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TO:

All Programmers

FROM:

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SUBJECT: An Abbreviated Description of the MAD Compiler Language

A. Introduction

The MAD compiler was designed and prepared originally for the IBM 704 computer by B. Arden, B. Galler and R. Graham at the University of Michigan. The compiler has subsequently been converted to the IBM 7090 by the Michigan group and with their cooperation been adapted by the MIT Computation Center staff to work within the FØRTRAN-FAP Monitor System used on many 7090 machines. The MAD language has many of the features found in FØRTRAN II, ALGEL 58, ALGEL 60 and the soom-to-be-introduced FØRTRAN IV. Since MAD has few restrictions and is highly flexible, it is well-suited as an initial computer language. Moreover, once a programmer has mastered one language, general experience has shown that subsequent languages are much easier to learn.

Only an essential subset of the MAD language is offered here but this subset is sufficient for complete, correct programs. The full language specification is given in the MAD Reference Manual (but with the final definition of the language, of course, implicitly contained in the MAD compiler program of approximately 12,000 instructions.)

B. Types of Constants

1. Integer:

of magnitude less than 227 and primarily used for counting and integer arithmetic. Integer constants are written in decimal with an optional plus sign if positive or an obligatory minus sign if negative but without a point. Examples: 4, 49, -100

2. Floating-point:

of magnitude zero or in the approximate range 10^{-38} to 10^{38} These numbers are primarily used in arithmetic calculations. Floating-point constants may be written with or without exponents. If written without an exponent, the constant contains from one to eight digits and a decimal point, which must be written, but which may appear anywhere in the number.

If the number is written with an exponent, it must contain from one to eight digits with or without a decimal point, followed by the letter E, followed by the exponent of the power of 10 that multiplies the number.

Examples: 3.0, 3. ,+.3 -123.45689 3.02+9, +97.58-30, 1008-5, -.983

C. Variables

of the corresponding mode. (Strictly speaking are 5 modes in MAD but the remaining three are 3e, below, for mode integer of floating-point and only takes on value: Each variable may be of one of the of variable names: l to 6 letters or digits of which the first is alphabetic. A variable is assumed to be of fil unnecessary here.) point mode unless declared otherwise. Variable names may consist ALPHA, MAKE, ANSWER, LANDES declaration statements.) (Strictly speaking there two modes? See section Tloating. Example:

D. Arrays

values between 1 and d where d is the fixed maximum subscript value to be deed for the 1th dimension. Subscript values of zero (or negative values) are not Subscripts s of the ith dimension are integer mode expressions (see Expressions below) which take on values between 1 and d where d is the fixed maximum subscript (1.0. vectors) where a subscript value of allowed here with the exception of arrays of one particular integer subscripts. The array name conventions used are the same as for variables. M-M-1), TABLE3 (5,7) has a fixed but arbitrary number of dimensions. in parentheses and separated by commis. For typographical reasons, subscripts are enclosed Arrays are sots of variables in which each variable were is allowed. mit of the array is distinguished by Examples: ALPHA(J), TABLE3(E+3, Rack array

E. External Functions

Function names follow the same rules as variable mames but must have a distinguishing terminal "." added. (In addition, all variables, arrays, and functions must have unique names.) The value of summetion can be of any mode. Examples: FIRSTF. ABCD.

. Expressions

exponentation, and absolute value respectively. The allowed integer floating-point constituents of arithmetic expressions are variables, scripted elements of arrays, constants, and functions. In general, all c stituents must be of the same mode but with the exception that if integer and floating-point modes are mixed, the expression is evaluated in the +,-,*,/,.P,, and .ABS. for addition, subtraction, multiplication, division, integer mode until conversion to the floating mode for example ((x,P,N)/T) + Qe,ABS,P-5, Arithmetic expressions may be formed using parentheses and the operations: The allowed integer or Secomes in general, all connecessary .

represents the expression

* 900

in the usual algebraic notation (where the wakue will be of the floating-point mode because of the floating-point five). All operations must be explicitly stated, e.g. (A+B)C is not correct, but (A+B)+C is correct.

