

Network Interference Management via **Interference Alignment**

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Network Interference Management via **Interference Alignment** (for Wireless Communications and Distributed Storage)

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m Equations, n Unknowns

$$\begin{aligned}y_1 &= v_{11}x_1 + v_{12}x_2 + \dots + v_{1n}x_n \\y_2 &= v_{21}x_1 + v_{22}x_2 + \dots + v_{2n}x_n \\&\vdots \\y_m &= v_{m1}x_1 + v_{m2}x_2 + \dots + v_{mn}x_n \\ \Leftrightarrow \mathbf{y} &= \mathbf{V}_1x_1 + \mathbf{V}_2x_2 + \dots + \mathbf{V}_nx_n\end{aligned}$$

We want to derive x_1, x_2, \dots, x_n from y_1, y_2, \dots, y_m

m Equations, n Unknowns

$$y_1 = v_{11}x_1 + v_{12}x_2 + \dots + v_{1n}x_n$$

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\vdots

$$y_m = v_{m1}x_1 + v_{m2}x_2 + \dots + v_{mn}x_n$$

$$\Leftrightarrow \mathbf{y} = \mathbf{V}_1x_1 + \mathbf{V}_2x_2 + \dots + \mathbf{V}_nx_n$$

We want to derive x_1, x_2, \dots, x_n from y_1, y_2, \dots, y_m

✓ $m \geq n$, $\mathbf{V}_1, \mathbf{V}_2, \dots, \mathbf{V}_n$ linearly independent

$m < n$,

m Equations, n Unknowns

$$y_1 = v_{11}x_1 + v_{12}x_2 + \dots + v_{1n}x_n$$

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We want to derive x_1, x_2, \dots, x_n from y_1, y_2, \dots, y_m

✓ $m \geq n$, $\mathbf{V}_1, \mathbf{V}_2, \dots, \mathbf{V}_n$ linearly independent

✗ $m < n$, In general, not resolvable

m Equations, n Unknowns - A slightly different problem

$$\mathbf{y} = \mathbf{V}_1 x_1 + \mathbf{V}_2 x_2 + \dots + \mathbf{V}_n x_n$$

We want to derive x_1 from y_1, y_2, \dots, y_m

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If $m < n$, then ...??

m Equations, n Unknowns - A slightly different problem

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If $m < n$, then ...?? ← Interference Alignment

Interference Alignment

Example: No. of Equations = 2, No. of Unknowns = 3

Interference Alignment

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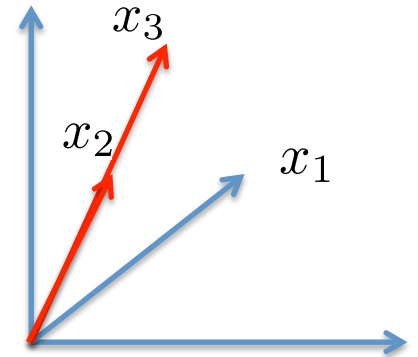
$$\begin{array}{r} 4 = x_1 + x_2 + 2x_3 \quad \times 2 \\ 7 = x_1 + 2x_2 + 4x_3 \quad \times (-1) \\ + \hline 1 = x_1 \end{array}$$

Interference Alignment

Example: No. of Equations = 2, No. of Unknowns = 3

$$\begin{pmatrix} 4 \\ 7 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \end{pmatrix} x_1 + \begin{pmatrix} 1 \\ 2 \end{pmatrix} x_2 + \begin{pmatrix} 2 \\ 4 \end{pmatrix} x_3$$

x_2, x_3 align



Interference Alignment

$$\mathbf{y} = \mathbf{V}_1 x_1 + \mathbf{V}_2 x_2 + \dots + \mathbf{V}_n x_n$$

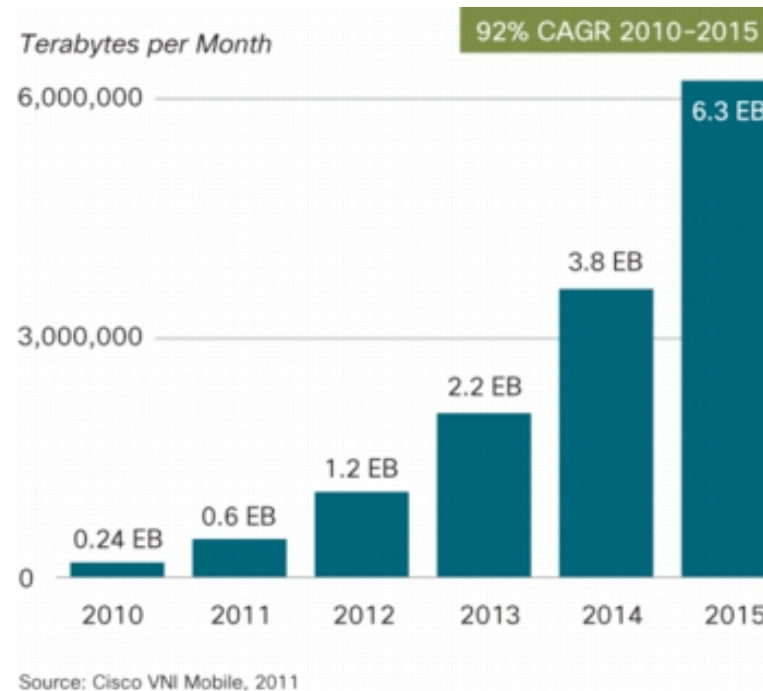
If $m < n$, then, to uniquely resolve x_1

$$\begin{aligned} \mathbf{V}_1 &\notin \text{span}(\{\mathbf{V}_2, \mathbf{V}_3, \dots, \mathbf{V}_n\}) \\ \Rightarrow \dim(\{\mathbf{V}_2, \mathbf{V}_3, \dots, \mathbf{V}_n\}) &\leq m - 1 \end{aligned}$$

Network Interference Management via **Interference Alignment** (for Wireless Communications and Distributed Storage)

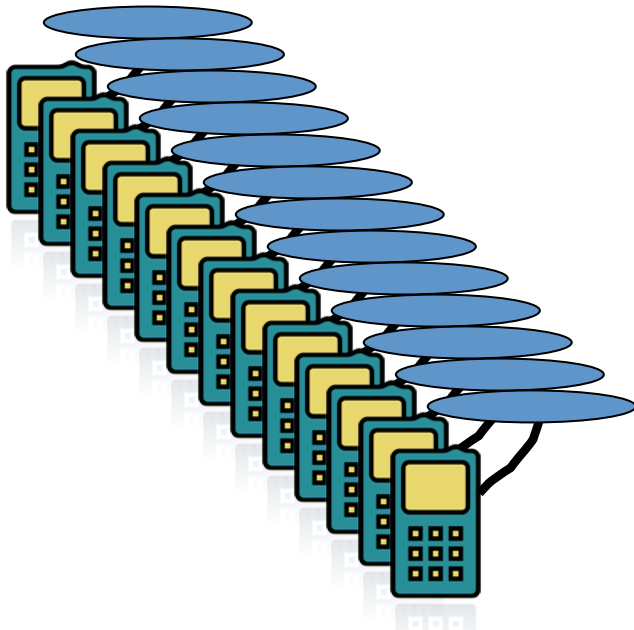
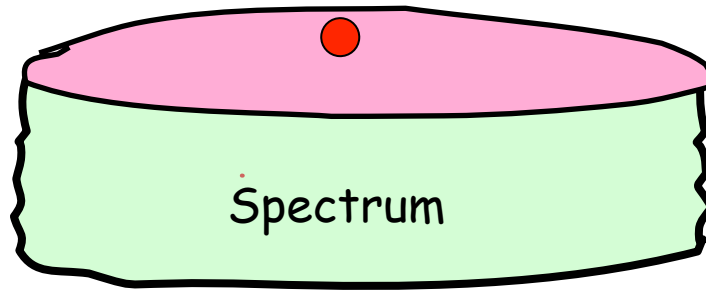
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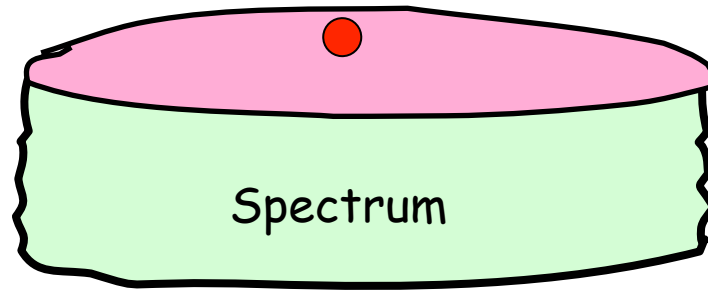
“Slim and sleek as it is, the iPhone is really the Hummer of cellphones.” *The NY Times*



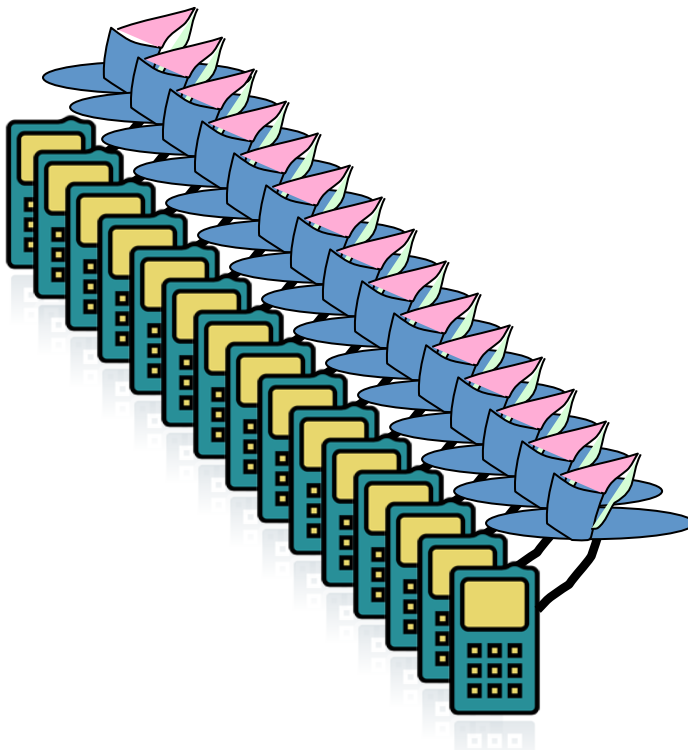
Worldwide mobile traffic projected to see a **26-fold increase** by 2015.

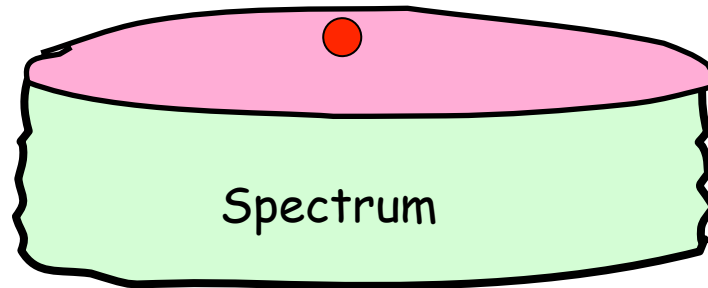
But the (usable) bandwidth remains the same!



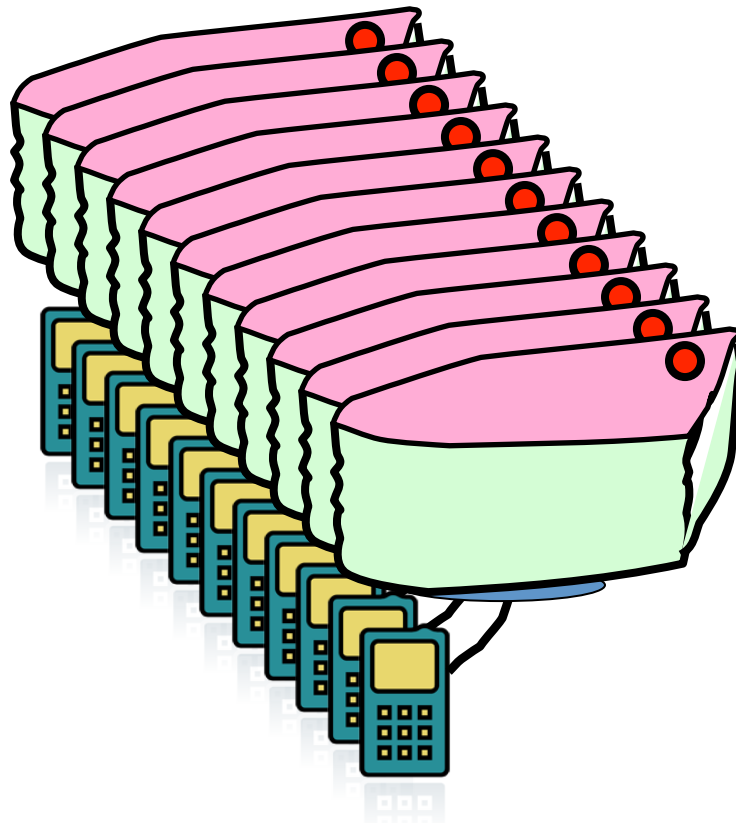


State of Art
Divide spectrum like a cake

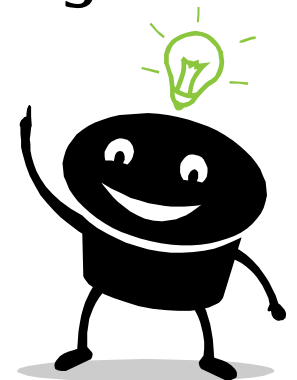




Everyone gets half the cake!!



