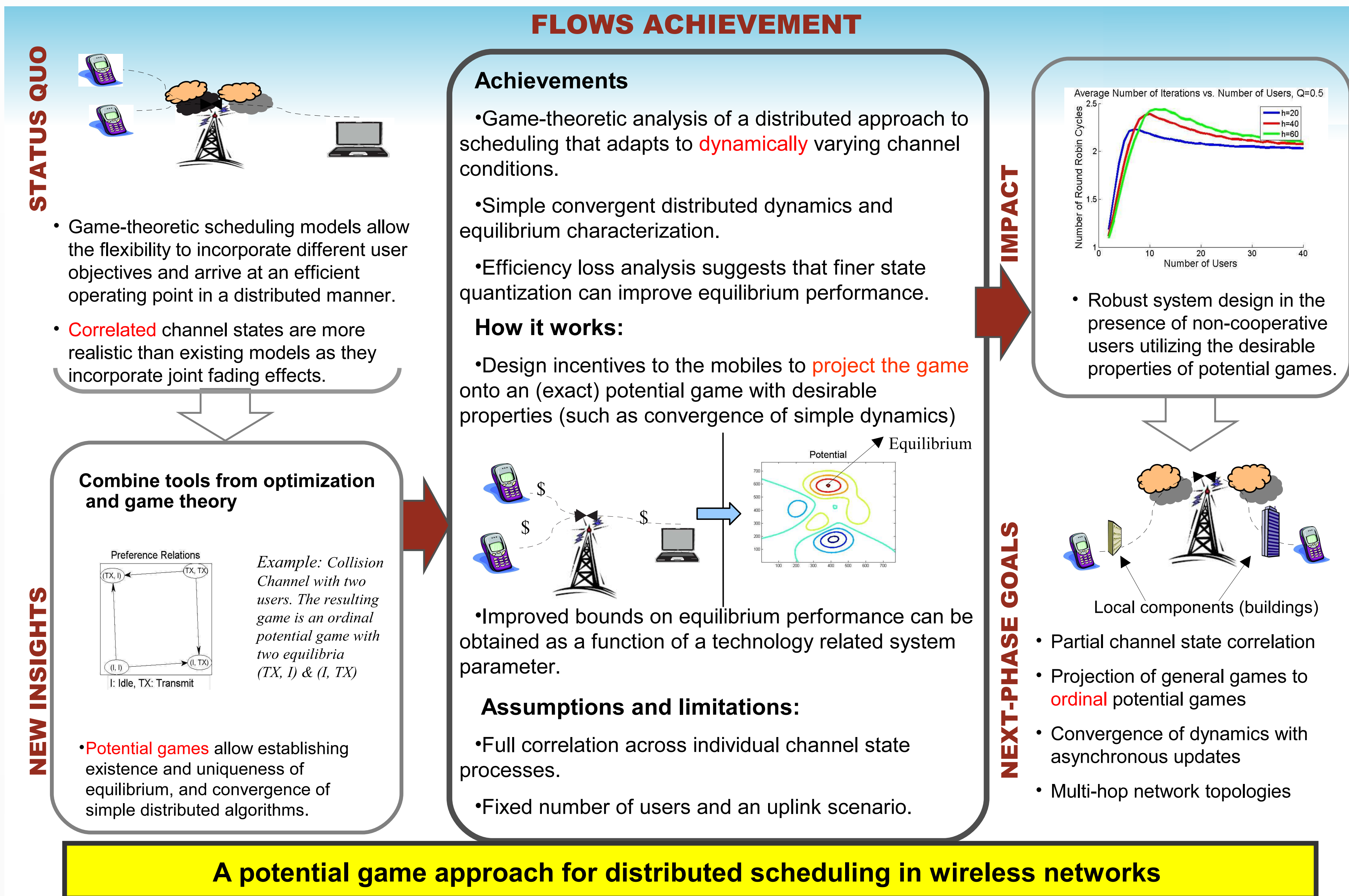
- A central entity may not have access to the required information on the end users, especially in dynamically evolving systems.
- Centralized optimization procedures might be computationally involved.

Distributed approaches are more natural. Competition for network resources can be modeled as a noncooperative **game** leading to a robust and distributed control paradigm.

