Abstract and Semicontractive Dynamic Programming

Dimitri P. Bertsekas

Department of Electrical Engineering and Computer Science Massachusetts Institute of Technology

Lecture 1 of 5

July 2016

System: $x_{k+1} = f(x_k, u_k, w_k)$

- *x_k*: State at time *k*, from some space *X*
- *u_k*: Control at time *k*, from some space *U*
- w_k : Random "disturbance" at time k, from a countable space W, with $p(w_k | x_k, u_k)$ given

Policies: $\pi = \{\mu_0, \mu_1, ...\}$

- Each μ_k maps states x_k to controls $u_k = \mu_k(x_k) \in U(x_k)$ (a constraint set)
- Cost of π starting at x_0 , with discount factor $\alpha \in (0, 1]$:

 $J_{\pi}(x_0) = \limsup_{N \to \infty} E\left\{\sum_{k=0}^{N} \alpha^k g(x_k, \mu_k(x_k), w_k)\right\}$

- Optimal cost starting at x_0 : $J^*(x_0) = \inf_{\pi} J_{\pi}(x_0)$
- Optimal policy π^* : Satisfies $J_{\pi^*}(x) = J^*(x)$ for all $x \in X$
- Stationary policies, those of the form {μ, μ,...}, play a special role (typically, there are stationary optimal policies that are optimal)

Bellman's Equation

The cost of a stationary policy μ starting from state x, denoted J_μ(x), typically satisfies

 $J_{\mu}(x) = E\{g(x,\mu(x),w) + \alpha J_{\mu}(f(x,\mu(x),w))\}, \quad \forall x \in X$

This is called Bellman's equation for policy μ

• The optimal cost starting from state x, denoted $J^*(x)$, typically satisfies

 $J^*(x) = \inf_{u \in U(x)} E\{g(x, u, w) + \alpha J^*(f(x, u, w))\}, \quad \forall x \in X$

This is called Bellman's equation

- Both types of Bellman's equation are functional equations in J_{μ} or J^*
- They can be viewed abstractly as having the form

$$J_{\mu} = T_{\mu}J_{\mu}$$
 or $J^* = TJ^*$

- In a given DP problem it is significant when Bellman's equation has a unique solution. This is true if all T_{μ} are contraction mappings (with a common modulus). If this is true, the problem is called contractive and otherwise noncontractive
- Contractive problems are much nicer!

Two Main Classes of Total Cost SOC Problems

Contractive Problems:

- $\alpha < 1$ and bounded g
- Date to 50s (Bellman, Shapley)
- Nicest results; key fact is the contraction property of the mapping in Bellman's equation

Noncontractive Problems - Stochastic Shortest Path (SSP):

- Date to 60s (Eaton-Zadeh, Derman, Pallu de la Barriere)
- Also known as first passage or transient programming
- Aim is to reach a special termination state at min expected cost
- Under favorable assumptions, the results are almost as strong as for the discounted case (when the noncontractive policies cannot be optimal)
- In general, very complex behavior is possible

Some Additional Noncontractive Problems:

- Discounted problems with unbounded g
- Undiscounted problems with positive and negative cost ($g \le 0$ or $g \ge 0$)

Intermediate Problem Types: Between Contractive and Noncontractive

- Problems where some policies are "well-behaved" and some are not
- "Well-behaved" has a problem-dependent meaning. The most common example of "well-behaved" policy is one that is contractive
- Pathological behaviors are due to policies that are not "well-behaved"

Our Approach

- Select a class of well-behaved policies (we call them regular and define them in a precise way later)
- Define a restricted optimization problem over the regular policies only
- Show that the restricted problem has nice theoretical and algorithmic properties
- Relate the restricted problem to the original
- Under reasonable conditions, obtain strong theoretical and algorithmic results

References

Research Monograph

D. P. Bertsekas, Abstract Dynamic Programming, Athena Scientific, 2013; updates on-line.

Subsequent Papers

- D. P. Bertsekas, "Regular Policies in Abstract Dynamic Programming," Lab. for Information and Decision Systems Report LIDS-P-3173, MIT, May 2015.
- D. P. Bertsekas, "Affine Monotonic and Risk-Sensitive Models in Dynamic Programming", Lab. for Information and Decision Systems Report LIDS-3204, MIT, June 2016.
- D. P. Bertsekas, 'Robust Shortest Path Planning and Semicontractive Dynamic Programming", Naval Research Logistics J., to appear.
- D. P. Bertsekas, "Value and Policy Iteration in Optimal Control and Adaptive Dynamic Programming," IEEE Transactions on Neural Networks and Learning Systems, to appear.
- D. P. Bertsekas and H. Yu, "Stochastic Shortest Path Problems Under Weak Conditions," Lab. for Information and Decision Systems Report LIDS-P-2909, MIT, January 2016.

- Semicontractive Examples.
- Semicontractive Analysis for Stochastic Optimal Control.
- Extensions to Abstract DP Models.
- Applications to Stochastic Shortest Path and Other Problems.
- Algorithms.