Interference Alignment

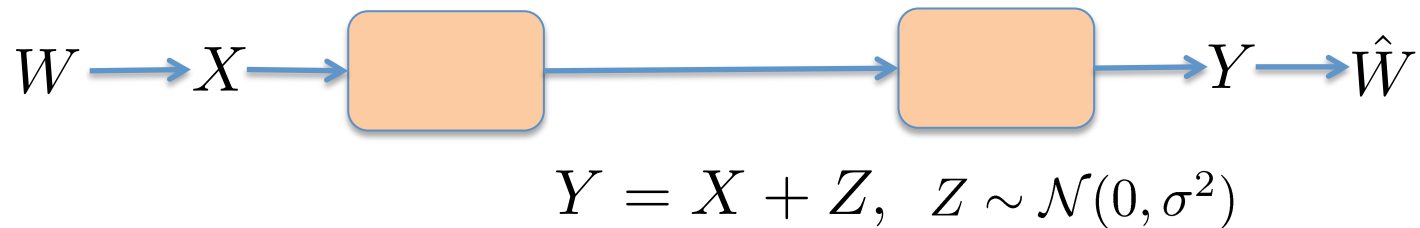




Claude E. Shannon



Claude E. Shannon

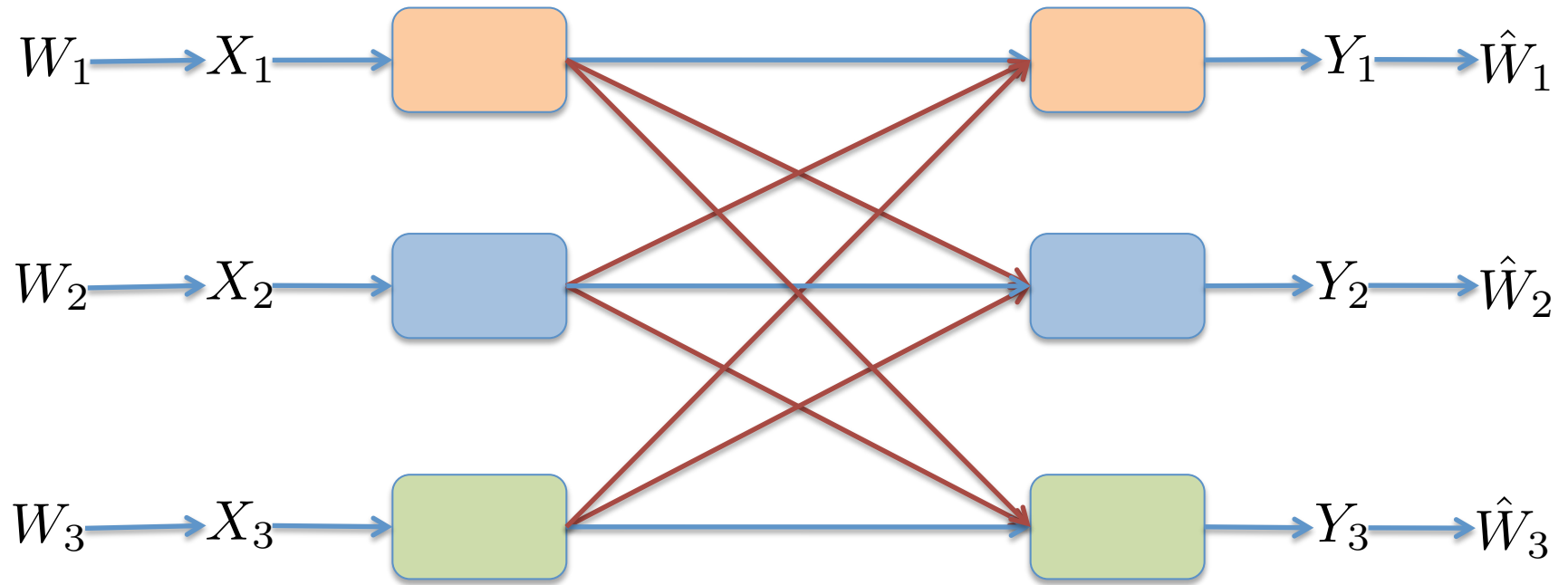


Max. Data Rate = Capacity = $B \log(1 + \text{SNR})$ bits/sec

B = Bandwidth

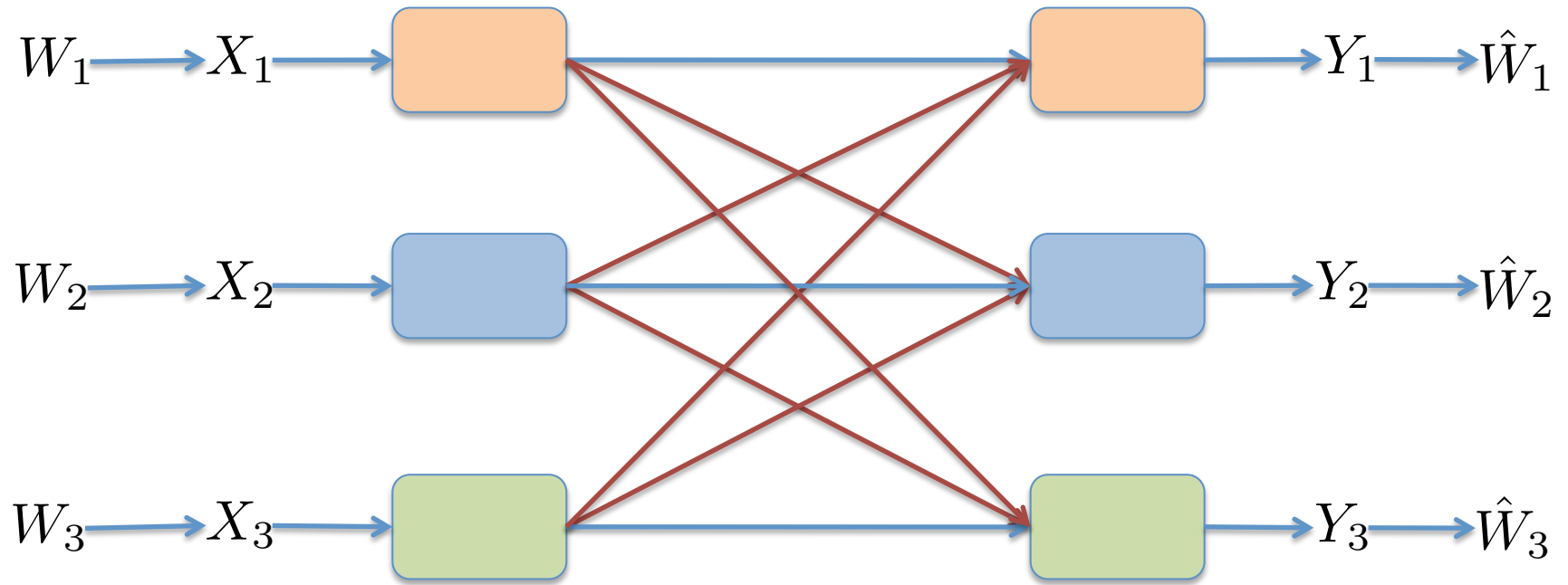
SNR = Ratio of Signal Power to Noise Power

K users sharing the same bandwidth and space



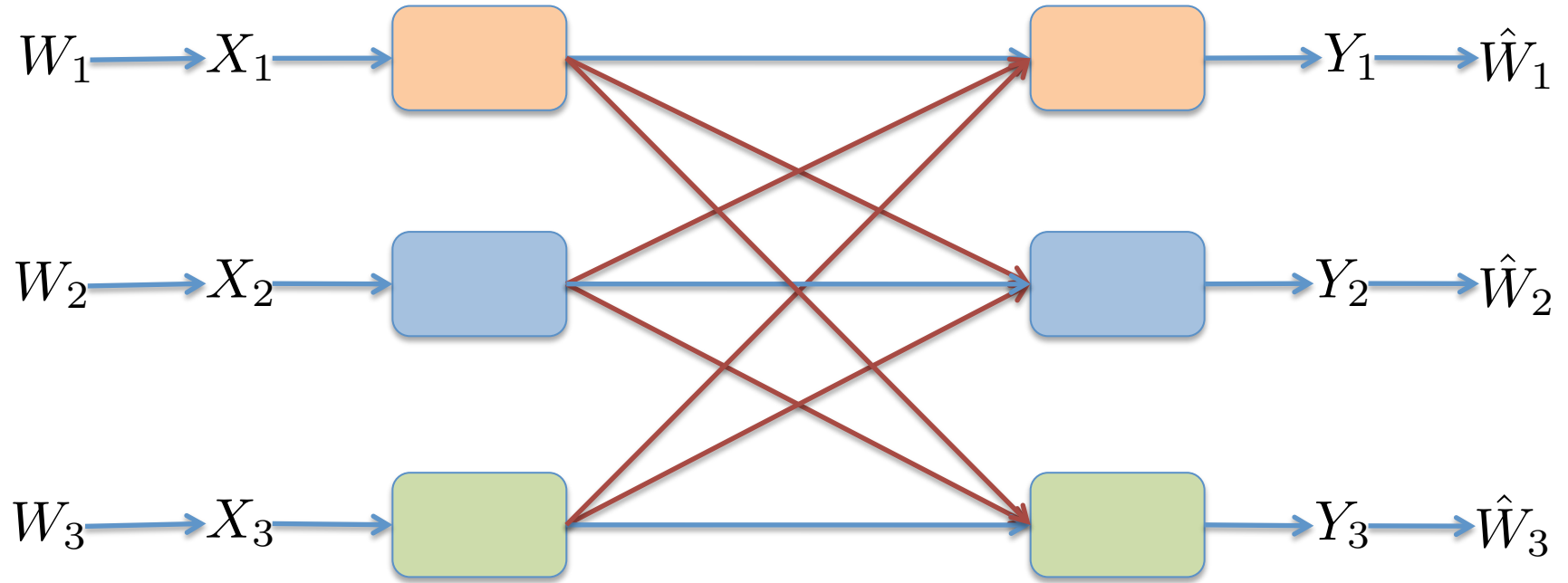
$$Y_i = \sum_{j,i} H_{j,i} X_i + Z_i$$

The K -User Interference Network



$$Y_i = \sum_{j,i} H_{j,i} X_i + Z_i$$

The K -User Interference Network



$$Y_i = \sum_{j,i} H_{j,i} X_i + Z_i$$

Capacity, 40 year old open problem, even for $K=2$!

Let us look at Capacity Approximations!

$$\text{Rate} = B \log(1 + \text{SNR})$$

$$\text{Rate} = B \log(1 + \text{SNR})$$

Degrees of Freedom = Capacity-Prelog

[Zheng-Tse 02]

Degrees of Freedom = Bandwidth Utilization Factor = Capacity-Prelog

[Zheng-Tse 02]

High SNR approximation of Network Capacity

$$\text{Degrees of Freedom} = \lim_{\text{SNR} \rightarrow \infty} \frac{\text{Network Capacity}}{\text{Single User Capacity}}$$

A network has d degrees of freedom iff

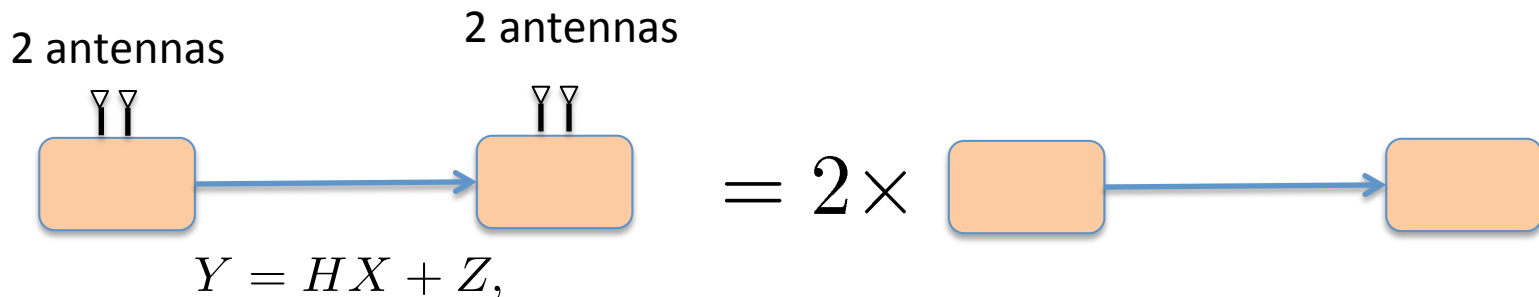
$$C(\text{SNR}) = dB \log(\text{SNR}) + o(\log(\text{SNR}))$$

$$\text{SNR} = \frac{\text{Total Network Signal Power}}{\text{Total Network Noise Power}}$$

