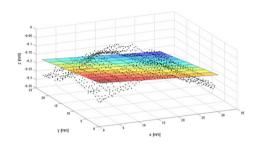


## High Dimensional Linear Regression (HDLR)

Recovering unknown coefficients  $\beta^*$  from few noisy observations and large number of features arises in a broad variety of contexts including

- pricing of a product in the digital economy (econometrics)
- GPS modeling and signal denoising (telecommunications)
- MRI analysis (compressive sensing)
- Generative Models and GANs (neural networks)



## The Model

**Setup:** Let  $\beta^* \in \mathbb{R}^p$ . For  $X \in \mathbb{R}^{n \times p}$  and  $W \in \mathbb{R}^n$  we get n noisy linear samples of  $\beta^*$ ,  $Y \in \mathbb{R}^n$ , given by,  $Y := X\beta^* + W$ . **Goal:** Given data (Y, X) with  $Y := X\beta^* + W$ recover  $\beta^*$  with  $n \ll p$  and  $p \to +\infty$ .

## **Regularity Assumptions and a Challenge**

To achieve  $n \ll p$  we need structural assumption on  $\beta^*$ .

• Sparsity!  $k \le p$  non zero coordinates.

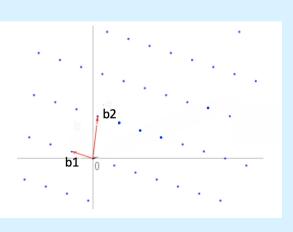
Vast literature. For X with iid N(0,1) entries and W with iid  $N(0,\sigma^2)$  entries  $(\sigma^2 \ll k)$  we need  $k \log \left(\frac{p}{k}\right)$  samples (Compressed Sensing) **Issue:**  $k \log \left(\frac{p}{k}\right)$  can still be **too large** for applications.

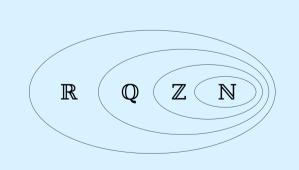
• Other assumptions: Block-sparsity [Baron et al',05], Tree-Sparsity [He et al '09] Ouput of a Generative Model [Bora et al '17] Similar issue: can achieve some n < p but not always small.

## This Work

**New efficient algorithm** for recovering  $\beta^*$  from (Y, X)under a new regularity assumption (**Q**-rationality assumption)

based on a connection with lattice-based algorithms.





*Guarantees:* works for any n (even n = 1) given sufficiently small noise!

## **The** *Q***-Rationality Assumption**

Every entry of  $\beta^*$  is a **rational number** with fixed denominator Q. Alternatively: For  $Q = 2^M$ ,  $\log Q = M$  bits after zero position per entry.

# High Dimensional Linear Regression using Lattice Basis Reduction

# Ilias Zadik, joint work with David Gamarnik

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## Why Q-rational?

- Large Q: Large but finite domain for the coefficients.
- Small Q: Standard in wireless communication. *Example:* Linear models for GPS ([Boyd, Hassibi '98]), **Physics laws** imply integer coordinates [Q = 1].



## **Under** *Q*-rationality, One Sample Suffices

**Lemma 1** Assume X with iid N(0,1) entries and W with iid  $N(0,\sigma^2)$  entries. Given one sample n=1 and small  $\sigma$  we can recover **exactly** the Q-rational  $\beta^*$ .

**Intuition for**  $\sigma = 0$ : Each row of X, X<sub>1</sub>, has iid N(0, 1) entries and therefore linearly independent entries over rationals. Hence, from  $(Y_1 = \langle X_1, \beta^* \rangle, X_1)$  we can recover  $\beta^*$ .

## **Previous Computational Results**

Sample size needs to grow!

• Convex Relaxations For  $\beta^* \in \{-1, 1\}^p$  (Q = 1),  $\sigma = 0$  consider  $\min \|\beta\|_{\infty}, \text{ s.t. } Y = X\beta.$ Works if and only if n > p/2, i.e. needs **linear samples** 

([Chandrasekaran et al '10], [Amelunxen et al '13])

- Statistical-Physics-based algorithm (AMP) [Donoho et al' '11] works for some n = o(p) and any Q but
- we only know n = o(p) (could be any sublinear quantity) and – needs delicate choice of X (not iid!)