Denote

$$H(x, u, J) = E\{g(x, u, w) + \alpha J(f(x, u, w))\}$$

• J* satisfies Bellman's equation

$$J^*(x) = \inf_{u \in U(x)} H(x, u, J^*), \qquad \forall \ x \in X$$

and if $\mu^*(x)$ attains the min for all x, μ^* is optimal

• The value iteration (VI) method converges: $J_k \rightarrow J^*$, where

$$J_{k+1}(x) = \inf_{u \in U(x)} H(x, u, J_k)$$

• The policy iteration (PI) method converges: $J_{\mu^k} \to J^*$, where $\{\mu^k\}$ is generated by

$$\begin{aligned} J_{\mu^k}(x) &= H\big(x, \mu^k(x), J_{\mu^k}\big), \quad \forall \ x \in X, \qquad \text{(policy evaluation)} \\ \mu^{k+1}(x) &\in \arg\min_{u \in U(x)} H(x, u, J_{\mu^k}), \quad \forall \ x \in X. \qquad \text{(policy improvement)} \end{aligned}$$

Four Pathological Examples: An Overview

 In all examples, we introduce a set of "well-behaved" or "regular" policies (in shortest path problems, regular policies will be the ones that reach the termination state in finite time).

Let

 $J^*(x)$: Optimal cost (over all policies) starting from x

 $\hat{J}(x)$: Optimal cost over the regular policies only, starting from x

The Four Examples

- A finite-state, finite-control deterministic shortest path problem. Here Bellman's equation may have multiple solutions (including *J*^{*} and *Ĵ*), and VI and PI may not converge to *J*^{*} or to *Ĵ*
- A finite-state, finite-control stochastic shortest path problem. Here *J*^{*} does not satisfy Bellman's equation, while \hat{J} does
- A finite-state, infinite-control stochastic shortest path problem. Here there is no optimal policy, and VI and PI exhibit some peculiarities
- A linear-quadratic optimal control problem. Here Bellman's equation has two solutions, *J*^{*} and *Ĵ*, and VI and PI typically converge to *Ĵ*

A Deterministic Shortest Path Problem



• VI for the irregular policy $\mu': J_{\mu', k+1}(1) = J_{\mu', k}(1)$ (fails)

Policy iteration (PI) starting from μ

If b < 0: Oscillates between μ and μ' . If b > 0: Converges to suboptimal μ

A Stochastic Shortest Path Problem (from Bertsekas and Yu, 2015)

A single policy μ . The only uncertainty is at the first stage starting at state 1.



The Bellman Eq. is violated at 1: $J_{\mu}(1) \neq \frac{1}{2}J_{\mu}(2) + \frac{1}{2}J_{\mu}(5)$

A peculiar phenomenon

Consider the deterministic optimal control problem where at state 1 we may choose either to go to 2 or to 5 at zero cost

- Then $J^*(x) = 1$ for all x, including $J^*(1) = 1$
- Belman's equation $J^*(1) = \min \{J^*(2), J^*(5)\}$ is satisfied
- Randomization lowers the optimal cost and invalidates Bellman's equation

The Blackmailer's Dilemma



- Every policy μ terminates with probability 1, and $J_{\mu}(1) = -\frac{1}{\mu(1)}$
- We have $J^*(1) = -\infty$ and there exists no optimal policy
- Bellman's equation is

$$J(1) = \min_{0 < u \le 1} \left\{ -u + (1 - u^2)J(1) \right\}$$

It is satisfied by $J^* = -\infty$ (also by $J = \infty$)

- VI converges to J* starting from any scalar J
- In PI we have $J_{\mu^k} \to J^*$, but $\mu^k(1) \to 0$ (which is not an admissible policy)
- A variation of the problem: Replacing the probability u^2 by u. Then $J^*(1) = -1$ is a solution of Bellman's Eq., but all $J \le -1$ are also solutions, and still there is no optimal policy

System: $x_{k+1} = \gamma x_k + u_k$, Cost per stage: $g(x, u) = u^2$

- Here $J^*(x) \equiv 0$ and the optimal policy is $\mu^*(x) \equiv 0$
- Bellman's equation is

$$J(x) = \min_{u \in \Re} \left\{ u^2 + J(\gamma x + u) \right\}, \qquad x \in \Re,$$

and is satisfied by J^* . Are there any other solutions?

Let $\gamma > 1$, so the system is unstable

- The optimal policy yields an unstable closed-loop system
- Bellman's equation has a second solution: $\hat{J}(x) = (\gamma^2 1)x^2$
- \hat{J} is the optimal cost function over the class of policies that stabilize the system (these are the "well-behaved" or "regular" policies)
- Both VI and PI typically converge to \hat{J} (not J^* !)

A Summary from the Examples

- Bellman's equation may have multiple solutions
- Often but not always, J* is a solution
- A restricted problem, involving "well-behaved" policies, is meaningful and plays an important role
- The appropriate set of "well-behaved" policies is problem-dependent (e.g., terminating in shortest path problems, or stabilizing in the linear quadratic case)
- The optimal cost function over all policies, J^* , may differ from \hat{J} , the optimal cost function over the "well-behaved" policies
- \hat{J} is the likely limit of the VI and the PI algorithms, starting from an appropriate set of initial conditions

In the next lecture, we will aim to:

- Explain this behavior through analysis
- Formalize the notion of "well-behaved" policy through a notion of regularity
- Introduce the kind of assumptions under which anomalous behavior can be avoided or mitigated
- Provide results of the type that are available for contractive problems