6.033 Spring 2018
Lecture #5

- Threads
- Condition Variables
- Preemption
operating systems enforce modularity on a single machine using virtualization in order to enforce modularity + build an effective operating system

1. programs shouldn’t be able to refer to (and corrupt) each others’ memory

2. programs should be able to communicate

3. programs should be able to share a CPU without one program halting the progress of the others
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1. programs shouldn’t be able to refer to (and corrupt) each others’ memory

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**today’s goal:** use threads to allow multiple programs to share a CPU
thread: a virtual processor

thread API:
suspend(): save state of current thread to memory
resume(): restore state from memory
send(bb, message):
    acquire(bb.lock)
    while bb.in - bb.out == N:
        release(bb.lock)
        acquire(bb.lock)
    bb.buf[bb.in mod N] <- message
    bb.in <- bb.in + 1
    release(bb.lock)
    return
send(bb, message):
    acquire(bb.lock)
    while bb.in - bb.out == N:
        release(bb.lock)
        yield()
        acquire(bb.lock)
    bb.buf[bb.in mod N] <- message
    bb.in <- bb.in + 1
    release(bb.lock)
return
yield():
  acquire(t_lock)

  id = cpus[CPU].thread
  threads[id].state = RUNNABLE
  threads[id].sp = SP
  threads[id].ptr = PTR

  do:
    id = (id + 1) mod N
  while threads[id].state != RUNNABLE

  SP = threads[id].sp
  PTR = threads[id].ptr
  threads[id].state = RUNNING
  cpus[CPU].thread = id

  release(t_lock)
**condition variables**: let threads wait for events, and get notified when they occur

**condition variable API**:
- **wait(cv)**: yield processor and wait to be notified of cv
- **notify(cv)**: notify waiting threads of cv
send(bb, message):
    acquire(bb.lock)
    while bb.in - bb.out == N:
        release(bb.lock)
        wait(bb.not_full)
        acquire(bb.lock)
        bb.buf[bb.in mod N] <- message
        bb.in <- bb.in + 1
        release(bb.lock)
    notify(bb.not_empty)
return

(threads in receive() will wait on bb.not_empty and notify of bb.not_full)

problem: lost notify
condition variable API:

wait(cv, lock): yield processor, release lock, wait to be notified of cv

notify(cv): notify waiting threads of cv
send(bb, message):
    acquire(bb.lock)
    while bb.in - bb.out == N:
        wait(bb.not_full, bb.lock)
    bb.buf[bb.in mod N] <- message
    bb.in <- bb.in + 1
    release(bb.lock)
    notify(bb.not_empty)
    return
wait(cv, lock):
    acquire(t_lock)
    release(lock)
    id = cpus[CPU].thread
    threads[id].cv = cv
    threads[id].state = WAITING
    yield_wait()
    release(t_lock)
    acquire(lock)

notify(cv):
    acquire(t_lock)
    for id = 0 to N-1:
        if threads[id].cv == cv &&
           threads[id].state == WAITING:
            threads[id].state = RUNNABLE
    release(t_lock)
yield_wait(): // called by wait()
    acquire(t_lock)

    id = cpus[CPU].thread
    threads[id].state = RUNNABLE
    threads[id].sp = SP
    threads[id].ptr = PTR

do:
    id = (id + 1) mod N
while threads[id].state != RUNNABLE

    SP = threads[id].sp
    PTR = threads[id].ptr
    threads[id].state = RUNNING
    cpus[CPU].thread = id

release(t_lock)

problem: wait() holds t_lock
yield_wait(): // called by wait()

```

id = cpus[CPU].thread
threads[id].state = RUNNABLE
threads[id].sp = SP
threads[id].ptr = PTR

do:
    id = (id + 1) mod N
while threads[id].state != RUNNABLE

SP = threads[id].sp
PTR = threads[id].ptr
threads[id].state = RUNNING
cpus[CPU].thread = id
```

**problem:** current thread’s state shouldn’t be RUNNABLE
yield_wait(): // called by wait()

    id = cpus[CPU].thread
    threads[id].sp = SP
    threads[id].ptr = PTR

    do:
        id = (id + 1) mod N
    while threads[id].state != RUNNABLE

    SP = threads[id].sp
    PTR = threads[id].ptr
    threads[id].state = RUNNING
    cpus[CPU].thread = id

problem: deadlock (wait() holds t_lock)
yield_wait():  // called by wait()

id = cpus[CPU].thread
threads[id].sp = SP
threads[id].ptr = PTR

do:
   id = (id + 1) mod N
   release(t_lock)
   acquire(t_lock)
while threads[id].state != RUNNABLE

SP = threads[id].sp
PTR = threads[id].ptr
threads[id].state = RUNNING
cpus[CPU].thread = id

problem: stack corruption
yield_wait(): // called by wait()

```c
id = cpus[CPU].thread
threads[id].sp = SP
threads[id].ptr = PTR
SP = cpus[CPU].stack

d: [id = (id + 1) mod N]
   release(t_lock)
   acquire(t_lock)
   while threads[id].state != RUNNABLE

   SP = threads[id].sp
   PTR = threads[id].ptr
   threads[id].state = RUNNING
   cpus[CPU].thread = id
```
**preemption:** forcibly interrupt threads

```
    timer_interrupt():
        push PC
        push registers
        yield()
        pop registers
        pop PC
```

**problem:** what if timer interrupt occurs while running `yield()` or `yield_wait()`?
preemption: forcibly interrupt threads

```c
void timer_interrupt()
{
    push PC
    push registers
    yield()
    pop registers
    pop PC
}
```

solution: hardware mechanism to disable interrupts
• **Threads** virtualize a processor so that we can share it among programs. `yield()` allows the kernel to suspend the current thread and resume another.

• **Condition Variables** provide a more efficient API for threads, where they `wait` for an event and are `notified` when it occurs. `wait()` requires a new version of `yield()`, `yield_wait()`.

• **Preemption** forces a thread to be interrupted so that we don’t have to rely on programmers correctly using `yield()`. Requires a special `interrupt` and hardware support to disable other interrupts.