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# Cooperation and Coding in MANETs

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# Vision

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- λ To develop new capacity definitions for the appropriate MANET metrics and models.
- λ To find novel communication techniques that expand the associated achievable rate regions.
- λ To exploit network degrees of freedom optimally
- λ To integrate security and robustness into the information theory of MANETs

# New Theory

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- λ New capacity definitions
- λ Scaling laws that make sense
- λ Adaptive coding over degrees of freedom
- λ Cooperation, relaying and conferencing
- λ Error exponents
- λ Mismatch detection
- λ Capacity with security and robustness

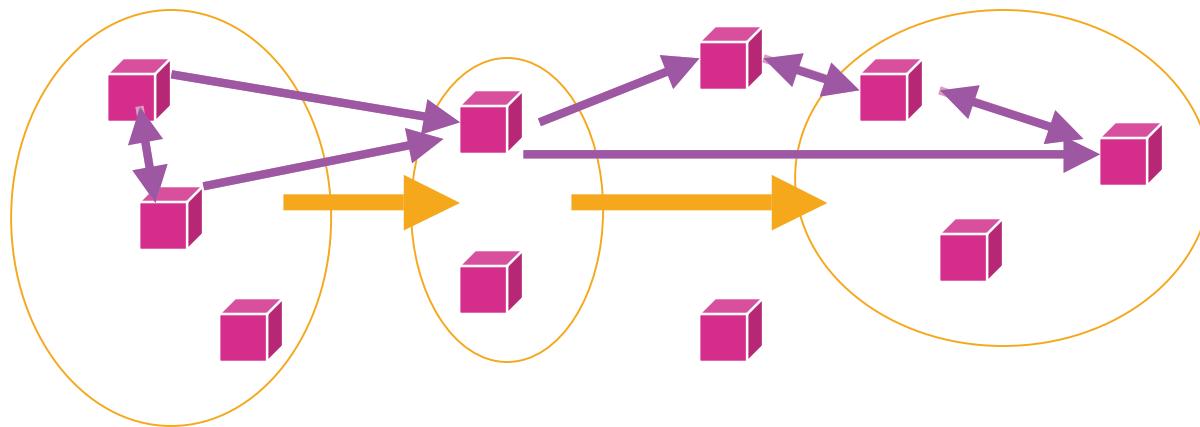
# New Capacity Definitions

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- λ New capacity definitions needed that allow for dynamics, errors, delay, etc.
  - λ Capacity with outage
  - λ Expected capacity
- λ Coding theorems to achieve capacity under these definitions are fundamentally different.
- λ The network application(s) dictate the “right” capacity definition

