# High Dimensional Linear Regression using Lattice Basis Reduction Ilias Zadik, joint work with David Gamarnik 

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High Dimensional Linear Regression (HDLR)
Recovering unknown coefficients $\beta^{*}$ from few noisy observations Recovering large number of features arises in a broad variety of contexts including

- pricing of a product in the digital economy (econometrics) - GPS modeling and signal denoising (telecommunications)
- MRI analysis (compressive sensing)
- Generative Models and GANs (neural networks)


## The Model

Setup: Let $\beta^{*} \in \mathbb{R}^{p}$. For $X \in \mathbb{R}^{n \times p}$ and $W \in \mathbb{R}^{n}$ we get $n$ noisy inear samples of $\beta^{*}, Y \in \mathbb{R}^{n}$, given by, $Y:=X \beta^{*}+W$. Goal: Given data $(Y, X)$ with $Y:=X \beta^{*}+W$ recover $\beta^{*}$ with $n \ll p$ and $p \rightarrow+\infty$.
Regularity Assumptions and a Challenge
To achieve $n \ll p$ we need structural assumption on $\beta$.

- Sparsity! $k \leq p$ non zero coordinates.

Vast literature.
For $X$ with iid $N(0,1)$ entries and $W$ with iid $N\left(0, \sigma^{2}\right)$ entries $\left(\sigma^{2} \ll k\right)$ we need $k \log \left(\frac{p}{k}\right)$ samples (Compressed Sensing) Issue: $k \log \left(\frac{p}{k}\right)$ can still be too large for applications.

- Other assumptions:

Block-sparsity [Baron et al',05], Tree-Sparsity [He et al '09] Ouput of a Generative Model [Bora et al '17]
Similar issue: can achieve some $n<p$ but not always small.

## This Work

New efficient algorithm for recovering $\beta^{*}$ from $(Y, X)$ under a new regularity assumption
(Q-rationality assumption)
based on a connection with lattice-based algorithms

Guarantees: works for any $n$ (even $n=1$ ) given sufficiently small noise!

## The $Q$-Rationality Assumption

Every entry of $\beta^{*}$ is a rational number with fixed denominator $Q$. Alternatively: For $Q=2^{M}, \log Q=M$ bits after zero position pe entry.

Why $Q$-rational?

- Large $Q$ : Large but finite domain for the coefficients.
- Small Q: Standard in wireless communication.

Example: Linear models for GPS ([Boyd, Hassibi '98]) Physics laws imply integer coordinates $[Q=1]$.

## (2) ${ }^{2}$

Under $Q$-rationality, One Sample Suffices Lemma 1 Assume $X$ with iid $N(0,1)$ entries and $W$ with iid $N\left(0, \sigma^{2}\right)$ entries. Given one sample $n=1$ and small $\sigma$ we can recover exactly the $Q$-rational $\beta^{*}$
ntuition for $\sigma=0$ : Each row of $X, X_{1}$, has iid $N(0,1)$ entries and therefore linearly independent entries over rationals.
Hence, from ( $Y_{1}=\left\langle X_{1}, \beta^{*}\right\rangle, X_{1}$ ) we can recover $\beta^{*}$.

## Previous Computational Results

## Sample size needs to grow

- Convex Relaxations For $\beta^{*} \in\{-1,1\}^{p}(Q=1), \sigma=0$ consider $\min \|\beta\|_{\infty}$, s.t. $Y=X \beta$.
Works if and only if $n>p / 2$, i.e. needs linear sample ([Chandrasekaran et al '10], [Amelunxen et al '13])
Statistical-Physics-based algorithm (AMP) [Donoho et al' '11] for some $n=o(p)$ and any $Q$ but
we only know $n=o(p)$ (could be any sublinear quantity) and needs delicate choice of $X$ (not iid!)


## Main Results

Theorem 1 (Efficient Recovery with $n=1$ ) Let $n$ samples, $n \ll$ $p$, and $0 \leq \sigma \leq \exp \left(-\frac{p \max \{p, \log Q\}}{n}\right)$.
There exists a polynomial in $)$, $Q$, input $(Y, X)$ ouputs exactly $\beta^{*}$ w.h.p. as $p \rightarrow+\infty$

- The theorem works with any $X$ with iid well-behaved contin uous entries (or uniform iid integer in a large domain) and any $W$ with $\|W\|_{\infty} \leq \sigma$ !
Theorem 2 Let $n$ samples, $n \ll p$, and $\sigma>\exp \left(-\frac{p \log Q}{n}\right)$
Then if $X$ has iid $N(0,1)$ entries and $W$ iid $N\left(0, \sigma^{2}\right)$ entries, its impossible to w.h.p. recover correctly any $Q$-rational $\beta^{*}$ with any algorithm with only access to $(Y, X)$.
- If $\log Q>p$ : our algorithm has optimal noise-tolerance!