B = Total Bandwidth available

De-emphasize noise, emphasize Interference

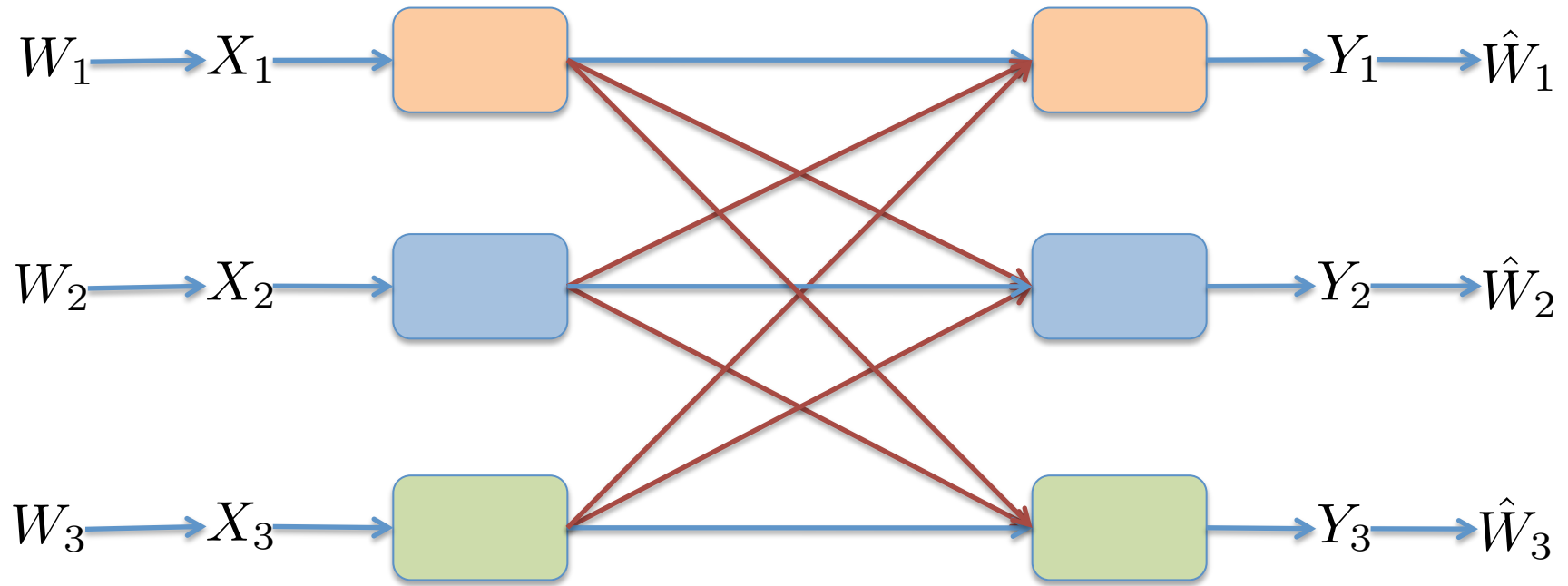
Intuitively, number of interference-free links (at high SNR).



$$C(\text{SNR}) = 2 \cdot B \log(\text{SNR}) + o(\log(\text{SNR}))$$

[Telatar 95]

The K -User Interference Network



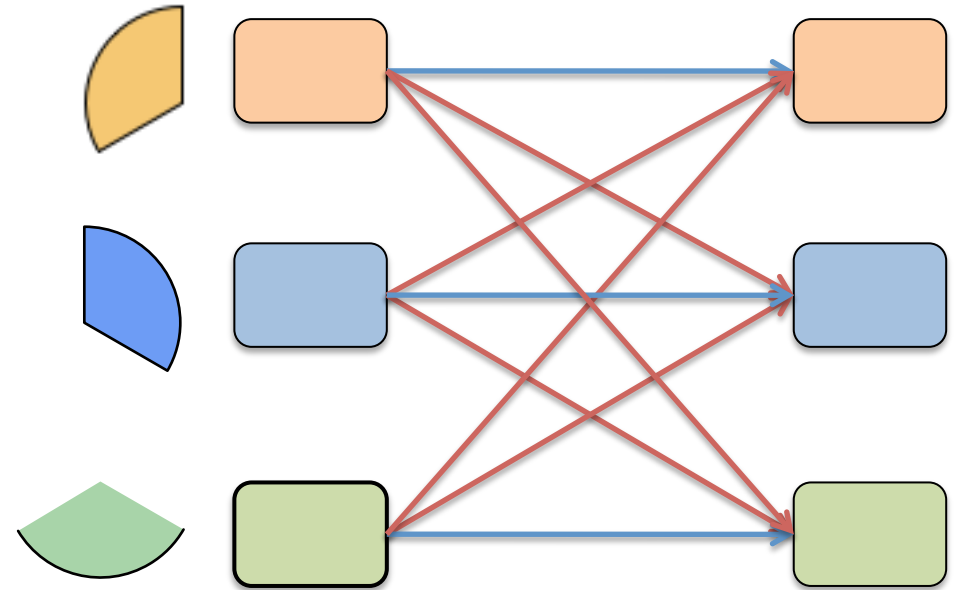
$$Y_i = \sum_{j,i} H_{j,i} X_i + Z_i$$

Question : What is the number of Degrees of Freedom?

The K -User Interference Network - Current Design



Spectrum





Orthogonalization (cake-cutting)

Each user gets $\frac{1}{K}$ of the cake

$$\text{Per user capacity} = \frac{1}{K} B \log(\text{SNR}) + o(\log(\text{SNR}))$$

SUBOPTIMAL

$$\text{Sum-Rate} = 1 \cdot B \log(\text{SNR}) + o(\log(\text{SNR}))$$

~~Conjectured to be DoF optimal~~



[Cadambe-Jafar 08]

Interference Alignment

Each user gets **half** the cake

$$\text{Per-User Capacity} = \frac{1}{2} B \log(\text{SNR}) + o(\log(\text{SNR}))$$

$$\text{Sum-Capacity} = \frac{K}{2} B \log(\text{SNR}) + o(\log(\text{SNR}))$$

Cake = Bandwidth-utilization factor of an interference-free user (high SNR).

Toy Example for Interference Alignment

[Cadambe-Jafar 08]

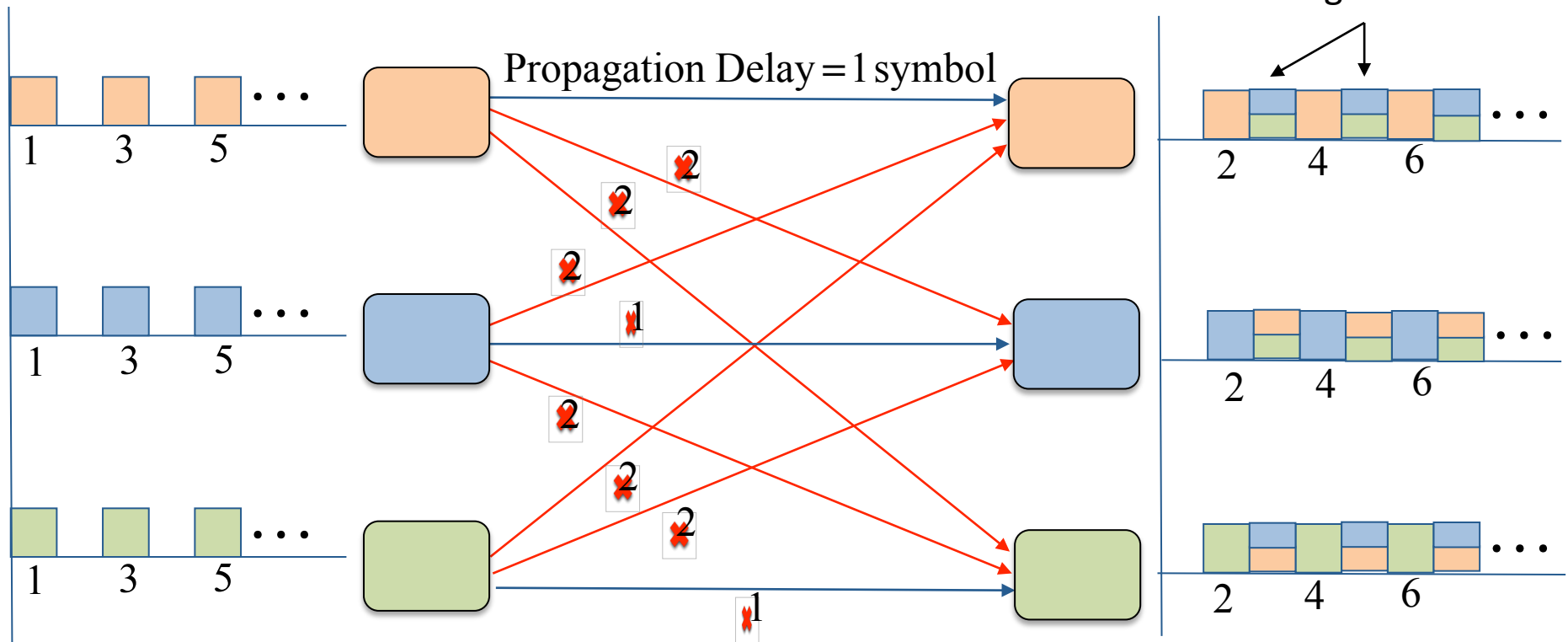
K speaker-listener pairs in a room, K conversations
How long can each conversation be active in one hour ?



Toy Example for Interference Alignment

[Cadambe-Jafar 08]

K speaker-listener in a room, K conversations
How long can each conversation be active in one hour ?



Each user gets **half the cake** because of interference alignment.

What is Interference Alignment?

[Maddah-Ali-Motahari-Khandani 06, Jafar-Shamai 06]

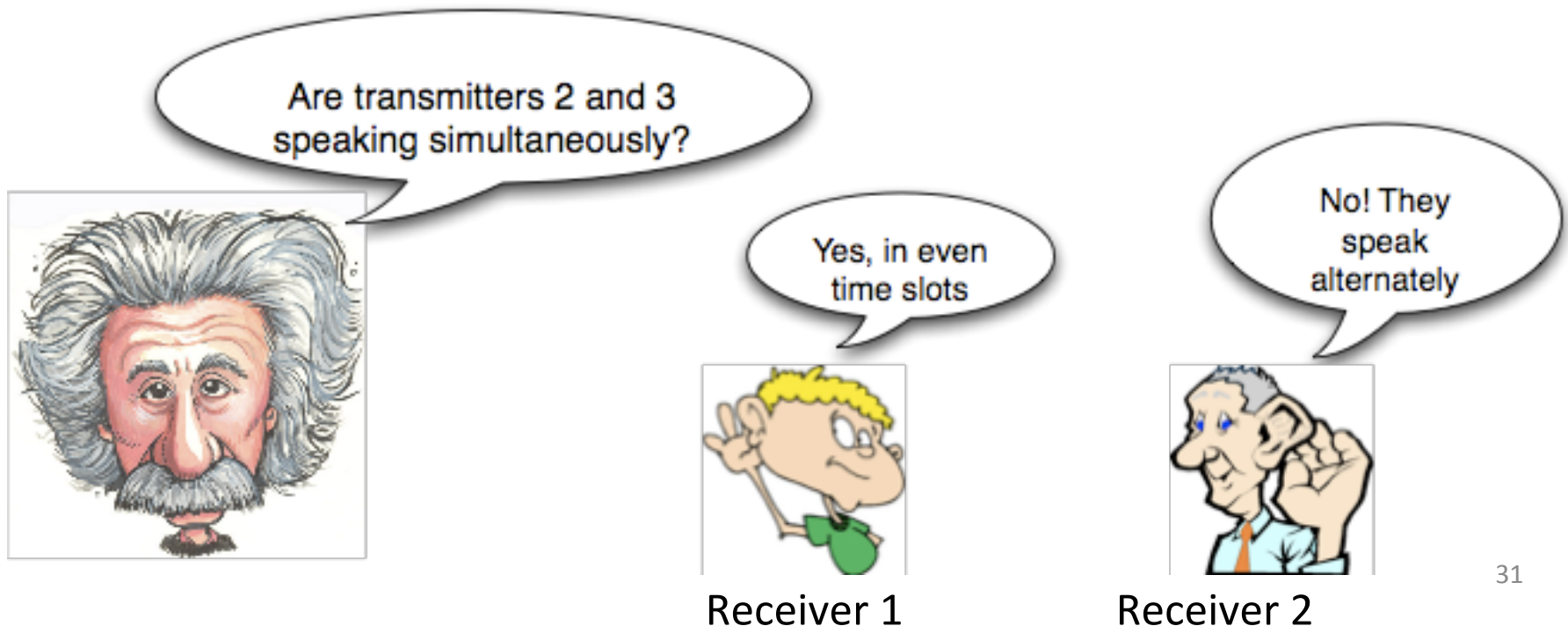
Signals are designed so that, at the receivers