## Main Results

Theorem 1 (Efficient Recovery with n = 1) Let n samples,  $n \ll 1$ p, and  $0 \leq \sigma \leq \exp\left(-\frac{p\max\{p,\log Q\}}{n}\right)$ .

There exists a **polynomial-in-**  $n, p, \log Q$  time algorithm which with input (Y, X) ouputs exactly  $\beta^*$  w.h.p. as  $p \to +\infty$ .

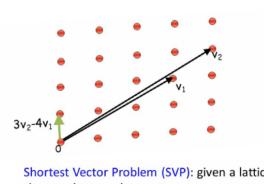
• The theorem works with any X with iid well-behaved contin**uous** entries (or uniform iid integer in a large domain) and any W with  $||W||_{\infty} \leq \sigma!$ 

**Theorem 2** Let n samples,  $n \ll p$ , and  $\sigma > \exp\left(-\frac{p \log Q}{n}\right)$ . Then if X has iid N(0,1) entries and W iid  $N(0,\sigma^2)$  entries, its impossible to w.h.p. recover correctly any Q-rational  $\beta^*$  with **any** algorithm with only access to (Y, X).

• If  $\log Q > p$ : our algorithm has **optimal noise-tolerance**!

# Shortest Vector Problem (SVP)

For a lattice  $\mathcal{L}$  (integer linear combinations of some vectors  $b_1, \ldots, b_m \in \mathbb{Z}^p$ ) the goal is to solve:  $\min_{x \in \mathcal{L} \setminus \{0\}} \|x\|_2$ 



Well-studied in Integer Programming and Cryptography.



# Using LLL for HDLR (General Scheme)

Step 1: Create a lattice  $\mathcal{L} = \mathcal{L}(Y, X)$  such that "approximately" shortest vectors of  $\mathcal{L} \leftrightarrow$  multiples of  $\beta^*$ Step 2: Use LLL and recover a multiple of  $\beta^*$ . Step 3: Recover  $\beta^*$  from a multiple (needs special structure!) **Note:** Step 1 is Inspired by the use of LLL in cryptography ([Lagarias, Odlyzko '83], [Frieze '86])

•  $X_1 \in \mathbb{Z}^p$  with iid **uniform in**  $[2^N]$  entries for large N Step 1: For M sufficiently large enough set  $\mathcal{L}_M(Y_1, X_1)$  produced by the columns of

# Intuition:

z =

- Choc
- We k

# The LLL Algorithm for SVP

Lattice Basis Reduction! SVP is NP-Hard but Lenstra-Lenstra-Lovasz (LLL), algorithm efficiently *approximates* it; finds  $\hat{x} \in \mathcal{L} \setminus \{0\}$  with

 $\|\hat{x}\|_2 \le 2^{\frac{p}{2}} \min_{x \in \mathcal{L} \setminus \{0\}} \|x\|_2.$ 

Time poly in p,  $\log \max_i \{ \|b_i\|_{\infty} \}$ .

## The Algorithm: Special Case

• n = 1,  $\sigma = 0$ ,  $\beta^*$  binary,  $Y_1 = \langle X_1, \beta^* \rangle$ .

$$A_M := \begin{bmatrix} MX_1 & -MY_1 \\ I_{p \times p} & 0 \end{bmatrix}$$

**Lemma:** Each  $z \in \mathcal{L}_M$ ,  $||z||_2 < M$  is a multiple of  $\begin{bmatrix} 0 \\ \beta^* \end{bmatrix}$ , w.h.p.

$$A_M \begin{bmatrix} \beta \\ \lambda \end{bmatrix} = \begin{bmatrix} M \langle X_1, \beta \rangle - M \lambda Y_1 \\ \beta \end{bmatrix} = \begin{bmatrix} M \langle X_1, \beta - \lambda \beta^* \rangle \\ \beta \end{bmatrix}$$

Either  $|z_1| \ge M \Rightarrow ||z||_2 \ge M$  or  $z_1 = 0 \Rightarrow \langle X_1, \beta - \lambda \beta^* \rangle = 0$ , low probability with  $\beta \neq \lambda \beta^*$ !

Step 2: Choose M appropriately so that LLL outputs a multiple of

ose 
$$M = \lceil 2^{\frac{p}{2}} \sqrt{p} \rceil + 1.$$
  
know  $A_M \begin{bmatrix} \beta^* \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ \beta^* \end{bmatrix} \in \mathcal{L}.$ 

• LLL outputs  $\hat{x}$  with norm at most  $2^{\frac{p}{2}} \| \begin{bmatrix} 0 \\ \beta^* \end{bmatrix} \|_2 \le 2^{\frac{p}{2}} \sqrt{p} < M$ .