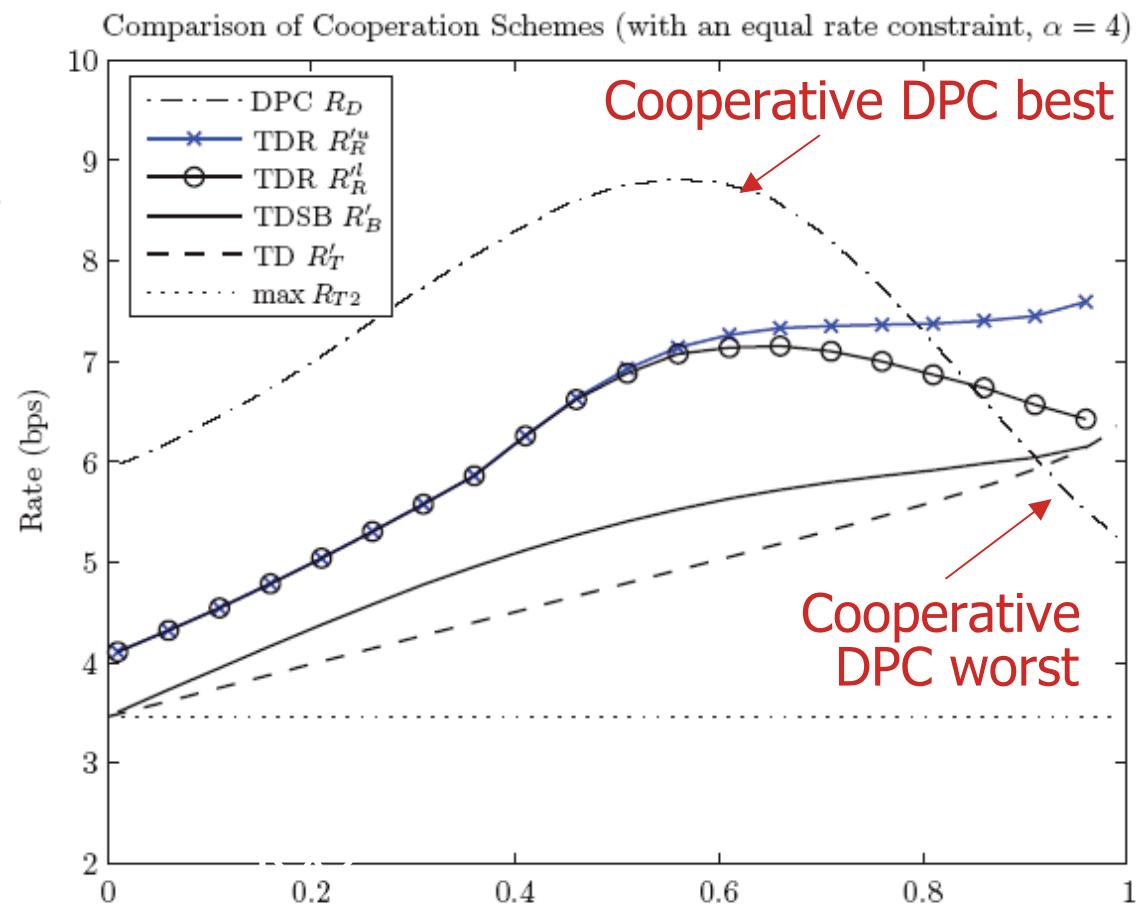
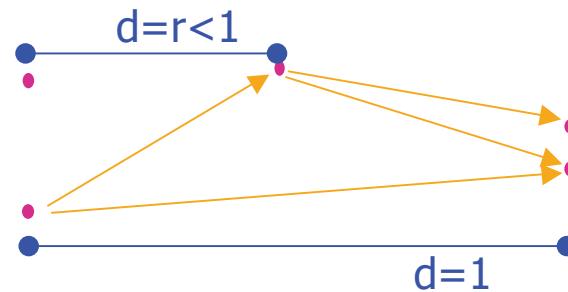
# Cooperation in Wireless Networks

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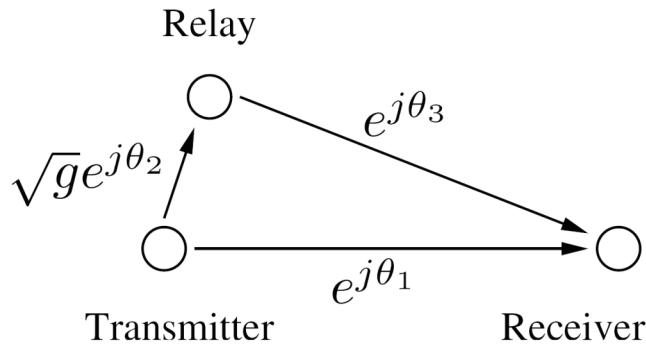
- λ Many possible cooperation strategies:
  - λ Virtual MIMO with DPC, relaying (DF, CF, AF), one-shot/iterative conferencing, and network coding
  - λ Nodes can use orthogonal or non-orthogonal channels.
  - λ Forwarding/routing/network “coding” should not be based on a point-to-point model

