

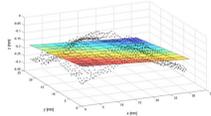
# High Dimensional Linear Regression with Binary Coefficients.

## Mean Squared Error and Phase Transitions

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The problem of recovering the sparsity pattern of an unknown vector  $\beta^*$  based on noisy observations arises in a broad variety of contexts including subset selection in regression, structure estimation in graphical models and signal denoising.



### Our Model

**Setup:** Let  $\beta^* \in \{0, 1\}^p$  be a **binary  $k$ -sparse** vector. For

- $X \in \mathbb{R}^{n \times p}$  consisting of entries i.i.d.  $N(0, 1)$  **random variables**
- $W \in \mathbb{R}^n$  consisting of entries i.i.d.  $N(0, \sigma^2)$  **random variables** with  $\sigma^2 = o(k)$

we get  $n$  noisy linear samples of  $\beta^*$ ,  $Y \in \mathbb{R}^n$ , given by,

$$Y := X\beta^* + W.$$

**Goal:** Given  $(Y, X)$ , recover **w.h.p.**  $\beta^*$  with the minimum  $n$  possible.

### Why binary?

- Discrete structure  $\Rightarrow$  easier to analyze.
- Keeps the challenge of **support recovery** (highly nontrivial)
- Best known information theoretic lower bound is **much smaller** than the best known algorithmic upper bound.

### Literature Review

- Best known **positive** results (e.g. [Donoho '06],[Wainwright '09]) If

$$n > 2k \log p$$

many efficient algorithms (including LASSO) recover exactly  $\beta^*$  w.h.p.

- Best known **negative** result ([Wang et al '10]) If

$$n < n^* := \frac{2k}{\log\left(\frac{2k}{\sigma^2} + 1\right)} \log p,$$

then there is no recovery mechanism of  $\beta^*$  which succeeds w.h.p.



### Main Question

There is a **gap** in the literature when  $n^* < n < 2k \log p$ . Is there **enough information/ efficient algorithms** to recover  $\beta^*$  in this regime?

### Maximum Likelihood Estimator - All or Nothing result

It has a **simple-to-state form**: the MLE  $\hat{\beta}$  is the optimal solution of

$$(\Phi_2) \min_{\beta \in \{0,1\}^p, \sum_{i=1}^p \beta_i = k} \|Y - X\beta\|_2.$$

**Definition 1** For  $\beta \in \{0, 1\}^p$ ,  $k$ -sparse we define

$$\text{Overlap}(\beta) := |\text{Support}(\beta^*) \cap \text{Support}(\beta)|.$$

**Theorem 1 ("All or nothing")** Set  $n^* := \frac{2k}{\log\left(\frac{2k}{\sigma^2} + 1\right)} \log p$  and let  $\epsilon > 0$  be arbitrary.

- If  $n < (1 - \epsilon)n^*$ , then w.h.p.  $\frac{1}{k} \text{Overlap}(\hat{\beta}) \rightarrow 0$ , as  $n, p, k \rightarrow +\infty$ .
- If  $n > (1 + \epsilon)n^*$ , then w.h.p.  $\frac{1}{k} \text{Overlap}(\hat{\beta}) \rightarrow 1$ , as  $n, p, k \rightarrow +\infty$ .

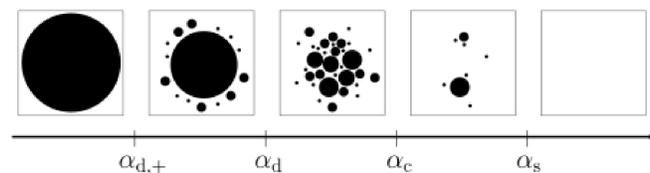
So, when  $n > n^*$  **information exists** and  $n^*$  is a **sharp phase transition point**.

### Algorithmic Difficulty

Why all known efficient algorithms seem to fail when  $n^* < n < 2k \log p$  and work only if  $n > 2k \log p$ ?

The picture from the analysis of random CSPs and spin glass theory suggests that a usual reason is an **"important change in the geometry of the space of solutions"** between the two regimes. [Achlioptas et al, 2008]

Such a geometrical property has been established for many problems such as *random  $k$ -SAT,  $k$ -coloring of a random graph, maximum independent set in a sparse random graph* and many others. (Figure below by [Krzakala et al '07])



### Overlap Gap Property in Linear Regression

We prove a geometrical property for the near-optimal feasible solutions of the problem  $(\Phi_2)$ . We call the property **Overlap Gap Property (OGP)** for high-dimensional linear regression. For  $r > 0$ , set

$$S_r := \{\beta \in \{0, 1\}^p : \|\beta\|_0 = k, n^{-\frac{1}{2}} \|Y - X\beta\|_2 < r\}.$$

**Definition 2 (The Overlap Gap Property)** Let  $r > 0$  and  $0 < \zeta_1 < \zeta_2 < 1$ . We say that the high-dimensional linear regression problem defined by  $(X, W, \beta^*)$  satisfies the Overlap Gap Property with parameters  $(r, \zeta_1, \zeta_2)$  if the following holds.

(a) For every  $\beta \in S_r$ ,

$$\frac{1}{k} \text{Overlap}(\beta) < \zeta_1 \text{ or } \frac{1}{k} \text{Overlap}(\beta) > \zeta_2.$$

(b) Both the sets

$$S_r \cap \{\beta : \frac{1}{k} \text{Overlap}(\beta) < \zeta_1\} \text{ and } S_r \cap \{\beta : \frac{1}{k} \text{Overlap}(\beta) > \zeta_2\}$$

are non-empty.

Intuitively, this means that the set of  $\beta^*$ s with closed to optimum objective value in  $(\Phi_2)$  **"shatters" in two components**, one with **low** overlap size with the ground truth  $\beta^*$  and one with **high** overlap size with  $\beta^*$ .

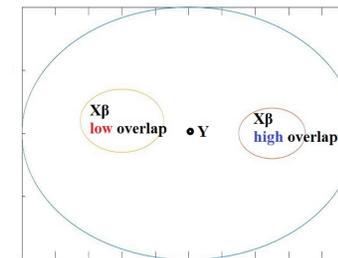


Figure 4: The OGP around  $Y$

**Theorem 2** Suppose the assumptions of Theorem 1 hold. There exists  $C > c > 0$  with the following properties.

- If  $n^* < n < ck \log p$  then there exists  $0 < \zeta_1 < \zeta_2 < 1$  and a sequence  $r_k > 0$  such that w.h.p. as  $k$  increases the high-dimensional problem defined by our model **satisfies** the Overlap Gap Property with parameters  $(r, \zeta_1, \zeta_2)$ .
- If  $n > Ck \log p$  then for any  $0 < \zeta_1 < \zeta_2 < 1$  and any sequence  $r_k > 0$  w.h.p. as  $k$  increases the high-dimensional problem defined by our model **does not satisfy** the Overlap Gap Property with parameters  $(r, \zeta_1, \zeta_2)$ .

**Corollary 1 (Informal)** If  $n < ck \log p$  then any "successful" local search algorithm needs in the worst case to increase the distance from  $Y$  in at least one step.

### Summary

- We **positively answer** the question of whether information for recovering  $\beta^*$  exists when  $n > n^*$ .
- We establish a certain Overlap Gap Property (OGP) in the space of binary  $k$ -sparse vectors when  $n < ck \log p$ . We conjecture that OGP is **the source of algorithmic hardness** of the problem when  $n^* < n < 2k \log p$ .

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