Interfering signals overlap,

Desired signals are distinguishable

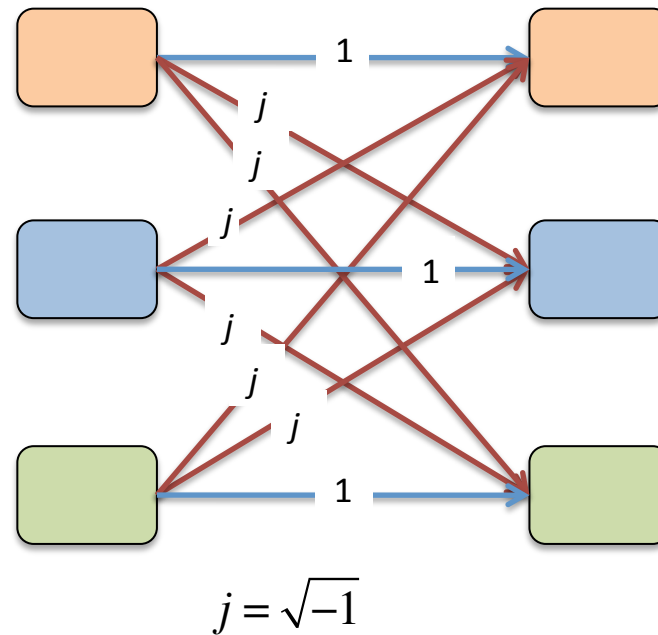
Works because

Each receiver sees a different picture



Toy Example- Interference Alignment in Phase

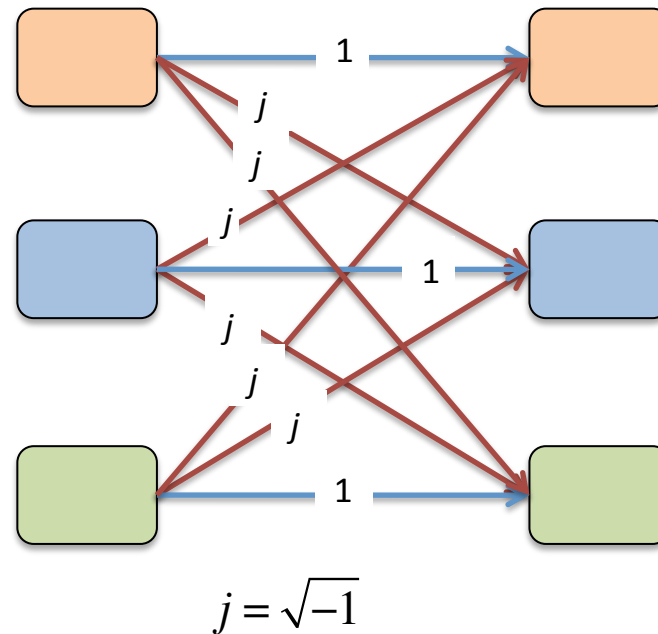
[Cadambe-Jafar 08]



Single-User Capacity = $B \log(1 + \text{SNR})$, Cake= B

Toy Example- Interference Alignment in Phase

[Cadambe-Jafar 08]



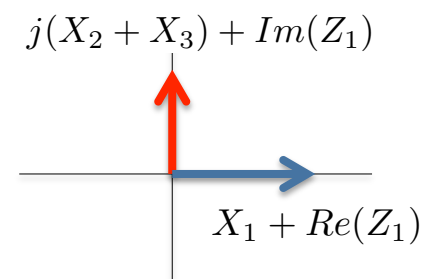
Single-User Capacity = $B \log(1 + \text{SNR})$, Cake = B

Transmit **Real** Symbols.

Interference **Aligns** along imaginary dimension.

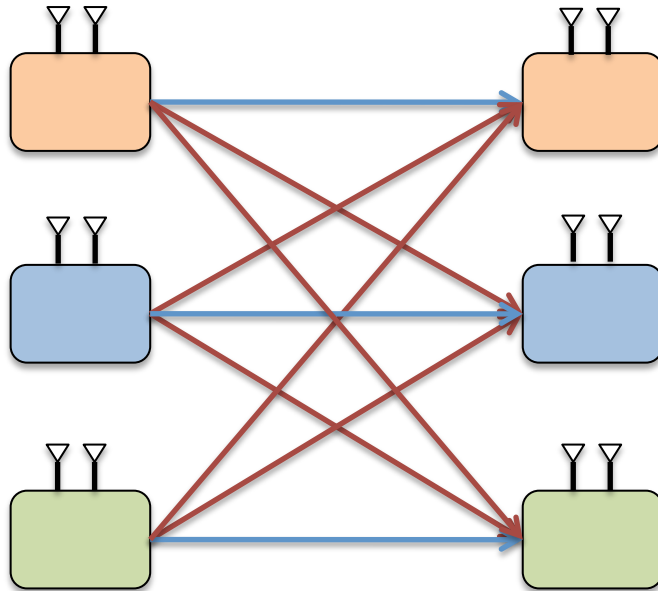
Each user gets **half** the cake!

$$\text{per-user-capacity} = \frac{1}{2} B \log(\text{SNR}) + o(\log(\text{SNR}))$$



3 Users, 2 antennas at each transmitter and receiver

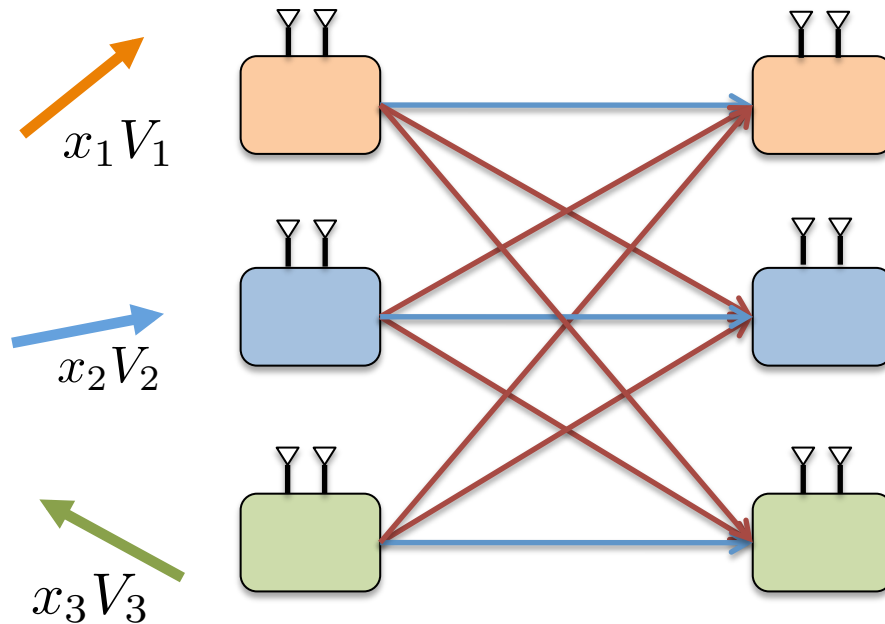
[Cadambe-Jafar 08]



Single-User Capacity = $2B \log(\text{SNR}) + o(\log(\text{SNR}))$, Cake = $2B$

3 Users, 2 antennas at each transmitter and receiver

[Cadambe-Jafar 08]



Single-User Capacity = $2B \log(\text{SNR}) + o(\log(\text{SNR}))$, Cake = $2B$

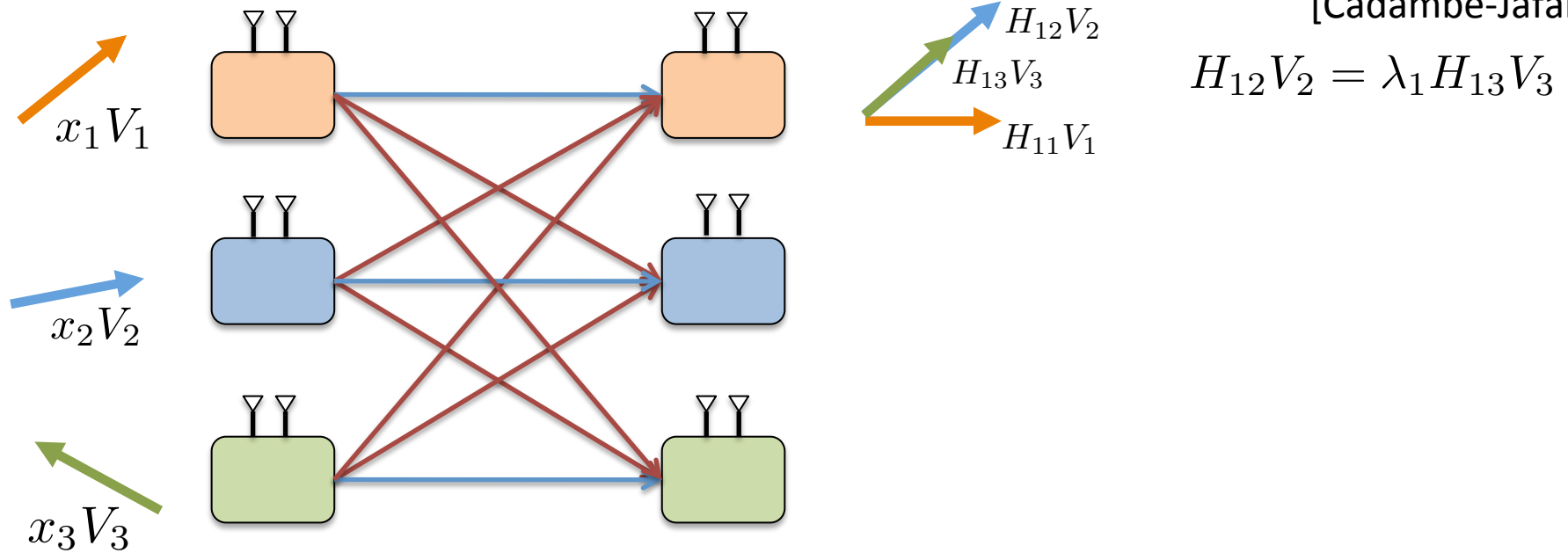
With Interference Alignment,

Each user gets **half** the cake!

Each user gets $1 \cdot B \log(\text{SNR}) + o(\log(\text{SNR}))$

3 Users, 2 antennas at each transmitter and receiver

[Cadambe-Jafar 08]



Single-User Capacity = $2B \log(\text{SNR}) + o(\log(\text{SNR}))$, Cake = $2B$

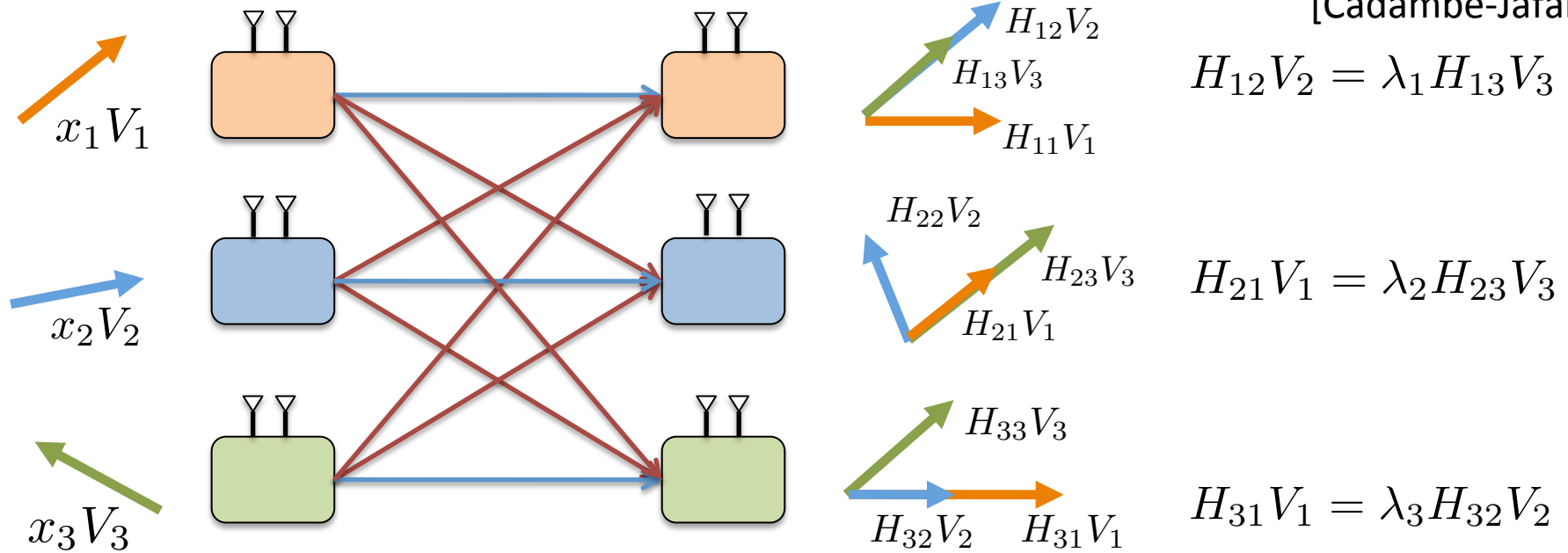
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[Cadambe-Jafar 08]