# Virtual MIMO vs. Relaying

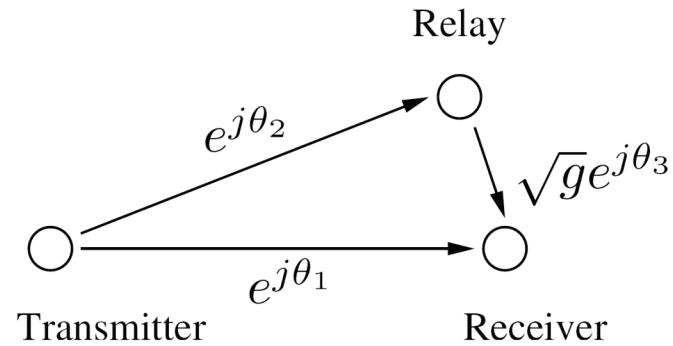


*Must quantify cost of synchronization and CSI*

# Relative Benefits of TX and RX Cooperation



(a) Transmitter cooperation

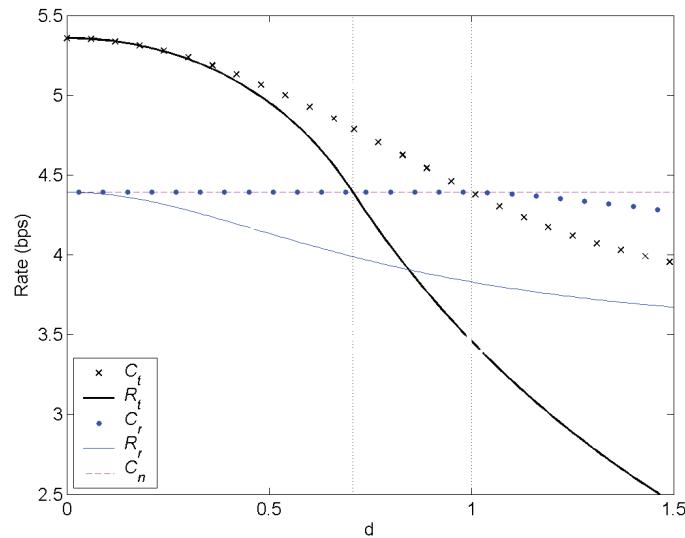


(b) Receiver cooperation

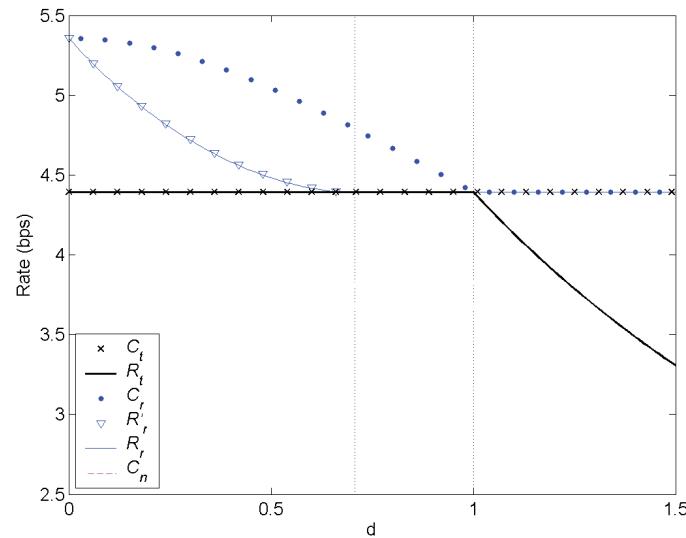
- λ Two possible CSI models:
  - λ Each node has full CSI (synchronization between Tx and relay).
  - λ Receiver phase CSI only (no TX-relay synchronization).
- λ Two possible power allocation models:
  - λ Optimal power allocation: Tx has power constraint  $aP$ , and relay  $(1-a)P$ ;  $0 \leq a \leq 1$  needs to be optimized.
  - λ Equal power allocation ( $a = 1/2$ ).