The result of any integer mode calculation is an integer. When two integer mode numbers are divided only the integer portion of the result is kept as the quotient; the fractional portion of the result is discarded e_{off} , 7/3 = 2

Boolean expressions have the values true, or false. They may be formed using parentheses and the operators .NOT., $_{\circ}$ /R., .AND., .THEN., and .EQV., for "negation", "inclusive or", "and", "implication", and "equivalence", respectively. The allowed constitutents of boolean expressions are relations which have boolean values. Boolean-valued relations are $>_{\circ}$ / $>_{\circ}$ / $>_{\circ}$, and which for typographical reasons are written as .G., .GE., .E., .NE., .LE. and .L. respectively; these relations may occur between any two arithmatic expressions and mixed modes are allowed. For example

represents the boolean expression

$$((a \ge b) \land (o \ne d)) \lor (e_1 + 5) > (a/1000.))$$

G. Statements

A program consists of a collection of subprograms. Each subprogram written in the MAD language consists of a sequence of MAD statements, These statements, which compile into segments of machine language instructions, fall into 4 categories:

1. Substitution Statements are of the form

a m b

where b is an expression and a is a variable. When executed the new value of the variable a (in the appropriate mode) is computed by evaluating the expression b. Thus it is possible to transform from fixed to floating vice-versa. Note that if the variable a appears in the expression b, the old value of a is used in the computation. Examples:

Y= ALPHA/(B+C)-BETA+2.00S J= J+1

2. Control Statements

a. TRANSFER TO s is a statement which when executed causes the program to transfer to the statement labeled s in its left-margin (cols. 1-10). Statements are labeled in the same way as variables and each label must be distinct from all other statement labels, variables, arrays, or function names. Example:

TRANSFER TØ ALPHA

* 3

b. WHENEVER bas is a simple conditional statement where b is a boolean expression, and s is any executable statement except: END OF PROGRAM, another WHENEVER, THROUGH or a function ENTRY. If the boolean expression is true the statement s will be executed. Otherwise control will pass to the next statement following the conditional. The comma in this statement must be written. Examples:

Whenever x.L.100, transfer to stop Whenever y.G.22... \$6.5.E.X1. W=U

The sequence

C.

WHENEVER DOWNERS BY WHENEVER BY WH

end of Conditional

is a compound conditional statement where the b₄ are boolean expressions and the dots are any sequence of statements. The compound conditional is a pattern of statements which allows the conditional execution of just one of the program segments bracketed by the indicated statements. The segment executed is determined by the first boolean expression b₄ which is true. Any or all of the SH WHENEVER statements or the STHERWISE statement may be omitted. After the conditional execution of one program segment, program control automatically transfers to the statement logically following the required END SF CSNDITISNAL statement. This compound conditional is distinguished from the simple conditional by the absence of a comma following the boolean expressions.

Examples:

WHENEVER X.LE.O
Y=0.
Z=1.
ØR WHENEVER X.G.10.
Y=10.
Z=100.
OTHERW ISE
Y=X
Z=10. *X

END OF CONDITIONAL

Whenever W.L.10,
Transfer TØ LØØP2
ØR Whenever W.L.20,
Transfer TØ LØØP3
END ØF CØNDITIØNAL

- d. THROUGH s, FOR v = 2, b is an iteration statement which operates as follows:
 - l. The variable y is set equal to the expression el.
 - 2. The boolean expression b is tested. If b is true, program control passes to the statement after the statement labeled s.
 - 3. If b is false, the statements up to and including s (the scope) are executed.
 - 4. v is incremented by the value of the expression e2; return to step 2.

A THROUGH statement may appear within the scope of any other THROUGH, provided that the scope of the nested THROUGH lies entirely within the scope of the higher level THROUGH statement. Note that the statement labeled s may not be a declaration statement.

Examples

THROUGH LOOP, POR Mel, 2, M, G, 48

LAME

THROUGH LOOPI, FOR M-J.K-1, M.G. N+1 .OR.K.E.O

røddi .

THROUGH LOOP2, FOR J=1,1,J.G.M THROUGH LOOP2, FOR K=1,1,K.G.J

Lane 2

a calling sequence to the subprogram name. The listed by a calling sequence to the subprogram name. The listed by an expressions. For example EXECUTE HIM. (.5, X+Y, Z.P.5) creates a calling sequence to the subprogram HIM. with the values of the three expressions as arguments. The names of subprograms follow the same rules as for external functions. Example:

EXECUTE SØRT. (A)

f. CONTINUE is a dummy statement which when executed causes no action. The statement normally has a label and is used to indicate a joining point in a program, to which another statement may transfer. It is sometimes used as the last statement in the scope of a THRSUGH, to indicate to the reader that the end of the scope has been reached. Example:

LOGP3 CONTINUE

3. DECLARATION STATEMENTS

These statements only convey information to the MAD compiler and would not appear on a flow diagram of a program.