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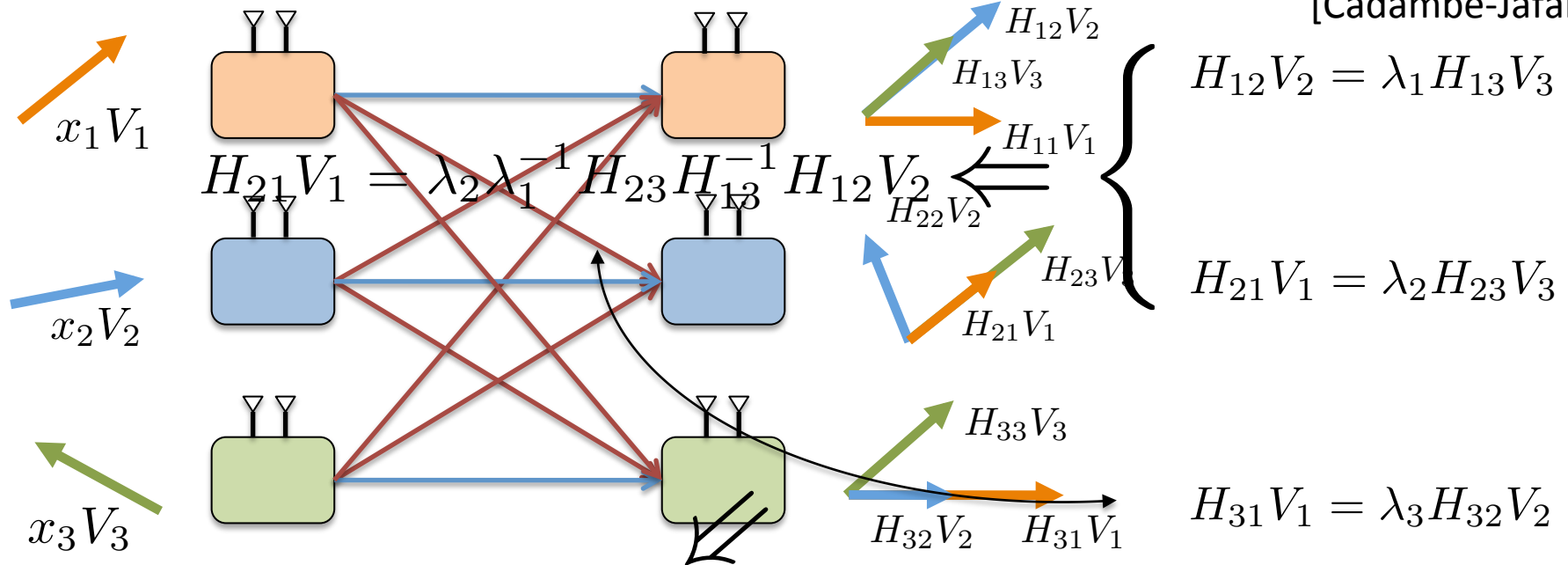
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3 Users, 2 antennas at each transmitter and receiver

[Cadambe-Jafar 08]



$$V_1 = \lambda T V_1$$

Single-User Capacity = $2B \log(\text{SNR}) + o(\log(\text{SNR}))$, Cake = $2B$

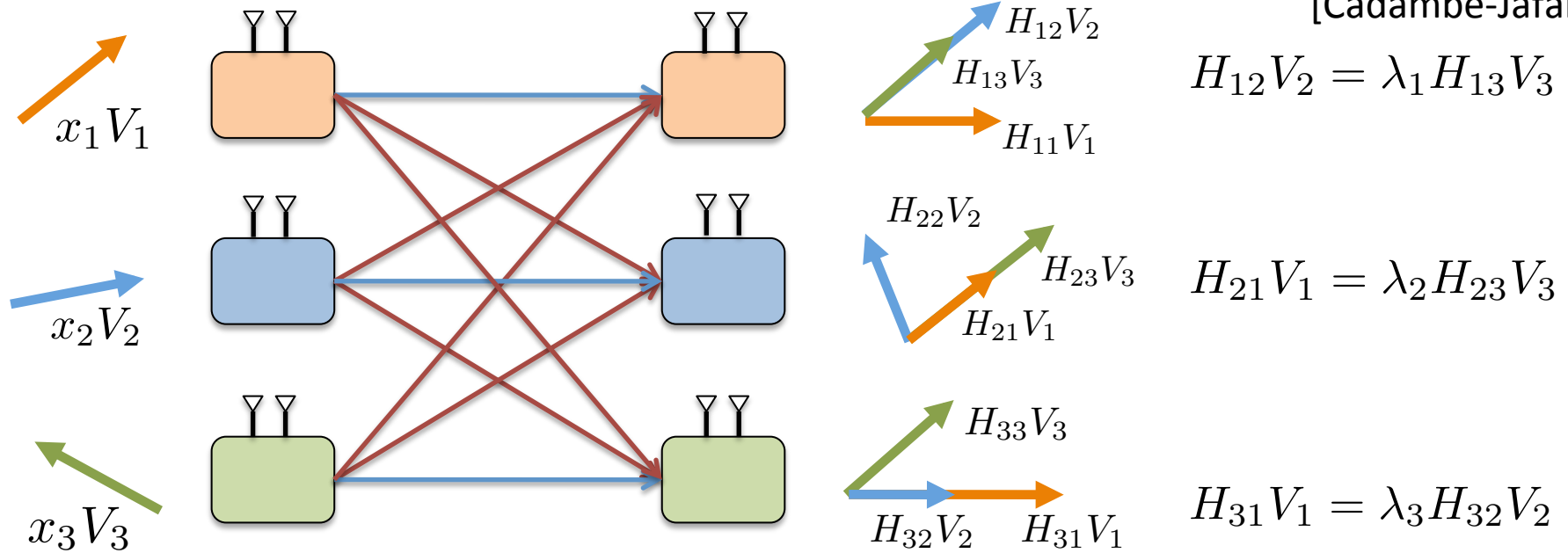
V_1 eigen-vector of T
 With Interference Alignment,

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3 Users, 2 antennas at each transmitter and receiver

[Cadambe-Jafar 08]



Single-User Capacity = $2B \log(\text{SNR}) + o(\log(\text{SNR}))$, Cake = $2B$

With Interference Alignment,

Each user gets **half** the cake!

Each user gets $1 \cdot B \log(\text{SNR}) + o(B \log(\text{SNR}))$

What about more than 3 users?

For every set of 3 users,

$$V_1 = \lambda TV_1$$

What about more than 3 users?

$$\text{span}(V) = \text{span}(T_1 V)$$

$$\text{span}(V) = \text{span}(T_2 V)$$

$$\vdots$$

$$\text{span}(V) = \text{span}(T_N V)$$

$$\text{rank}(V) = \frac{\text{rank}(T_i)}{2}$$

V invariant subspace of T_1, T_2, \dots, T_N

Infeasible!



How about finding approximately invariant subspaces?

Asymptotic Interference Alignment

[Cadambe-Jafar 08]

$$\begin{aligned} \text{span}(V) &\approx \text{span}(T_1 V) \\ \text{span}(V) &\approx \text{span}(T_2 V) \\ &\vdots \\ \text{span}(V) &\approx \text{span}(T_N V) \end{aligned} \quad \text{rank}(V) = \frac{\text{rank}(T_i)}{2}$$

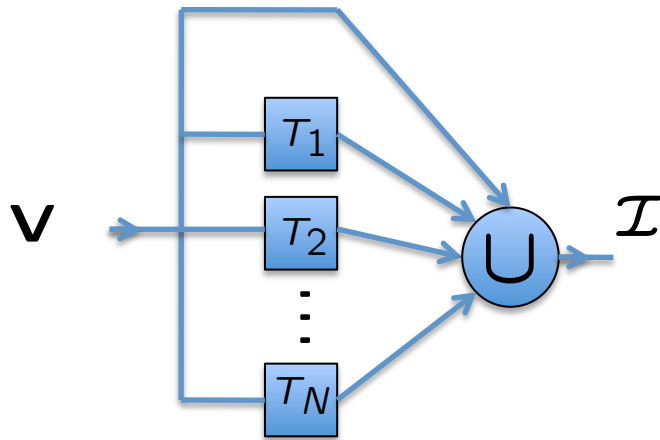
Asymptotic Solution \rightarrow Diagonal T_i s of arbitrarily large dimension

Diagonal T_i s \rightarrow Code over multiple frequencies (**Vector Coding**)
Frequency-Selective Channels

Goal: Simultaneously satisfy “N” Alignment Constraints:
 $\text{span}(\mathbf{V}) \equiv \text{span}(T_1\mathbf{V}) \equiv \text{span}(T_2\mathbf{V}) \equiv \cdots \equiv \text{span}(T_N\mathbf{V})$

[Cadambe, Jafar, IT08]

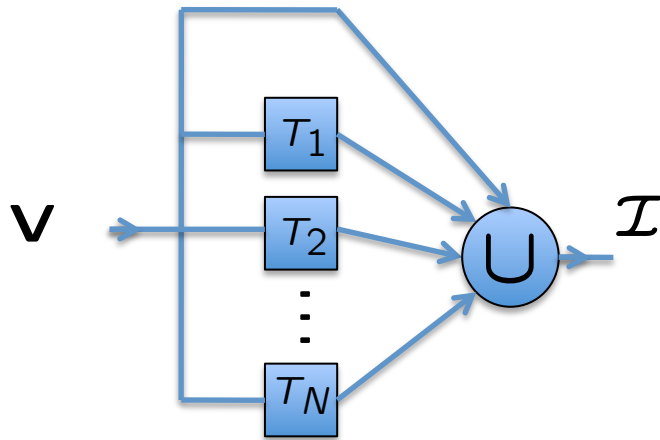
Goal: Make $\mathbf{V} \equiv \mathcal{I}$



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[Cadambe, Jafar, IT08]

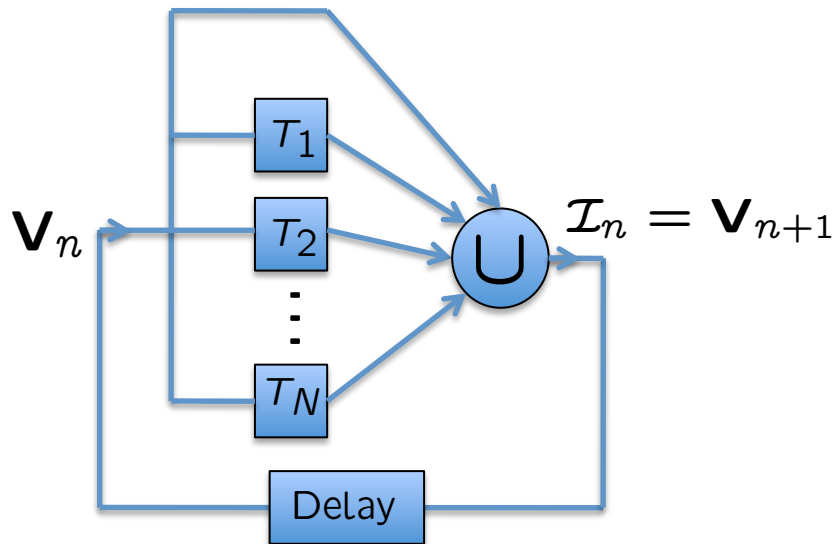
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[Cadambe, Jafar, IT08]

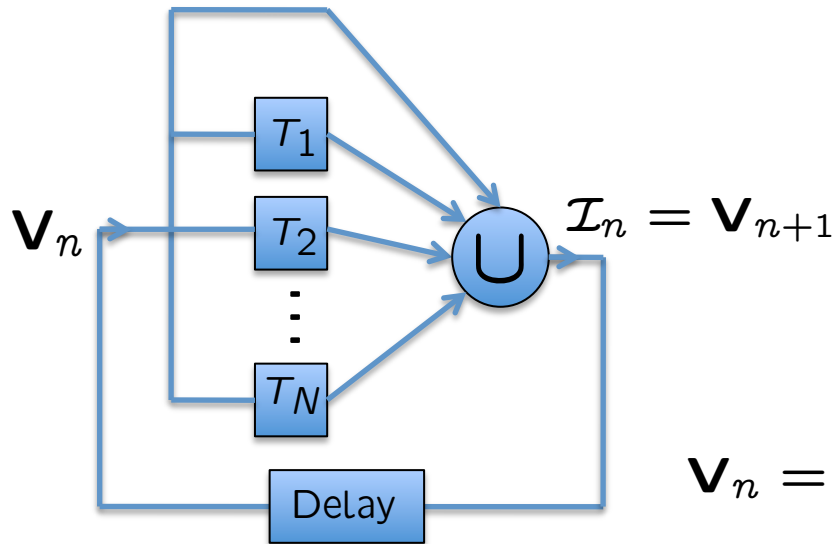
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[Cadambe, Jafar, IT08]