# Effective Cooperation

Equal Power Allocation with Full CSI

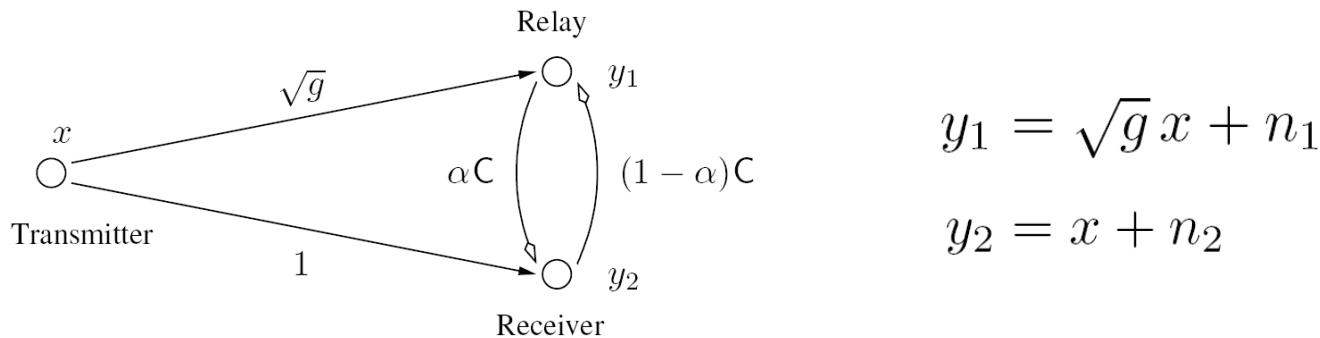


Optimal Power Allocation with Receiver CSI



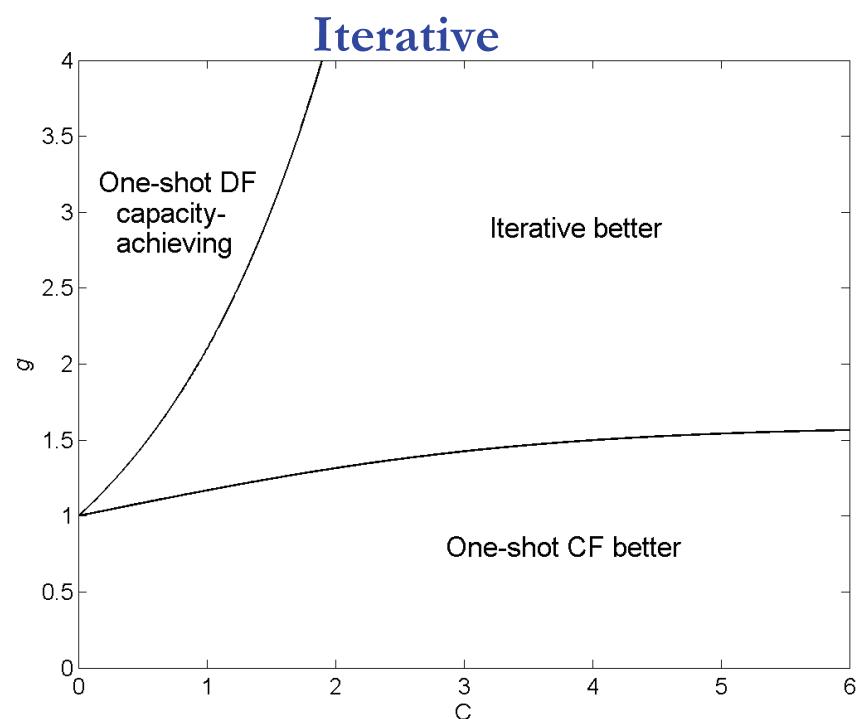
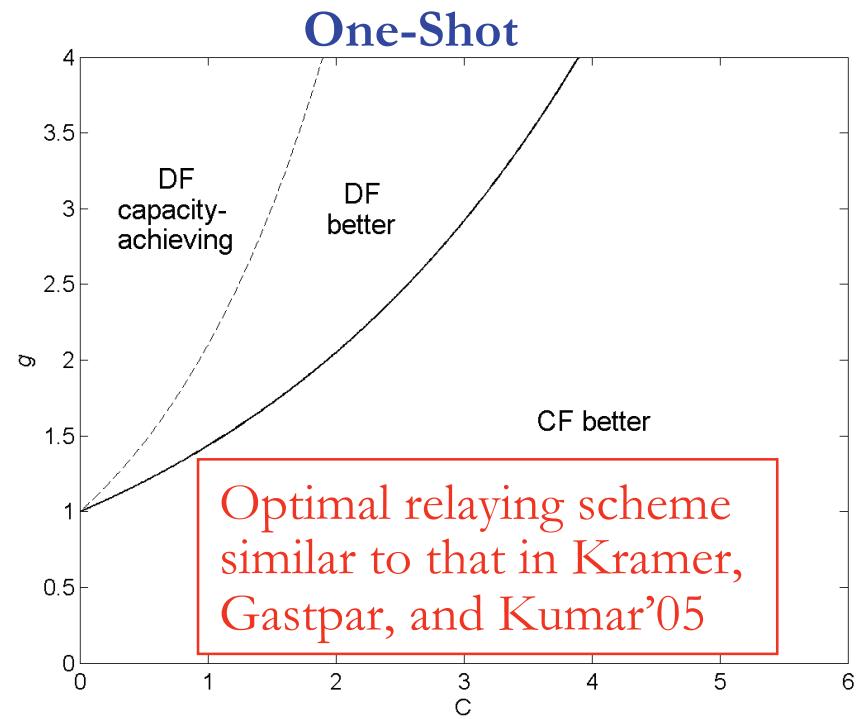
- λ Gain only realized with the right cooperation strategy
- λ With full CSI, Tx co-op is superior.
- λ With optimal power allocation and receiver phase CSI, Rx co-op is superior.
- λ With equal power and Rx phase CSI, no gain from cooperation.

# Conferencing Relay Channel



- λ Willems introduced conferencing for MAC (1983)
  - λ Transmitters conference before sending message
- λ We consider a relay channel with conferencing between the relay and destination
- λ The conferencing link has total capacity  $C$  which can be allocated between the two directions

# Iterative vs. One-shot Conferencing



- λ Weak relay channel: the iterative scheme is disadvantageous.
- λ Strong relay channel: iterative outperforms one-shot conferencing for large  $C$ .