- a. The following monitor control card (fully defined later) and statement;
 - MAD

END OF PROGRAM

bracket a "Main" subprogram; upon subsequent loading of a program composed of several subprograms, the program is started at the beginning of the "Main" subprogram.

- h. Similarily the monitor control card and statements
 - B MAD

EXTERNAL FUNCTION (argument list)

ENTRY TO Da

FUNCTION RETURN V

END OF FUNCTION

bracket an external function subprogram which depending on how it is programed may be used in one of two different ways. In the first case the function named n. has a single value, \mathbf{v}_{θ} and is used in an expression. An example is the function SIN, in

Y=Y +3.5 *8IN. (X-.5)

in the second case, the single value v is meaningless (e.g. in a "function" to sort a table) and is omitted from the FUNCTION RETURN statement which is still required. This type of external function subprogram may only be used by an EXECUTE statement. In either case, the arguments listed (separated by commas) in an EXTERNAL FUNCTION statement must all be of the simple form of variable ames, unsubscripted array names or function names (without arguments). These names as they occur in the defining function subprogram act as dummies for compilation purposes only. In the use of the function, the arguments used may be expressions of the appropriate mode.

It also should be noted that, depending on the programming of the function definition, the argument list may not only contain input values but also may contain variable and array names which will contain output values of the external function subprogram after it is executed. Example:

EXECUTE SORT. (TABLES, N)

could sort the array named TABLES, which has N elements.

c. Every array used in a program must have adequate space reserved for it. For example, if the array A(J) may be used with J between 1 and 50, 50 locations must be set aside for the array A. For vectors (one-dimensional arrays) this setting aside may be done by the statement:

DIMENSION v(d)

where y is the name of the vector and d is a constant integer equal to the amount of storage needed. Actually, dil locations will be reserved, to allow use of the subscript zero. Several dimension statements may be combined, as in the example:

DIMENSIAN A(50), IARRAY(245), Z1(24)

Arrays of 2 or more dimensions require a slightly more complicated specification, including information about the structure of the array. With every multidimensional array is associated a vector (the "auxiliary vector") containing this information. This vector's name appears in the dimension specification of the multidimensional array. For example, if an array $A(J_0K_0L)$ may have all three subscripts between 1 and 10, 1000 spaces are needed. The dimension statement may appear as follows:

DIMENSION A(1000, AY) AV(3)

Note that the auxiliary vector must have space reserved for it, also, thus its name appears twice in the DIMENSION statement. The numbers stored in the auxiliary vector must be as follows:

AV(0) = number of dimensions in the array

AV(1) = 1 ocation of (1,1,...,1) element, usually 1.

AV(2) = maximum size of second subscript.

AV(3) = " " third "

AV(4) = " " fourth "

etc.

In the example AV(3) would be the last entry, since the array bad only 3 dimensions. The maximum size of the first subscript is not specified in the auxiliary vector. The auxiliary vector may be filled in by substitution statements at run time, or preset with a VKCTSR VALUES statement at compile time. (see below)

d. It is often desirable to preset an array with constant values at the compilation time of a program. This may be done by the statement:

VECTOR VALUES v(s)=cs, cs+1°cs+2°°°

where v is the array name and the c₁ are constants, all of the same mode. If s is zero, the subscript on v may be omitted. Examples:

VECTOR VALUES PLO(23) = 15.5, 21.7, 3E2 VECTOR VALUES AV = 3,1,10,10

It is not necessary to provide a DIMENSION statement for vectors preset with a VECTOR VALUES statement, if the element with the highest possible subscript has been preset.

A multidimensional array is stored in consecutive registers in memory. Thus it is possible to calculate an equivalent linear subscript using the following formula:

$$r = d_1 + (...,((s_1-1)-d_2 + (s_2-1))-d_3 +...)-d_n + (s_n-1)$$

where the di are the ith elements in the array's auxiliary vector (see above) and the si are the values of the subscripts.