Goal: Make $\mathbf{V} \equiv \mathcal{I}$



Initialize: $\mathbf{V}_0 = \mathbf{1}$

$$\mathbf{V}_1 = \{\mathbf{1}, T_1\mathbf{1}, \dots, T_N\mathbf{1}\}$$

$$\mathbf{V}_2 = \{\mathbf{1}, \dots, T_i\mathbf{1}, \dots, T_i T_j\mathbf{1}, \dots, T_i^2\mathbf{1}\}$$

$$\mathbf{V}_n = \{T_1^{\alpha_1} T_2^{\alpha_2} \dots T_N^{\alpha_N} \mathbf{1}, \alpha_1 + \dots + \alpha_N \leq n\}$$

$$\mathbf{V}_{n+1} = \{T_1^{\alpha_1} T_2^{\alpha_2} \dots T_N^{\alpha_N} \mathbf{1}, \alpha_1 + \dots + \alpha_N \leq n + 1\}$$

$$|\mathbf{V}_n| = \binom{n+N}{n}$$

$$|\mathcal{I}| = \binom{n+N+1}{n+1}$$

$$\frac{|\mathbf{V}|}{|\mathcal{I}|} = \frac{n+1}{n+N+1} \rightarrow 1 \text{ as } n \rightarrow \infty$$

Key : T_i s commute

Asymptotic Interference Alignment

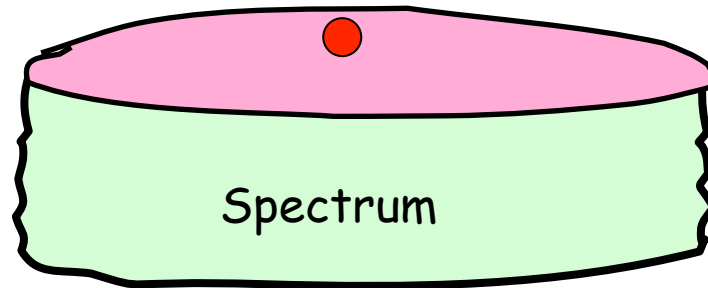
[Cadambe-Jafar 08]

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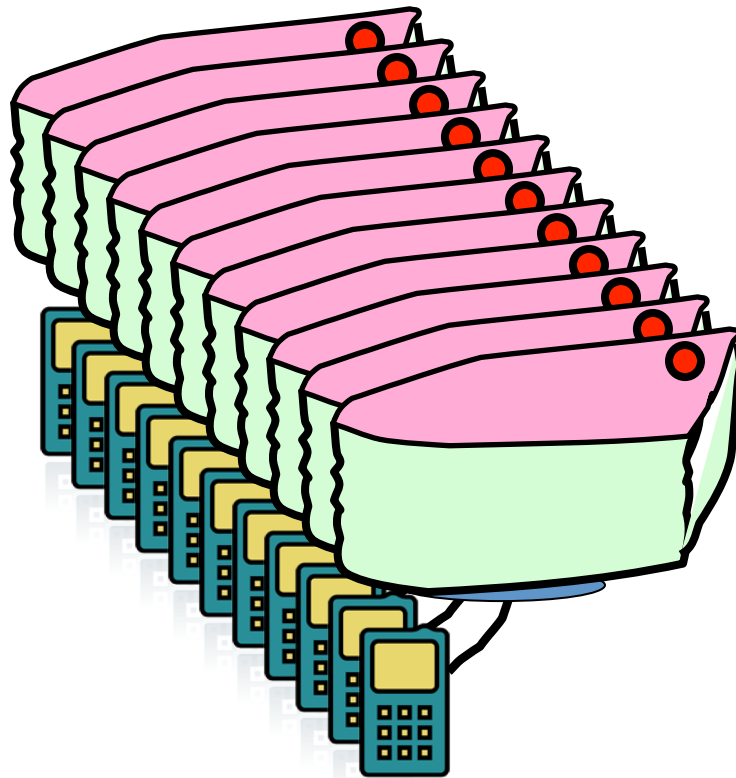
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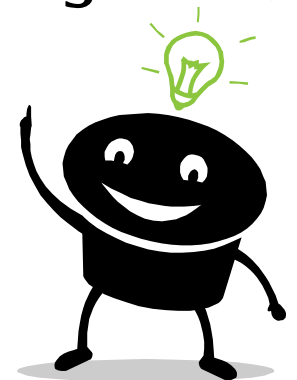
Can be done over constant (non-frequency-selective) channels by exploiting rational dimensions [Motahari et. al. 09]



Everyone gets half the cake!!



Interference Alignment



Asymptotic Alignment scheme widely applicable

X channels [Cadambe-Jafar 09], MIMO Int. Channels [Gou-Jafar 10], Compound Broadcast and Interference Networks [Wang et. al 10], Distributed Storage Problems [Cadambe-Jafar-Maleki 10, Suh-Ramchandran 10], Network Coding, [Ramkrishnan et. al. 10]

Design principles of current wireless systems inspired by

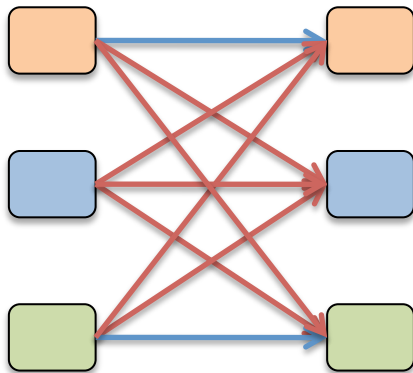


Claude E. Shannon

information theory of single user systems



Theoretical foundations of today's networks

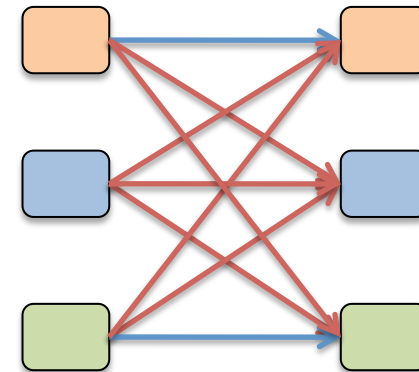


Random Codes

Circularly Symmetric Signaling

Multicarrier Systems,
Whole = sum of parts

New Theoretical Insights for Networks because of alignment



Structured Codes*

Asymmetric Signaling*

Multicarrier Systems,
Whole > sum of parts

*[Cadambe-Jafar-Shamai 08, Nazer-Gastpar 08, Bresler-Parekh-Tse 07, Cadambe-Jafar 09, Cadambe-Jafar-Wang 09, Cadambe-Jafar-08, Shankar et. al 08, Cadambe-Jafar 10]

Single User Channels versus Multi-User Channels

Main Challenge in Single-User Channels - Robustness to Noise

Insight : Use Random codes. [Shannon 1948]

(Eg. Turbo Codes)

Main Challenge in Multi-User Channels - Robustness to Interference

Insight : Interference Alignment, **Structured** codes!

(Eg. Lattice Codes)

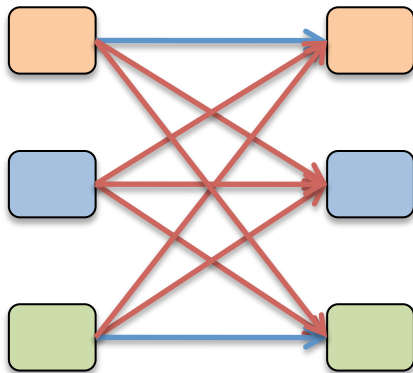
Random codes do not align interference

[Cadambe-Jafar-Shamai Trans. IT 08, Nazer-Gastpar 08, Bresler-Parekh-Tse Trans. IT 10]

except under special circumstances.

[Cadambe-Jafar Allerton-10]

Theoretical foundations of today's networks

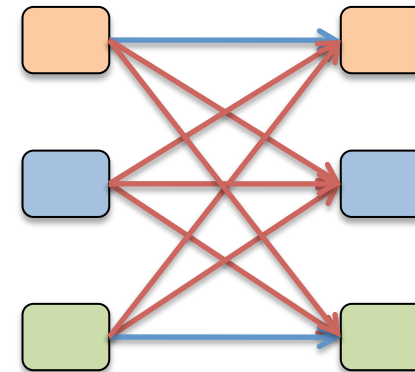


Random Codes

Circularly Symmetric Signaling

Multicarrier Systems,
Whole = sum of parts

New Theoretical Insights for Networks because of alignment



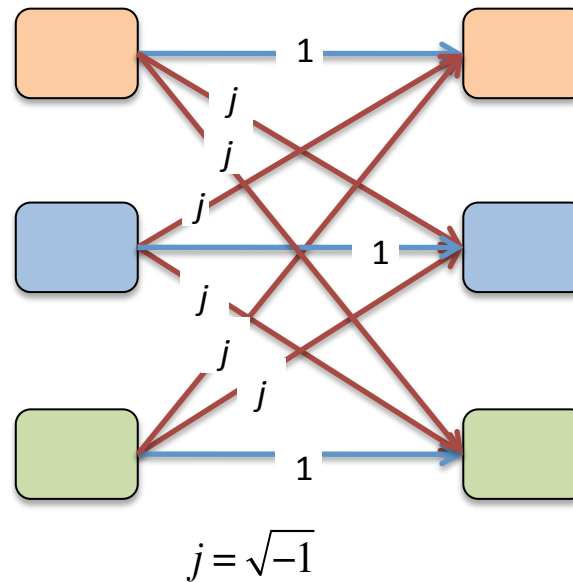
Structured Codes*

Asymmetric Signaling*

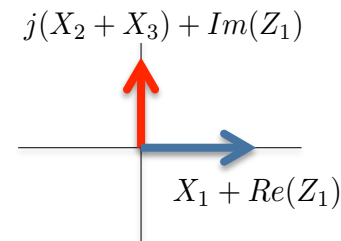
Multicarrier Systems,
Whole > sum of parts*

*[Cadambe-Jafar-Shamai 08, Nazer-Gastpar 08, Bresler-Parekh-Tse 07, Cadambe-Jafar 09, Cadambe-Jafar-Wang 09, Cadambe-Jafar-08, Shankar et. al 08, Cadambe-Jafar 10]

Interference Alignment in Phase



Transmit **Real** Symbols.
 Interference **Aligns** along imaginary dimension.

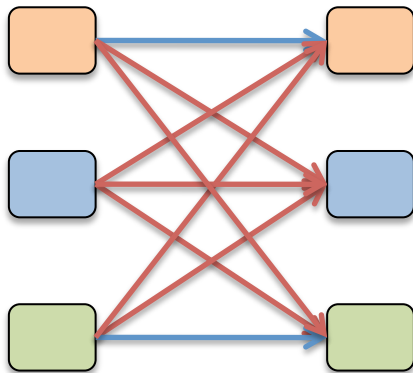


Unlike single user channels,

Asymmetric Complex Signals are useful in aligning interference

Characterized the gains of asymmetric complex signaling for networks with *random* channels

Theoretical foundations of today's networks

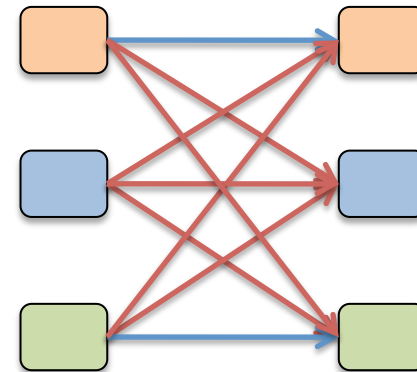


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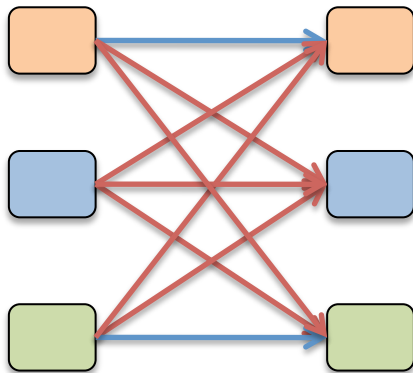
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Theoretical foundations of today's networks