## Shortest Vector Problem (SVP)

For a lattice $\mathcal{L}$ (integer linear combinations of some vectors
$b_{1}, \ldots, b_{m} \in \mathbb{Z}^{p}$ ) the goal is to solve: $\min _{r \in \mathcal{L} \backslash\{ \}}\|x\|_{2}$


Well-studied in Integer Programming and Cryptography

The LLL Algorithm for SVP
Lattice Basis Reduction!
SVP is NP-Hard but Lenstra-Lenstra-Lovasz (LLL), algorithm efficiently approximates it; finds $\hat{x} \in \mathcal{L} \backslash\{0\}$ with

$$
\|\hat{x}\|_{2} \leq 2^{\frac{p}{2}} \min _{x \in \mathcal{L} \backslash\{0\}}\|x\|_{2} .
$$

$$
\text { Time poly in } p, \log \max _{i}\left\{\left\|b_{i}\right\|_{\infty}\right\} \text {. }
$$

## Using LLL for HDLR (General Scheme)

Step 1: Create a lattice $\mathcal{L}=\mathcal{L}(Y, X)$ such tha Step 2. Us L liple of $\beta^{*}$.
Step 3: Recover $\beta^{*}$ from a multiple (needs special structure!)
Note: Step 1 is Inspired by the use of LLL in cryptography
([Lagarias, Odlyzko '83], [Frieze '86])

## The Algorithm: Special Case

- $n=1, \sigma=0, \beta^{*}$ binary, $Y_{1}=\left\langle X_{1}, \beta^{*}\right\rangle^{\prime}$
- $X_{1} \in \mathbb{Z}^{p}$ with iid uniform in $\left[2^{N}\right]$ entries for large $N$

Step 1: For $M$ sufficiently large enough set $\mathcal{L}_{M}\left(Y_{1}, X_{1}\right)$ produced by the columns of

$$
A_{M}:=\left[\begin{array}{cc}
M X_{1} & -M Y_{1} \\
I_{p \times p} & 0
\end{array}\right]
$$

Lemma: Each $z \in \mathcal{L}_{M},\|z\|_{2}<M$ is a multiple of $\left[\begin{array}{c}0 \\ \beta^{*}\end{array}\right]$, w.h.p. Intuition
$z=A_{M}\left[\begin{array}{l}\beta \\ \lambda\end{array}\right]=\left[\begin{array}{c}M\left\langle X_{1}, \beta\right\rangle-M \lambda Y_{1} \\ \beta\end{array}\right]=\left[\begin{array}{c}M\left\langle X_{1}, \beta-\lambda \beta^{*}\right\rangle \\ \beta\end{array}\right]$,
Either $\left|z_{1}\right| \geq M \Rightarrow\|z\|_{2} \geq M$ or
$z_{1}=0 \Rightarrow\left\langle X_{1}, \beta-\lambda \beta^{*}\right\rangle=0$, low probability with $\beta \neq \lambda \beta^{*}$ ! Step 2: Choose $M$ appropriately so that LLL outputs a multiple of $\beta^{*}$.

- Choose $M=\left\lceil 2^{\frac{p}{2}} \sqrt{p}\right\rceil+1$.
- We know $A_{M}\left[\begin{array}{c}\beta^{*} \\ 1\end{array}\right]=\left[\begin{array}{c}0 \\ \beta^{*}\end{array}\right] \in \mathcal{L}$.
- LLL outputs $\hat{x}$ with norm at most $2^{\frac{p}{2}}\left\|\left[\begin{array}{c}0 \\ \beta^{*}\end{array}\right]\right\|_{2} \leq 2^{\frac{p}{2}} \sqrt{p}<M$.
- Using the lemma we are done!

Step 3: Rescale to get $\beta^{*}$.

## Special Case $\rightarrow$ General Case

1) One sample $n=1 \rightarrow$ many samples $n>1$. Way: Redesign the Lattice)
(2) Noiseless $\sigma=0 \rightarrow$ noisy $\sigma>0$
(Way. Redesign the Latice)
2) Integer $Y, X \rightarrow$ real $Y, X$
(Way: Truncate first bits and Rescale the data $(Y, X)$ )
(4) Binary coefficients $\beta^{*} \rightarrow Q$-rational $\beta^{*}$.
(Way: Translate and Rescale the samples $Y$ )
(Preliminary) Experiments
(joint work with Patricio Foncea and Andrew Zheng) Integer Data
Assume $X$ iid uniform in $\left[2^{N}\right], \beta^{*}$ iid uniform in $[100]$ and no noise iccess is exact recovery.

$$
\text { Plot: Avg Success/ Running Time against input size } N \text {. }
$$




Figure 5: (20 instances per dot) $p=30, n=1,10,30, \alpha \sim 1 / N$

Assume $X$ iid $U(0,1), W$ Reai-valued $U(-\sigma, \sigma)$ and $\beta^{*}$ iid uniform in [100].
Plot: Avg Success against noise level $\sigma$ and truncation leve.


Figure 6: (20 instances per dot) $p=30$ and $n=10$.

## Conclusion

- High dimensional linear regression with rational co efficients can be efficiently solved with one sam ple $n=1$, under small noise!
- New algorithm for high dimensional linear regression using lattice based methods (LLL algorithm).
- The algorithm has guarantees for large $p$, but also works well for small $p$.


## Open Questions

- Can lattice-based methods also be used for nonlinear inference problems?
Example: Phase-Retrieval where $Y_{i}=\left|\left\langle X_{i}, \beta^{*}\right\rangle\right|$
(many applications in Crystallography and MRI)
- Can we tolerare higher noise levels for smaller $Q$ ?