An example of multidimensional to linear mapping is the presetting of entries in the middle of a multidimensional array. The multidimensional array a may be preset starting with the rth linear element by:

VECTER VALUES a(r) = C

where r has been calculated from the si⁰s and the auxiliary vector with the above formula. Example:

array $B(J_0K)_0$ 1 \le J \le 12, 1 \le K \le 6 to preset $B(4,3) \approx 15.5 \times 10^{12}$

Dimension B(72,BV)
VECTOR VALUES BV= 2,1,6
VECTOR VALUES B(21) = 15,5212

The dimension statement reserves 6K12=72 locations for B_0 and names the auxiliary vector BV_0 . The first VECTOR VALUES presets the auxiliary vector for B: 2 dimensions, B(1,1) is at 1, Max(K) = 6Since all entries of BV are preset, it is not dimensioned.

Since all entries of BV are preset, it is not dimensioned. The second VECTER VALUES presets B(4,3). The equivalent linear subscript from the formula is

$$r = d_1 + (s_1-1) \cdot d_{2+s_2}-1 = 1+(4-1) \cdot 6+3-1 = 21$$

e. All variables are assumed to be of the mode unless explicitly declared otherwise. Mode declarations hold for the entire subprogram and are made by the statement

INTEGER list

where <u>list</u> consists of one or more variable or array names, separated by commas. More than one mode declaration statement of a given type may be used. Example:

INTEGER A,B, IHAT

4. Input-Output Statements

Although it is possible to specify elaborate formats by which information can be brought in and out of the computer, for simplicity only three specific formats will be offered here. (The allowed formats of the Computation Center version of MAD are the same as those of FORTRAN II.) In all these statements, magnetic tape is used for input-output efficiency on the 7090 computer; an IBM 1401 computer is used to preprocess input tapes from cards and to post-process output tapes to the printed page.

a. To print out one or more lines of output (or read one or more input data cards) of integers which are right-justified in 4 adjacent fields 18 characters wide, one writes either of the statements:

PRINT PORMAT INT, list

or

READ FORMAT INT. list

where consecutive numbers on the in/out medium correspond to
the variables and array elements specified in the list. A list
may contain any number of variables, including none. In addition,
whenever integers are printed or read by a subprogram there must
be somewhere in the subprogram the format description statement:
Example:

VECTOR VALUES INT = \$(4118)\$

PRINT FORMAT INT.A,B(J).ARRAY(23),C,D,E

b. Similarly, to print (or read) one or more lines of real numbers which are right-justified in 4 adjacent fields 18 characters wide, the statements used are:

PRINT FORMAT REAL, list

or

READ PORMAT REAL, list

where list has the same meaning as for integers and somewhere in the subprogram there is also the format statement:

VECTOR VALUES REAL = \$(4E18.8)\$

c. To print a one line message of arbitrary text one must use the statement pair:

PRINT FORMAT 1

VECTOR VALUES $f = (ngh_1h_2h_3 ...h_n)$ \$

where n is an integer such that $1 \le n \le 72$ and the successive Hollerith characters h_1 contain the message. (The first character h_1 must be a blank.) The format vector name, \underline{f}_0 may be any unique name such as HOLL, MESS1, etc. Example:

PRINT FORMAT END1 VECTOR VALUES END1 = \$(20H END OF COMPUTATION.)\$

d. Preparation of Source Subprograms

The statements of a subprogram are called a "source subprogram" and are card-punched on IBM cards using the usual Hollerith codes to represent the letters of the alphabet, digits, and special characters. Statements are punched anywhere in the field of columns 12-72 (blanks are ignored with the obvious exception of the text included between \$ quote marks in VECTOR VALUES statements.) If a statement will not fit on one card it may be continued to another

card by punching consecutive digits from 1 to 9 in col. 11 of each continuing card. Statement labels are punched anywhere in the field of columns 1-10. If the letter "R" appears in column 11, that card will be ignored by the compiler, although it will be printed on the compilation listing (with the "R" deleted). This feature is used to make explanatory comments to a person reading the program. If a comment is to be continued "R" must appear in column 11 of each successive card.

e. Processing of a Program

Computation Center is usually done under the control of a system called the FØRTRAN-FAP-MAD Monitor System (FMS). Control of the monitor is by means of special control cards which have an asterisk punched in column 1 and key words in the cols. 7-72 of the card.