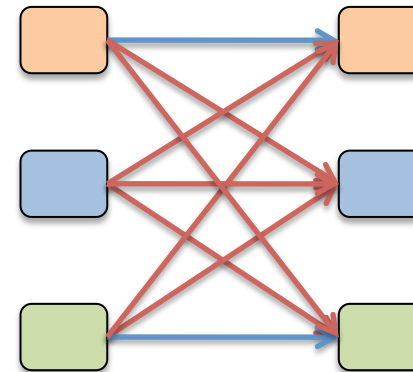


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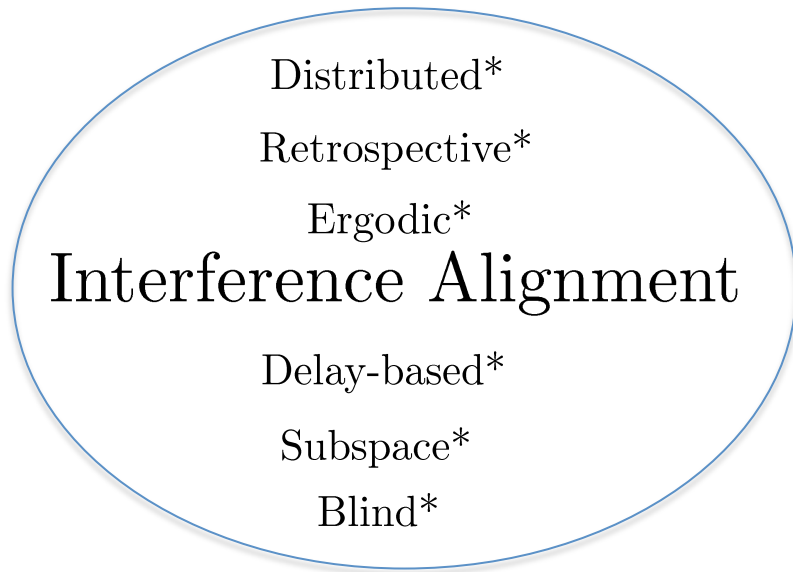
The road ahead - making Interference Alignment Practical

Interference Alignment

faces challenges of

Limited SNR, Channel State Information,
Synchronization, Signaling Dimensions

The road ahead - making Interference Alignment Practical

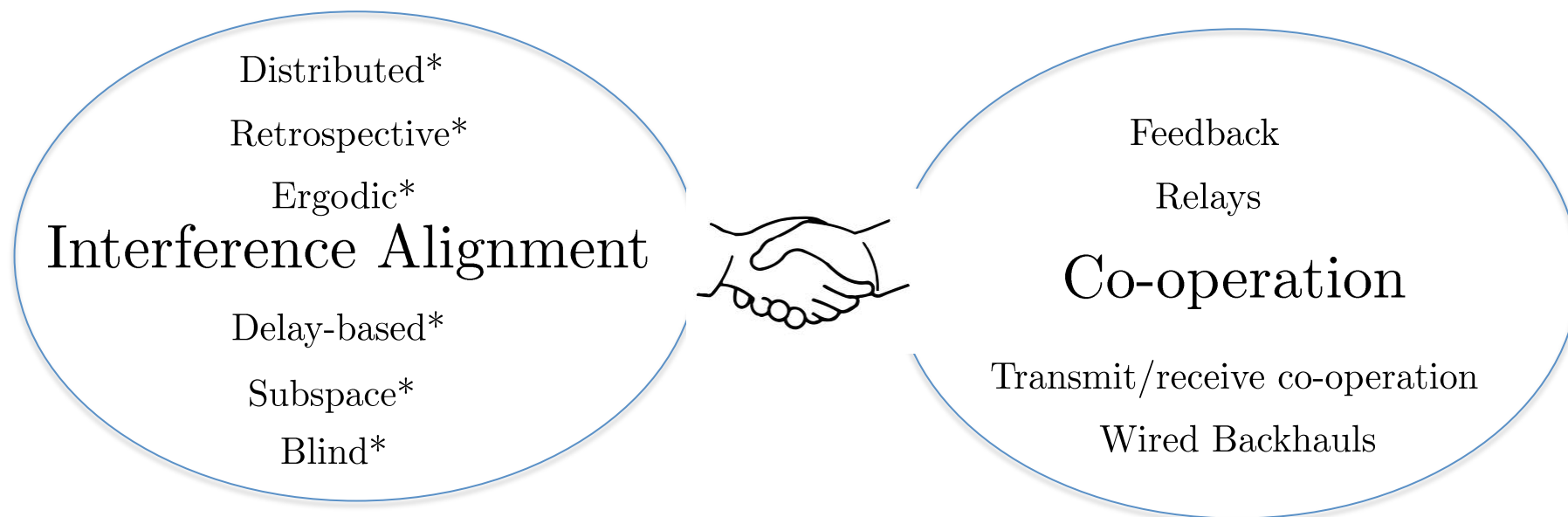


faces challenges of

Limited SNR, Channel State Information,
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*[Cadambe-Jafar 07, 08, Suh-Tse 08, Gropop-Tse-Yates 08, Nazer et. al. 09, Jafar 09, Wang-Guo-Jafar 10, Maddah-Ali-Tse 10, Maleki-Jafar-Shamai 10]

The road ahead - making Interference Alignment Practical



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Network Interference Management via **Interference Alignment** (for Wireless Communications and Distributed Storage)

Viveck R. Cadambe
University of California, Irvine

Servers (RAID-type systems)



Servers (RAID-type systems)



Data Centers



Servers (RAID-type systems)



Data Centers

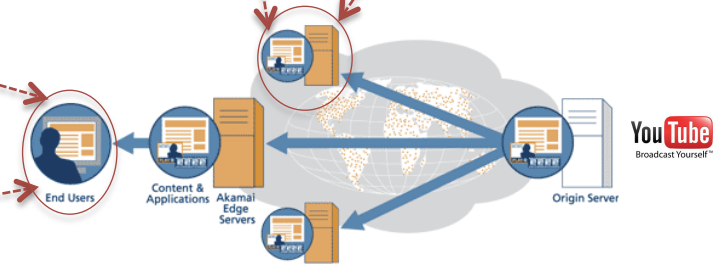


Content Delivery Networks

Servers (RAID-type systems)



Data Centers



Peer-to-Peer Storage

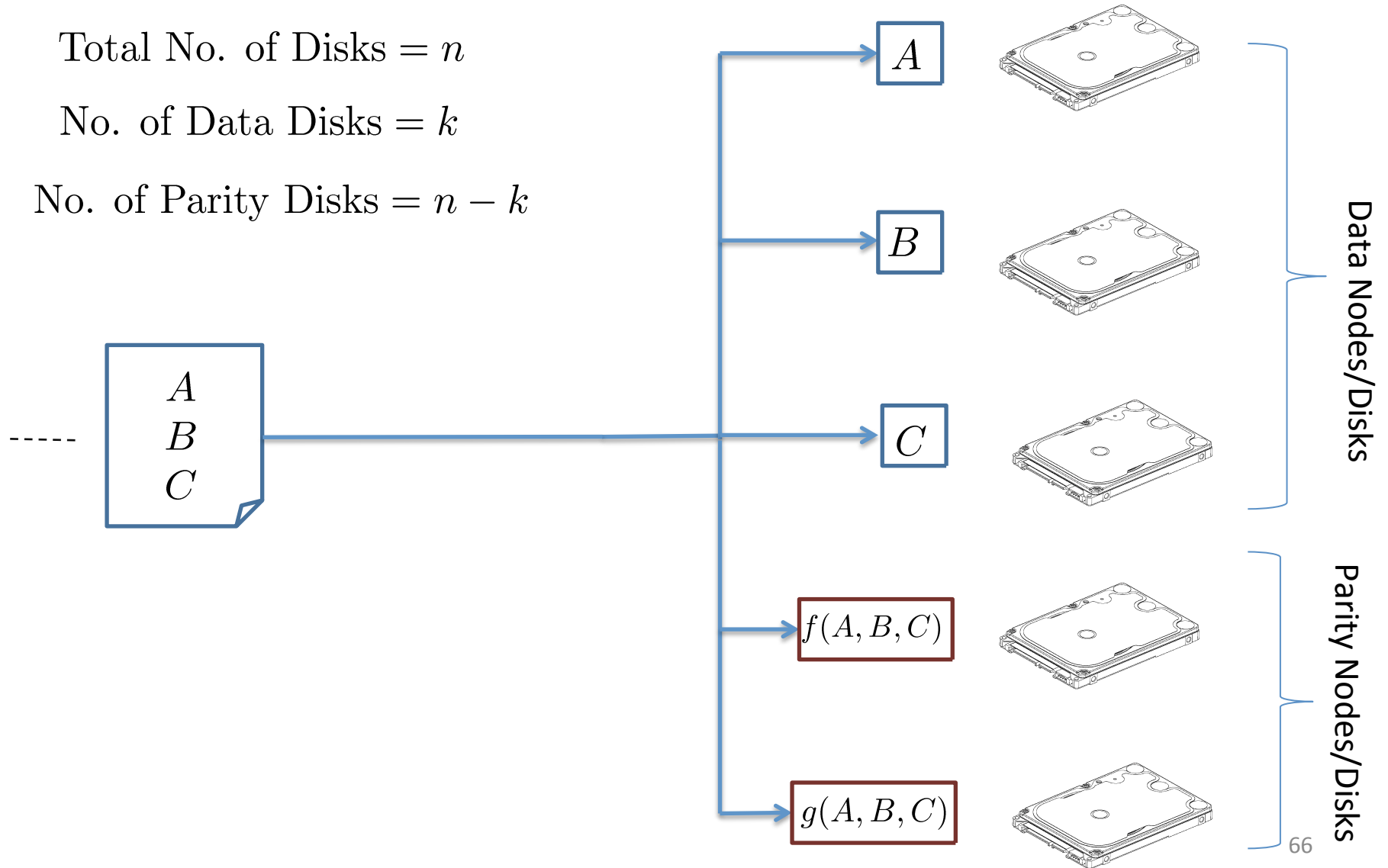
Content Delivery Networks

Distributed Storage System

Total No. of Disks = n

No. of Data Disks = k

No. of Parity Disks = $n - k$



Trade-offs in Distributed Storage Systems

Reliability

v/s

 **Cost of Storage**



“My data is my life...”

Replication widely used today.



Use Erasure Coding

Replication

$$n = 4$$

$$k = 2$$

$A, B \rightarrow 1$ unit

A

B

A

B

MDS* Codes

A

B

$A + B$

$A + 2B$

1 Disk Failure	Download 1 unit 😊	Download 2 unit 😞
2 Disk Failure	Data Loss 😞	Download 2 unit 😊

*MDS = Maximum Distance Separable⁶⁸

Replication

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*MDS = Maximum Distance Separable⁶⁹

(n, k) MDS Code - Pros and Cons

1 Disk Fails	Download k units
$n - k$ Disks Fail	Download k units



Highest redundancy (from worst-case failure perspective)



Slow Recovery of single disk failure

- Single node failure, most common failure scenario
- Recovery important before other nodes fail, so speed of recovery matters.
- Avoid Network Congestion!

(n, k) MDS Code - Pros and Cons

1 Disk Fails	Download k units
--------------	--------------------

Reduced? If so, how small? [Dimakis et. al 08, Wu-Dimakis 09]

Special case of multi-source (non-multicast) network coding

(n, k) MDS Code - Pros and Cons

1 Disk Fails	Download k units
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Reduced? If so, how small? [Dimakis et. al 08, Wu-Dimakis 09]

Special case of multi-source (non-multicast) network coding

Yes*, download only $\frac{n-1}{n-k} < k$ units

[Cadambe-Jafar-Maleki 10, Suh-Ramchandran 10]

Matches cut-set bound of [Dimakis et. al. 08]

$(n = 9, k = 6)$ Speed of repair \rightarrow doubly faster

*[Suh-Ramchandran 09, Shah et. al. 08] previously solved it for $n \geq k/2$

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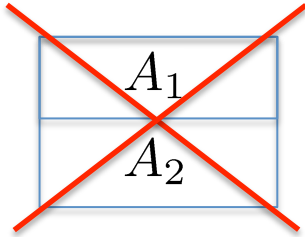
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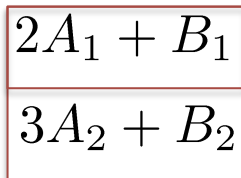
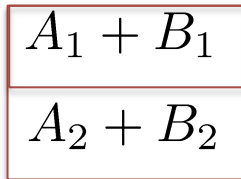
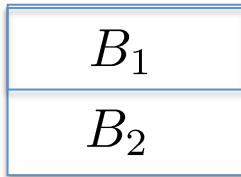
Use Interference Alignment, Vector Coding