• Using the lemma we are done!

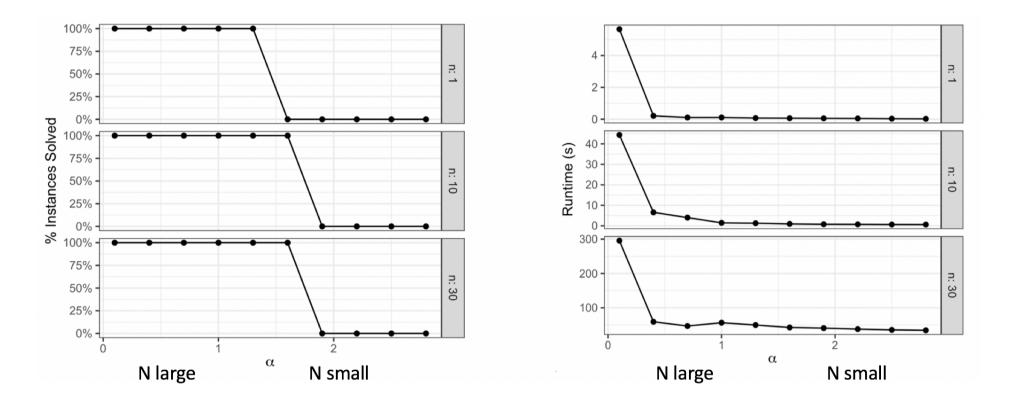
Step 3: Rescale to get  $\beta^*$ .

## Special Case $\rightarrow$ General Case

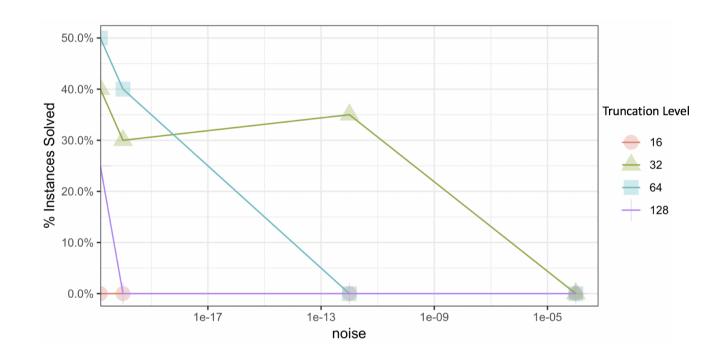
```
(1) One sample n = 1 \rightarrow many samples n > 1.
    (Way: Redesign the Lattice)
(2) Noiseless \sigma = 0 \rightarrow \text{noisy } \sigma > 0.
    (Way: Redesign the Lattice)
(3) Integer Y, X \rightarrow \text{real } Y, X.
    (Way: Truncate first bits and Rescale the data (Y, X))
(4) Binary coefficients \beta^* \to Q-rational \beta^*.
    (Way: Translate and Rescale the samples Y)
```

(joint work with Patricio Foncea and Andrew Zheng) Integer Data

Assume X iid uniform in  $[2^N]$ ,  $\beta^*$  iid uniform in [100] and no noise. Success is exact recovery.



Assume X iid U(0,1), W iid  $U(-\sigma,\sigma)$  and  $\beta^*$  iid uniform in [100]. Success is exact recovery. *Plot:* Avg Success against noise level  $\sigma$  and truncation level.



- works well for small p.

## (Preliminary) Experiments

*Plot:* Avg Success/ Running Time against input size N.

**Figure 5:** (20 instances per dot) p = 30, n = 1, 10, 30,  $\alpha \sim 1/N$ .

## **Real-valued Data**

Figure 6: (20 instances per dot) p = 30 and n = 10.

## Conclusion

 High dimensional linear regression with rational coefficients can be efficiently solved with one sam**ple** n = 1, under small noise!

• New algorithm for *high dimensional linear regression* using lattice based methods (LLL algorithm).

• The algorithm has guarantees for large p, but also

## **Open Questions**

 Can lattice-based methods also be used for non**linear inference** problems?

Example: *Phase-Retrieval* where  $Y_i = |\langle X_i, \beta^* \rangle|$ (many applications in Crystallography and MRI). • Can we tolerare **higher noise levels** for smaller Q?