Model

- A finite set of mobiles, $\mathcal{M} = \{1, \dots, M\}$.
- **Fading** effect: A global channel state process $H(t) \in \mathcal{H} = (1, 2, \dots, h)$. Example of global effects may include common weather conditions in satellite communications, and thermal noise at the base station.
- The state process is stationary (e.g., block-fading), state i is observed with probability π_i .
- Collision channel: Simultaneous transmissions collide and data is lost.
- The average rate user m can sustain in state i (assuming no collision) is denoted by R_i^m .
- Per-user average power constraint \bar{P}^m .
- We consider stationary strategies: p_i^m is the (stationary) transmission probability of user m at state i .
- Utilities capture tradeoff between power (P^m) and throughput (T^m): $u^m(\mathbf{p}^m) = T^m(\mathbf{p}^m, \mathbf{p}^{-m}) - \lambda^m P^m(\mathbf{p}^m)$

Summary



Main Results

- Characterization of the social welfare problem, and useful reduction of our game to a finite game.
- Let $\pi_{max} = \max_{i \in \mathcal{H}} \pi_i$, $P_{min} = \min_{m \in \mathcal{M}} \bar{P}^m$, and define a technology-related parameter: $Q = \frac{\pi_{max}}{P_{min}}$.

Theorem 1. Fix $Q < 1$. Then,

- The performance ratio between the best equilibrium and the social optimum is bounded above by $(1-Q)^{-2}$.
- The performance of the worse equilibrium point could be arbitrarily bad.

- Implication: With finer quantization, equilibrium efficiency can be improved.

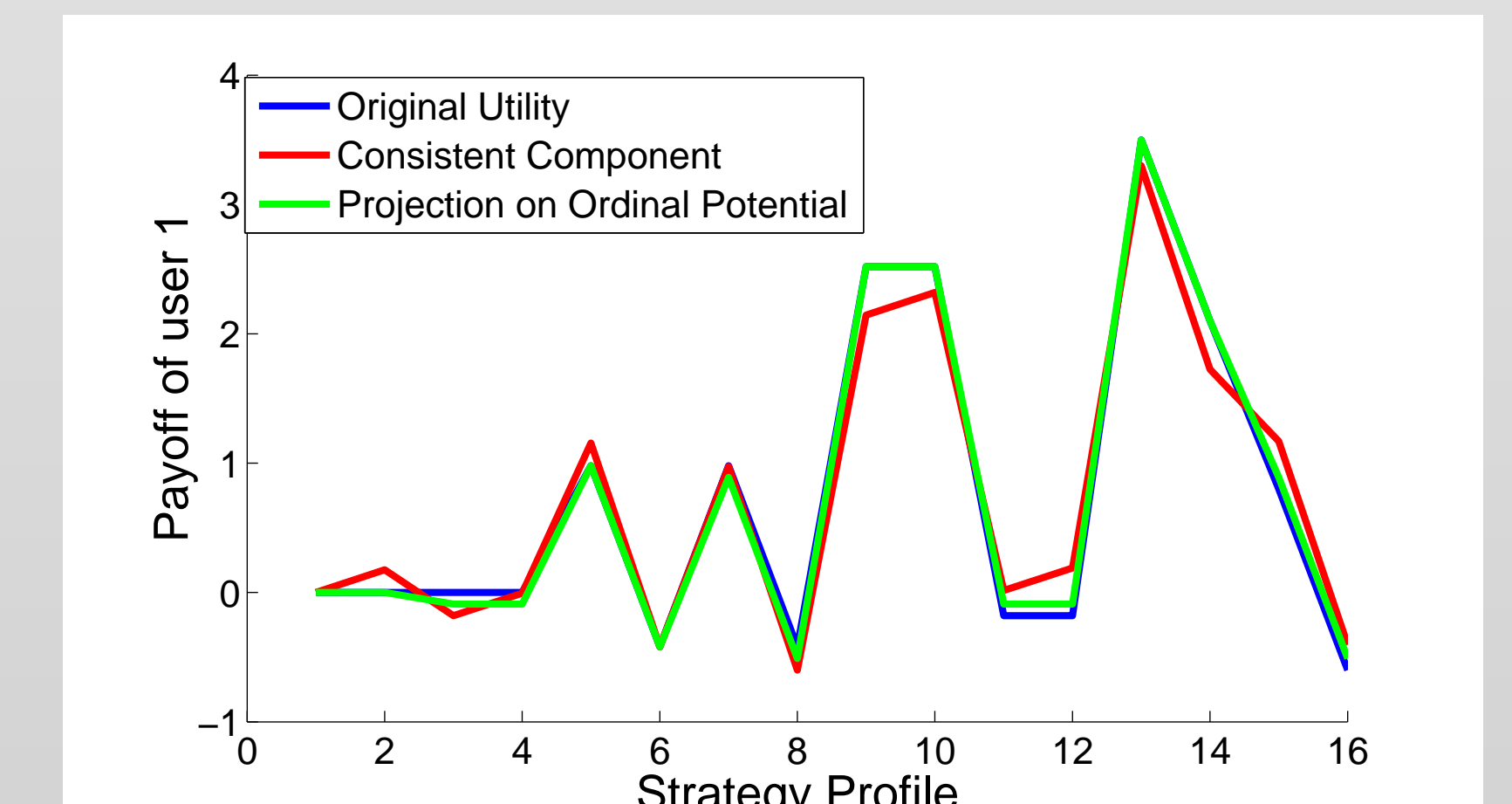
Theorem 2. If $R_i^m = R_i$ for all users m , then the game is a **potential** game, which in our case implies convergence of best-response mechanism to an equilibrium point in **finite** time.

