

Encoder Side Information *Is* Useful In Source Coding

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Abstract — We introduce the idea of distortion side information, which does not directly depend on the source but instead affects the distortion measure. Such side information is not only useful at the encoder, but under many conditions of interest, knowing it at the encoder alone is sufficient and knowing it at the decoder alone is useless.

I. INTRODUCTION

We consider side information that is independent of the source but affects the distortion measure. Such distortion side information is a natural complement to the signal side information studied by Wyner and Ziv that depends on the source but does not affect the distortion measure. When both types of side information are considered, knowing the former only at the encoder and the latter only at the decoder is often asymptotically as good as complete knowledge of all side information. Potential examples of distortion side information include reliabilities for sensor readings and perceptual effects such as distortion masking, and sensitivity to context [1, 2].

II. PROBLEM MODEL

We consider a source $\mathbf{x} = (x_1, x_2, \dots, x_n)$, which is independent and identically distributed (i.i.d.) according to the distribution $p_x(x)$. The signal side information \mathbf{w} is i.i.d. generated according to $p_{w|x}(w|x)$ and the distortion side information \mathbf{q} is i.i.d., independent of (\mathbf{x}, \mathbf{w}) , and generated according to $p_q(q)$. Distortion between a source \mathbf{x} and reconstruction $\hat{\mathbf{x}}$ is measured via $(1/n) \sum_{i=1}^n d(x_i, \hat{x}_i; q_i)$. Encoders, decoders, and rate-distortion functions are defined in the usual way.

As an example, consider a sensor application where each observation is observed at a different signal-to-noise ratio (SNR). We could let q_i represent the SNR for observation x_i and measure distortion via $d(x, \hat{x}; q) = q \cdot (x - \hat{x})^2$. Intuitively, the more reliable observations where q_i is large, should be communicated to the decoder more accurately than noisier observations where q_i is small. If \mathbf{q} is known by both encoder and decoder, a water-pouring distortion allocation achieves this goal. In the next section we analyze performance for distributed knowledge of \mathbf{q} .

III. MAIN RESULTS

We denote the various possible rate-distortion functions by describing where the side-information is available and derive various relationships between them. For example, $R_{Q-NONE-W-NONE}(D)$ denotes the rate-distortion function without side information and $R_{Q-NONE-W-DEC}(D)$ denotes the rate-distortion function where \mathbf{w} is available at the decoder studied by Wyner and Ziv [3]. Similarly, when all information is available at both encoder and decoder

$R_{Q-BOTH-W-BOTH}(D)$ describes Csiszár and Körner’s [4] generalization of Gray’s conditional rate-distortion function $R_{Q-NONE-W-BOTH}(D)$ [5] to the case where the side information can affect the distortion measure.

We need various technical conditions corresponding to a “smooth” source and distortion measure; see [1, 6] for details. Our first result (generalizing [7]) is that having \mathbf{q} only at the encoder and \mathbf{w} only at the decoder is asymptotically optimal in high-resolution; see also [1, 2] for more general results.

Theorem 1 *Let \mathbf{q} be statistically independent of (\mathbf{x}, \mathbf{w}) . For any source and scaled difference distortion measure $d(x, \hat{x}; q) = d_0(q) \cdot d_1(x - \hat{x})$ satisfying the conditions in [1]*

$$\lim_{D \rightarrow 0} R_{Q-ENC-W-DEC}(D) - R_{Q-BOTH-W-BOTH}(D) = 0. \quad (1)$$

Our next result shows when side information can be useless.

Theorem 2 *Let \mathbf{q} be statistically independent of (\mathbf{x}, \mathbf{w}) and consider a difference distortion measure of the form $d(x - \hat{x}; q)$. Then \mathbf{w} (respectively \mathbf{q}) provides no benefit when known only at the encoder (resp. decoder), i.e.,*

$$R_{Q-DEC-W-ENC}(D) = R_{Q-NONE-W-NONE}(D). \quad (2)$$

Furthermore, in [1] we generalize these results to show that the value of a given type of side information at the encoder/decoder is essentially independent of where the other type of side information is known. For example, under the conditions of the above theorems we derive results like $R_{Q-ENC-W-ENC}(D) = R_{Q-ENC-W-NONE}(D)$, $R_{Q-DEC-W-DEC}(D) = R_{Q-NONE-W-DEC}(D)$, and

$$\lim_{D \rightarrow 0} R_{Q-ENC-W-ENC}(D) - R_{Q-BOTH-W-ENC}(D) = 0.$$

Thus knowing distortion side information, \mathbf{q} , at the encoder and signal side information, \mathbf{w} , at the decoder is asymptotically optimal in a wide variety of scenarios.

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