*[Suh-Ramchandran 09, Shah et. al. 08] previously solved it for $n \geq k/2$

$$n = 4, k = 2 \quad [\text{Wu-Dimakis 09}]$$

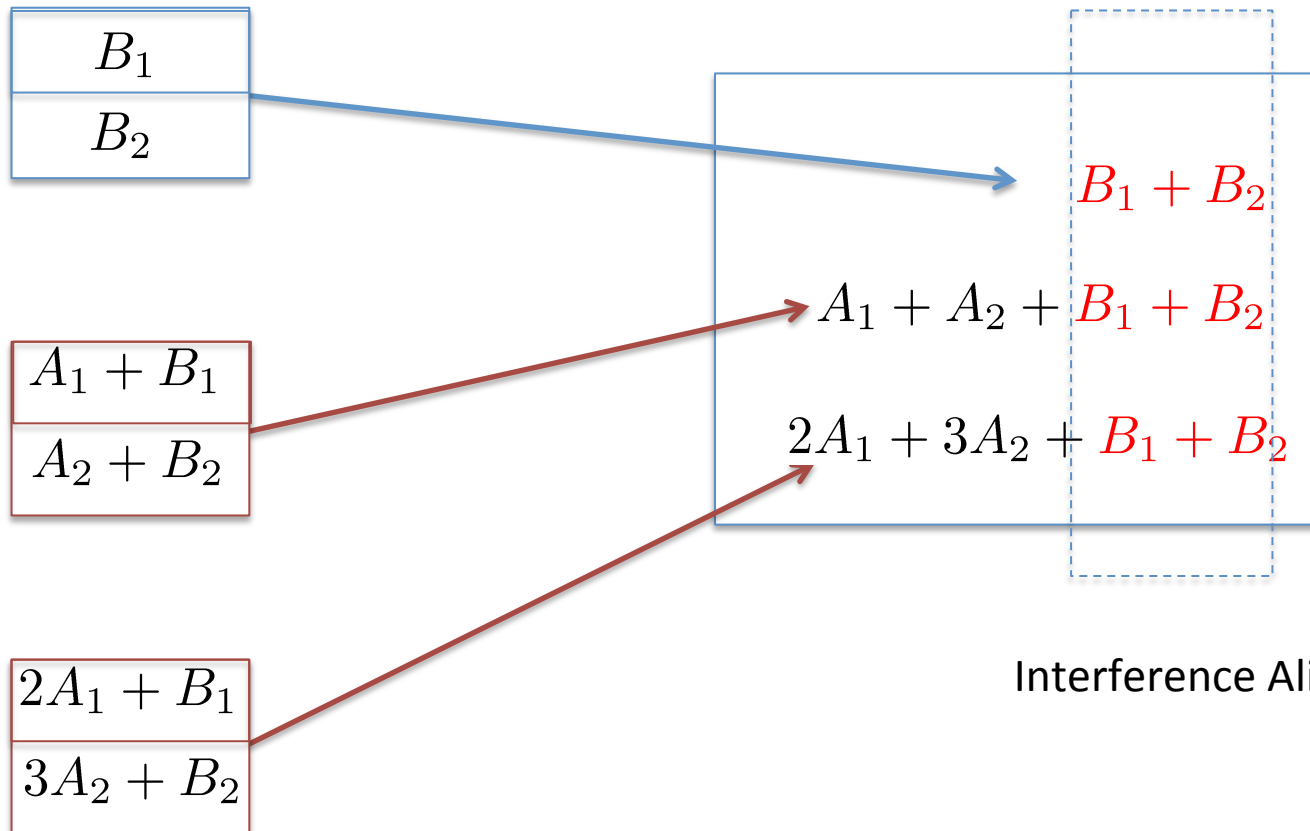
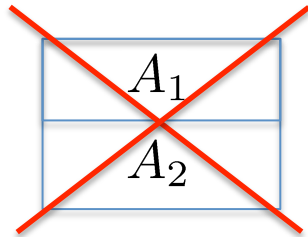


Trivial Repair : 4 linear combinations



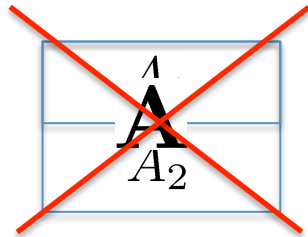
$n = 4, k = 2$, Repair with 3 Linear Combinations

[Wu-Dimakis 09]

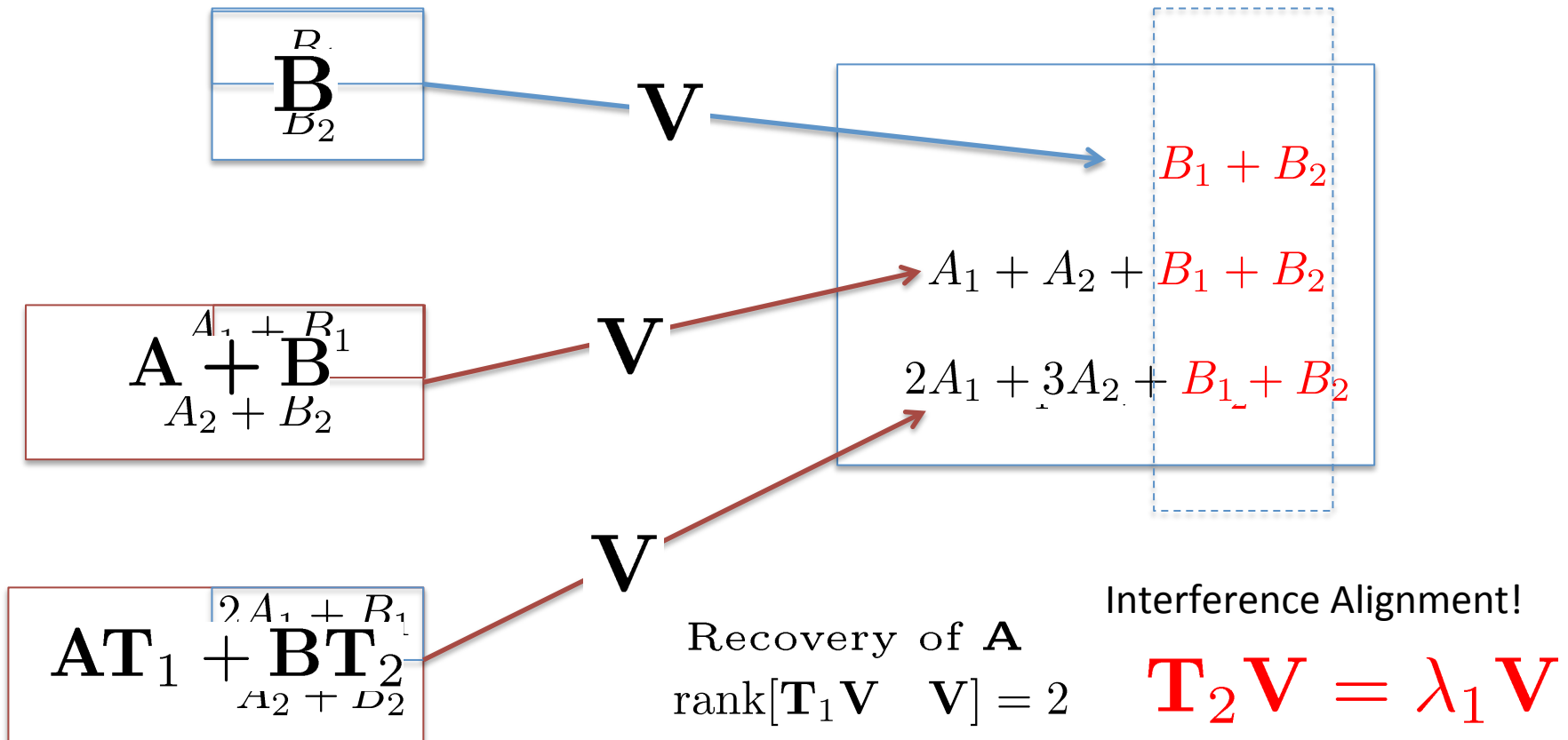


Interference Alignment!

$n = 4, k = 2$, Repair with 3 Linear Combinations



$$\begin{aligned} \mathbf{A}, \mathbf{B} &\rightarrow 1 \times 2 \\ \mathbf{T}_1, \mathbf{T}_2 &\rightarrow 2 \times 2 \\ \mathbf{V} &\rightarrow 2 \times 1 \end{aligned}$$



Repair for $n=4, k=2$

Repair of Node 1

$$\mathbf{T}_2 \mathbf{V} = \lambda_1 \mathbf{V}$$

$$\text{rank}[\mathbf{T}_1 \mathbf{V} \quad \mathbf{V}] = 2$$

$$\mathbf{T}_1, \mathbf{T}_2 \rightarrow 2 \times 2$$

$$\mathbf{V}, \mathbf{W} \rightarrow 2 \times 1$$

Repair of Node 2

$$\mathbf{T}_1 \mathbf{W} = \lambda_2 \mathbf{W}$$

$$\text{rank}[\mathbf{T}_2 \mathbf{W} \quad \mathbf{W}] = 2$$

\mathbf{V} eigen-vector of \mathbf{T}_2

\mathbf{W} eigen-vector of \mathbf{T}_1

Repair Vectors $\mathbf{V}_1, \mathbf{V}_2, \rightarrow$ Beamforming Vectors in Wireless Comm.

Coding matrices $\mathbf{T}_1, \mathbf{T}_2, \rightarrow$ Channel Matrices in Wireless Comm.

Artificial Channels applicable!

Code=Channel Matrices, Repair vectors = Beamforming vectors

$$\text{span}(T_i \mathbf{V}) = \text{span}(\mathbf{V}), i = 1, 2, \dots, N$$

Code (Channel Matrices) can be designed!

1. Asymptotic Alignment Codes

[Cadambe-Jafar-Maleki, 2010]

- Diagonal Channel (Code) Matrices

[Suh-Ramchandran, 2010]

- Approximate Alignment

2. Exact Finite Alignment-Based Codes

[Cadambe-Huang-Li ISIT 2011, Tamo-Wang-Bruck ISIT 2011]

- Non-trivial extensions of “Toy Alignment” Examples
- Tensor Product based Subspace Alignment

[Cadambe-Huang-Li-Jafar Arxiv 2011]

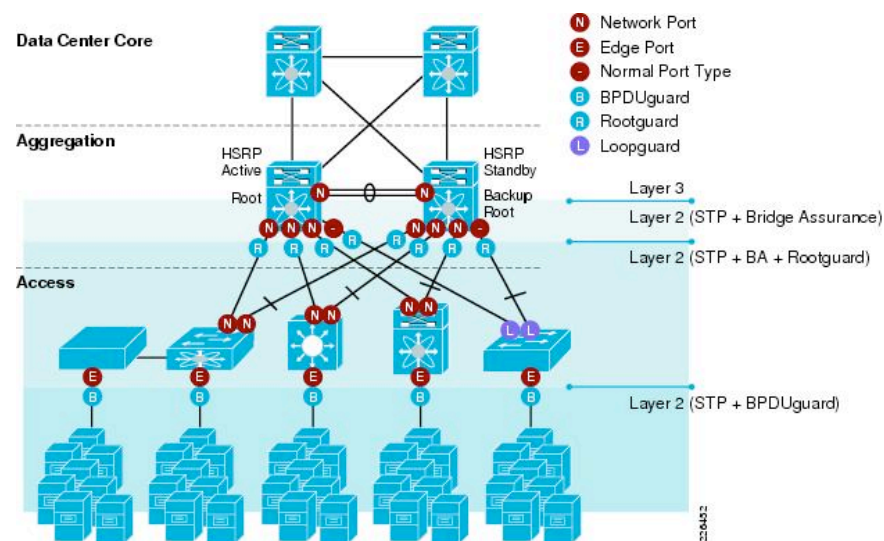


Future Work

Modeling Hierarchical Architecture of Data Centers for Repair

- Racks
- Machines
- Disks

all connected via switches/routers.

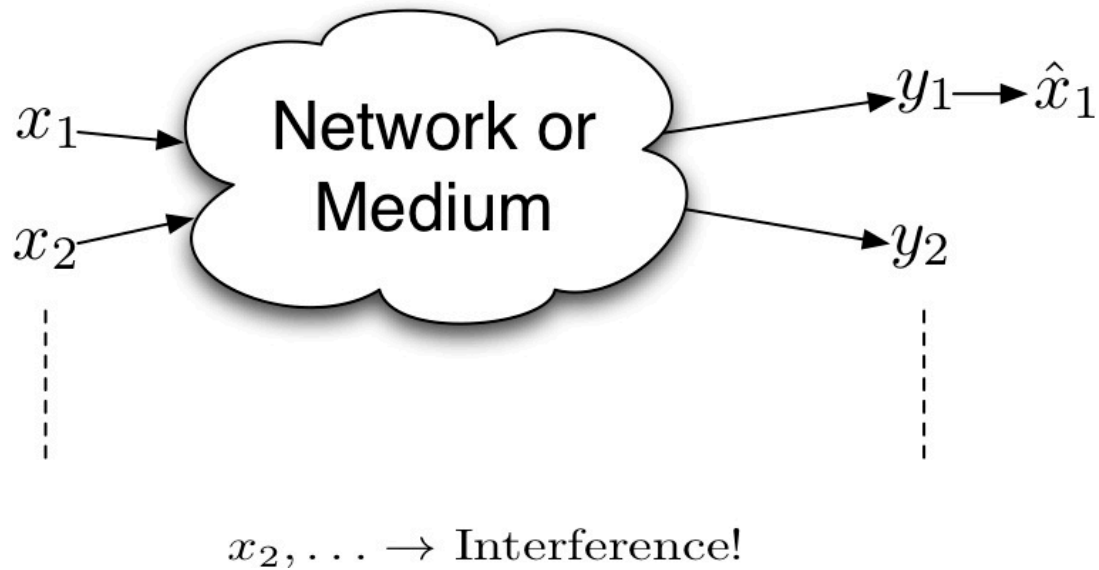


Multi-Source Network Coding

Routing \leftrightarrow Orthogonalization

Structured Network Coding \leftrightarrow Interference Alignment

Future Work



Broader view of alignment as an ingredient of Distributed Systems.

Thank you!

3 users 2 antennas