- We further show that under general rate values R_i^m the game is not a potential game.

Potential Games

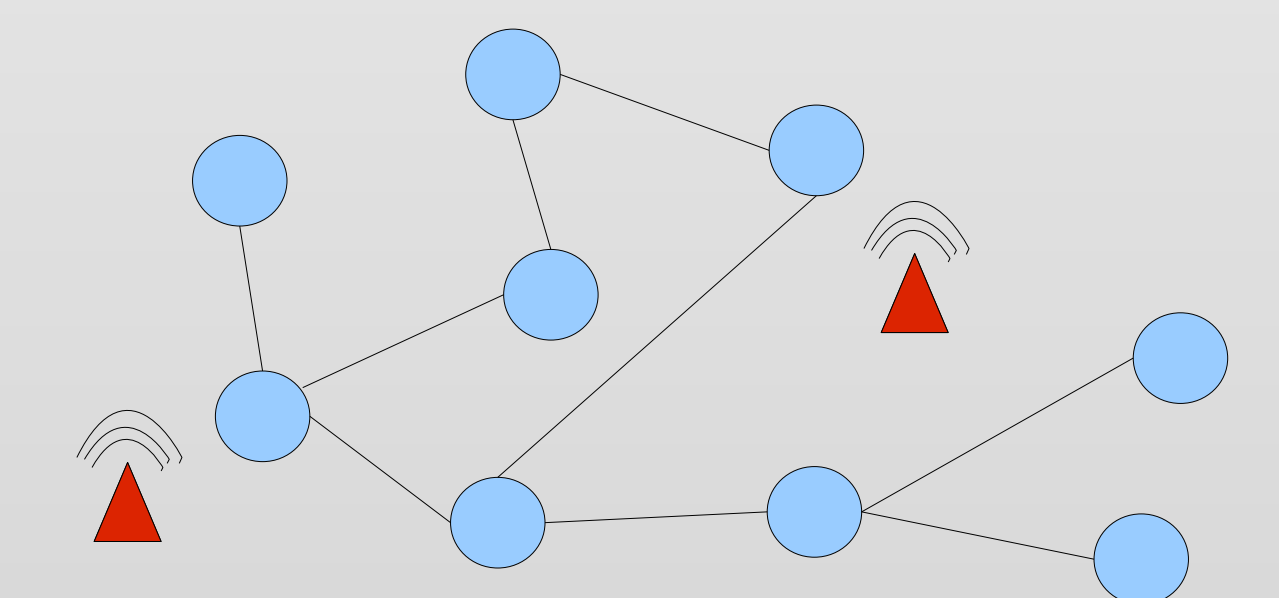
Potential games have desirable convergence properties.

- Thus, in order to guarantee convergence of simple myopic dynamics, a game can be **projected** onto a potential game.
- Ongoing research focuses on a general framework, in which the original user utilities are slightly modified (e.g., in the form of incentives) to form a potential game.



Future Work

- Exploit the idea of projection onto a potential game in our framework.
- Partial correlation of the state processes.
- Multi-hop architectures.



- Additional channel models (e.g., CDMA)
- Convergence of dynamics with asynchronous updates