MAD source subprograms are compiled by the monitor using the MAD compiler program. The compiler examines each subprogram for various syntactical errors and may give one or more "error diagnostics."

When a subprogram does compile (it may still contain logical errors, however?) it is punched off-line as an "object subprogram" in the form of relocatable binary cards. Off-line printed output also is produced giving the machine language instructions of the compiled subprogram as well as other reference details useful in debugging.

To operate a program in the most elementary manner, the following procedure is used. A run deck is prepared with the following sections in sequence. (See attached example)

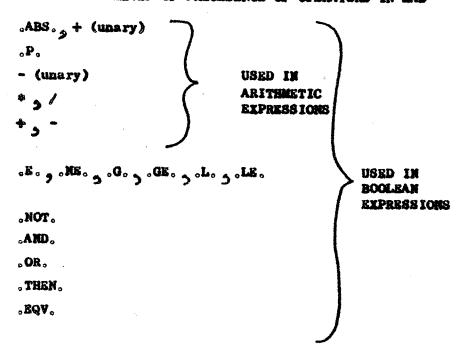
- l. Identification Card (ID) in a format given by the Center.
- An PMS control card * XEQ indicating that, if all source programs correctly translate, the program is to be executed.
- 3. The FMS control card * MAD indicating that the following program is to be compiled by the MAD translator.
- 4. "Main" subprogram and all other subprograms (in either source or object form) which are explicitly required by EXECUTE statements in these subprograms and which are not on the library tape.

 (The FMS system automatically supplies all needed Library subprograms including those implicitly required by inputoutput statements such as PRINT, etc., and by external functions such as SQRT, EXP., etc.)
- 5. The FMS control card * DATA
- 6. Data cards (if any) to be read from the input tape by the program.

This run deck is written on the off-line input tape (probably along with several other runs) and the FMS system is started. The monitor works on one run at a time, first compiling all the subprograms which are in source language. If there are no compiling diagnostics, the monitor brings in the "BSS loader" which then proceeds to read in each object subprogram, storing them in sequence starting from a lower address of core memory. During loading the loader maintains a storage map giving the location and symbolic name of each subprogram. When the FMS control card "*DATA" is encountered, a second phase of loader processing begins. The loader examines each subprogram to determine what other subprograms are needed by it (listed symbolically in a "transfer vector" at the beginning of each subprogram). If necessary, missing subprograms are obtained from the library tape. Using the completed storage map, the loader converts each symbolic name in a transfer vector to a corresponding transfer instruction to the memory location of the specified subprogram. Thus, the entry links between subprograms are established, and the program is started at the entry point of the "main subprogram,

The MIT version of MAD automatically inserts at an BND #F PROGRAM statement the necessary instruction to return to the Monitor system. When the program reaches this point, control returns to the monitor, and it then proceeds to the next job.

LEVEL OF PRECEDENCE OF OPERATORS IN MAD



```
*M2802=9000,FMS,DEBUG,1,1,0,0 JOHN DOE, 18 JUNE, 1963
           XEQ
           MAD
          R
                JOHN DOE
                INTERCHANGE SORT
          R
           VECTOR VALUES INT = $(4118)$
           DIMENSION KVL (100), KVA(100)
                READ IN INPUT
           READ FORMAT INT, N, KVL(1) ... KVL(N),
          1
                              KVA(1)...KVA(N)
                SORT ON KVL
           THROUGH LIST, FOR I=1,1,1.G.N
           KMIN = KVL (I)
           THROUGH SCAN, FOR J=I,I,J.G.N
           WHENEVER KVL(J).LE.KMIN
                KMIN = KVL(J)
                IMIN = J
             END OF CONDITIONAL
SCAN
           CONTINUE
           TEMP = KVL(IMIN)
           KVL(IMIN) = KVL(I)
           KVL(I) = TEMP
           TEMP = KVA(IMIN)
           KVA(IMIN) * KVA(I)
LIST
           KVA(I) = TEMP
                PRINT RESULTS
           PRINT FORMAT MESS1
           VECTOR VALUES MESS1 = $17H SORTED)$
           THROUGH PLIST, FOR K=1,1,K.G.N
PLIST
           PRINT FORMAT INT, KVL(K), KVA(K)
           END OF PROGRAM
           DATA
```