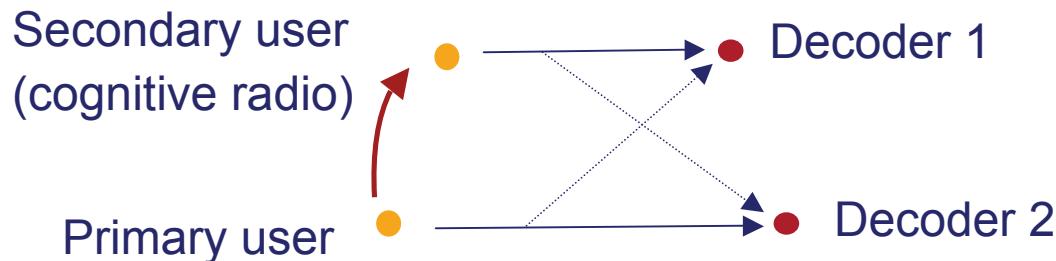
# Cognitive radio cooperation

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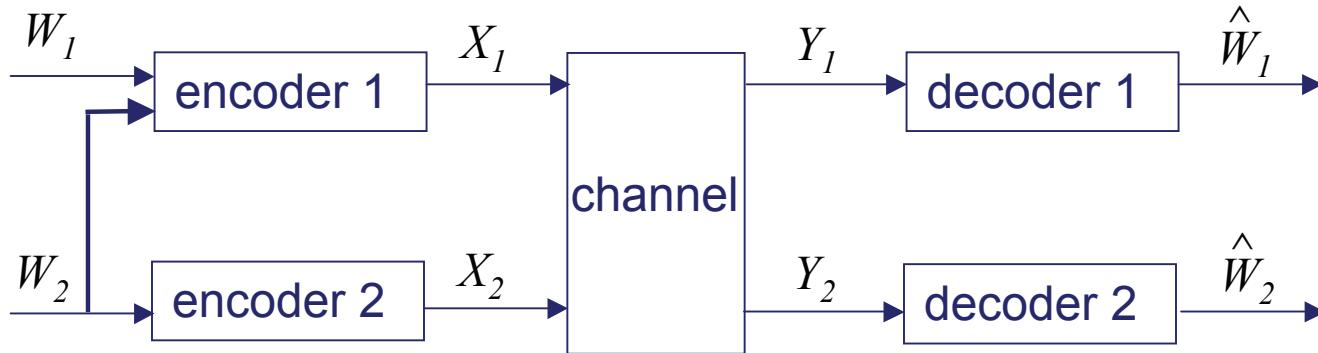
- λ One of the transmitters is a cognitive radio
- λ It “overhears” transmission of primary user
- λ It *partially* obtains primary user’s message

apple icon it can cooperate

- λ The setting becomes



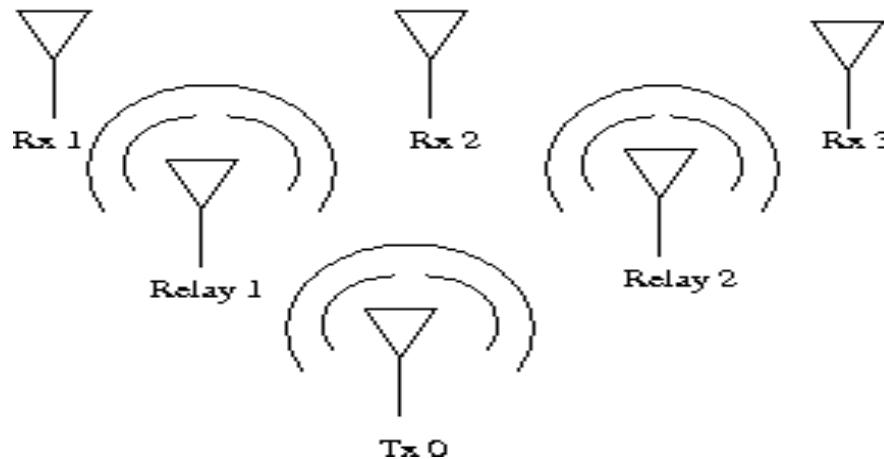
# Idealized model



- λ Secondary user learns the whole message  $W_2$
- λ Capacity region determined:
  - λ For the Gaussian channel in weak interference  
[Wu, Vishwanath & Arapostathis, 2006 ]
  - λ For strong interference [Maric, Yates & Kramer, 2005]
- λ We have upper and lower bounds on capacity for this model.

# Crosslayer Cooperative Broadcasting via Dualized Erasure Correction Codes

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- Major issue in wireless networks: interference.
- High level: noise well understood
- Many network designs suppress interference by orthogonalization and successive cancellation.
- If cooperation is enabled, interference can provide benefits.

