

OUTLINE

- Convexity issues in optimization
- Common geometrical framework for duality and minimax
- Unifying framework for existence of solutions and duality gap analysis
- Use of duality in algorithms

SOME HISTORY

- Late 19th-Early 20th Century:
 - Caratheodory, Minkowski, Steinitz, Farkas
- 40s-50s: The big turning point
 - Game Theory: von Neumann
 - Optimization-related convexity: Fenchel
 - Duality: Fenchel, Princeton group (Nash, Gale, Kuhn, Tucker)
- 60s-70s: Consolidation
 - Rockafellar
- 80s-90s: Extensions to nonconvex optimization and nonsmooth analysis
 - Clarke, Mordukovich, Rockafellar-Wets
- 2000- ... Rejuvenation: Many applications of large scale optimization using duality; resource allocation, combinatorial optimization, machine learning

WHY IS CONVEXITY IMPORTANT IN OPTIMIZATION I

- A convex function has no local minima that are not global
- A nonconvex function can be "convexified" while maintaining the optimality of its minima



WHY IS CONVEXITY IMPORTANT IN OPTIMIZATION II

- A polyhedral convex set is characterized by its extreme points and extreme directions
- Minima of linear functions over constraint sets can be found among extreme points





 A real-valued convex function is continuous and has nice differentiability properties



 Convex functions arise prominently in duality (a different but equivalent view of the same object)

DUAL DESCRIPTION OF CONVEX SETS



Primal description: Points

Dual description: Hyperplanes

CONJUGACY: DUAL DESCRIPTION OF CONVEX FUNCTIONS



Primal description: Values f(x) Dual description: Crossing points $f^{\star}(y)$









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Nice sets: Polyhedral, or convex and compact



I MIN COMMON/MAX CROSSING DUALITY



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CONVEX PROGRAMMING DUALITY

- Primal problem:
 min f(x) subject to x∈X and g_j(x)≤0, j=1,...,r
- Dual problem:

max q(μ) subject to $\mu \ge 0$

where the dual function is

$$q(\mu) = \inf_{x \in X} \left\{ f(x) + \mu' g(x) \right\}$$

• Consider common/max crossing framework: $M = epi(p), \qquad p(u) = inf_{x \in X, gj(x) \le uj} f(x)$

MINIMAX / ZERO SUM GAME THEORY ISSUES

 Given a function Φ(x,z), where x∈X and z∈Z, under what conditions do we have

 $\inf_x \sup_z \Phi(x,z) = \sup_z \inf_x \Phi(x,z)$

- Assume convexity/concavity, semicontinuity of Φ
- Min common/max crossing framework:
 M = epigraph of p

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p(u) = \inf_{x} \sup_{z} \{ \Phi(x,z) - u'z \}
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 $inf_x sup_z \Phi = Min \ common \ value$

 $sup_{z}inf_{x} \Phi = Max crossing value (can be shown)$

II UNIFICATION OF EXISTENCE AND NO DUALITY GAP ISSUES

INTERSECTIONS OF NESTED FAMILIES OF CLOSED SETS

- We will connect two basic problems in optimization
 - Attainment of a minimum of a function f over a set X
 - Existence of a duality gap
- The 1st question is a set intersection issue:

The set of minima is the intersection of the nonempty level sets $\{x \in X \mid f(x) \le \gamma\}$

 The 2nd question is also a set intersection issue (but not obvious). It is related to another fundamental question:

When is the function

 $f(x) = inf_z F(x,z)$

lower semicontinuous, assuming F(x,z) is convex and lower semicontinuous?



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PRESERVATION OF SEMICONTINUITY UNDER PARTIAL MINIMIZATION

• Key observation: For $f(x) = \inf_z F(x,z)$, we have $P(epi(F)) \subset epi(f) \subset cl(P(epi(F)))$

where $P(\cdot)$ is projection on the space of z. So if projection preserves closedness, f is l.s.c.



Given C, when is P(C) closed?

If y_k is a sequence in P(C) that converges to y, we must show that the intersection of the C_k is nonempty

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III PROBLEM STRUCTURES AND ALGORITHMS

SEPARABLE PROBLEMS & DECOMPOSITION

minimize
$$\sum_{i=1}^{n} f_i(x_i)$$

subject to $a'_j x \leq b_j, \quad j = 1, \dots, r,$

maximize
$$\sum_{i=1}^{n} q_i(\mu) - \sum_{j=1}^{r} \mu_j b_j \qquad q_i(\mu) = \inf_{x_i \in \Re} \left\{ f_i(x_i) + x_i \sum_{j=1}^{r} \mu_j a_{ji} \right\}$$
subject to $\mu \in \Re^r$,

Dual function calculated by decomposition

DISCRETE/INTEGER LP

minimize c'xsubject to $Ax \leq b$, $x_i = 0$ or 1, i = 1, ..., n

Typically there is duality gap

Solution of dual problem provides a lower bound Use in Lagrangian relaxation and branch-and-bound Dual problem is nondifferentiable/polyhedral

ADDITIVE COST PROBLEMS

$$f(x) = \sum_{i=1}^{m} f_i(x)$$

- Huge number m of terms
- This is a common structure:
 - Dual problems of separable problems
 - Expected values (e.g. in stochastic programming) have this structure
- Calculation of gradient or subgradient of the sum is very time consuming
- Need for an incremental algorithmic approach
 - Move x along the gradient/subgradient of a single component f_i

LARGE NUMBER OF CONSTRAINTS

minimize f(x)subject to $a'_j x \leq b_j, \quad j = 1, \dots, r,$

- Calls for approximation of the constraint set
- Outer approximation
- Inner approximation
- Inner and outer approx are dual to each other
- An alternative: Penalty approach converts to a minimization of a large sum

minimize
$$f(x) + c \sum_{j=1}^{r} P(a'_j x - b_j)$$

CLASSES OF ALGORITHMS

- Descent methods
 - Subgradient methods and incremental versions
 - Optimal algorithms (gradient methods with extrapolation) -Nesterov's methods
- Approximation methods
 - Cutting plane
 - Simplicial decomposition
 - Proximal and bundle method
 - Interior point methods
- All these methods rely on convexity concepts

CONCLUDING REMARKS

- Optimization has become a universal tool in applications
- Convexity is the "soul" of optimization
- Geometry is the "soul" of convexity
- Very few simple geometric ideas are sufficient to unify/clarify most of convex optimization
- Theoretical/algorithmic research on convex optimization is still very active