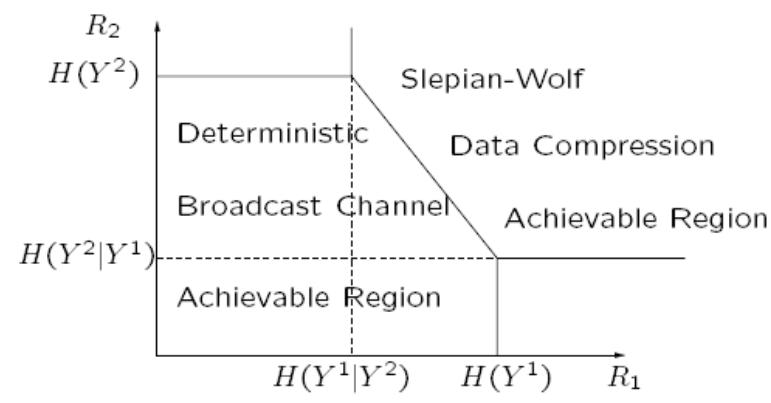
# Interference has value

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- λ Allow noise to be managed by physical layer, interference info exposed to app layer for handling.
  - λ *Interference carries information*

(Rel1,Rel2)	Rx1	Rx2	Rx3	
(-1,-1)	-1	-1	1	• Rx1 only hears Rel1, Rx3 only hears Rel2, Rx2 hears both
(-1,1)	-1	*	1	<i>Physical scenarios:</i>
(1,-1)	1	*	-1	• BPSK with additive combining
(1,1)	1	1	-1	• Non-coherent modulation with FSK

# Deterministic BC Channel



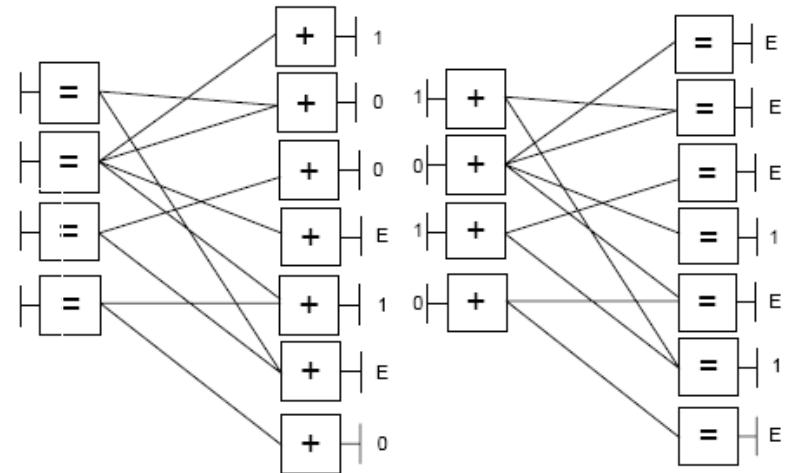
(Rel1, Rel2)	Rx1	Rx2	Rx3
(-1, -1)	-1	-1	1
(-1, 1)	-1	*	1
(1, -1)	1	*	-1
(1, 1)	1	1	-1

- λ Although DBC dual to Slepian Wolf, developing practical codes is non-trivial.
- λ Encoding is hard, decoding easy
- λ On opposite side of entropy boundary; so using off-the-shelf LDPCs yields exponentially many codewords consistent with what observed.

# Rate Splitting Approach

*A low-complexity  
capacity-achieving strategy*

- λ Coding at vertices
- λ First stage of pipeline, can use Cover’s “enumerative source coding” technique
- λ Later stages: use Luby’s generator-form LT codes
- λ Dualize the algorithm + dualize the linear code
- λ In 1-to-1 correspondence with Digital Fountain coding on the binary erasure channel



# Cooperation and Coding Posters

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- λ "Capacity Gain from Transmitter and Receiver Cooperation," C. Ng. and A. Goldsmith
- λ "Fundamental Limits of Networks with Cognitive Users," I. Maric and A. Goldsmith
- λ "Capacity Definitions of General Channels with Receiver Side Information," M. Effros, A. Goldsmith, and Y. Liang
- λ "Crosslayer Cooperative Broadcast Communication via Dualized Erasure Correction Codes", T. Coleman