## Tensor Product Representations

 and Holographic Reduced Representations Smolensky, 1990 \& Plate, 1991Tiwalayo Eisape, Joey Velez-Ginorio, Pedro Colon-Hernandez \{eisape, joeyv, pe2517\}@mit.edu

## Outline

Introductions (us +3 others) (11:35-11:40)

TPRs - why/what? (11:40-11:55)


Break out room (11:55-12:10 mins)
Discussion (12:10-12:20 mins)
[Early] Break (12:20-12:35)

TPR tutorial (12:35-12:50)

Discussion (12:50-12:55)

TPR Shortcomings; HRRs (12:55-1:10)

Discussion (1:10-1:25)

## Outline

1. Is variable binding necessary?
2. Do humans use a TPR-like mechanism?
3. Do current models approximate faithfulness?
4. Small group technical questions


Full Group Discussion (12:50-12:55)

TPR Shortcomings; HRRs (12:55-1:10)

Full Group Discussion (1:10-1:25)

## Tensor Product Representations - why?

A one-sentence summary of the implications of this view for Al :
connectionist models may well offer an opportunity to escape the brittleness of symbolic Al systems ...

This paper offers an example of what such a collaboration might look like.

## Tensor Product Representations - why?

Jay is loved by Kay. Who loves Jay? Kay.

- Jay in role: subject of passive sentence
- Jay in role: object of wh-question
- Kay in role: object of passive by-phrase
- Kay in role: answer to wh-question


## Tensor Product Representations - why?



## Tensor Product Representations - why?



## Tensor Product Representations - what?

Representing Structured Objects

(1) Decomposing the structures via roles (2) representing variable/value bindings (3) representing conjunctions

## Tensor Product Representations - what?

## Tensor Product Representations - what?



## Tensor Product Representations - what?



## Tensor Product Representations - what?



## Tensor Product Representations - what?



## Tensor Product Representations - what?


(3) representing conjunctions

## Tensor Products

'Faithful' Tensor Product Representations

## Tensor Products

## Faithfulness

$$
\begin{array}{l|llll}
\text { © } & & & & \\
0 & & & & \\
0 & & & & \\
0 & & & & \\
\text { (2) } & & & & \\
\hline & \odot & 0 & 0 & 0
\end{array}
$$

## Faithfulness



## Faithfulness




## Faithfulness



## Faithfulness



## Faithfulness



## Faithfulness



Orthogonality
Theorem 3.3, Section 3.2

Linear Independence
Definition 2.8, Section 2.2.2

## Tensor Products



## 'Faithful' <br> Tensor Product Representations

Graceful Decay

## Variable Binding

1. Is Variable Binding Necessary?
2. Do humans use a TPR-like mechanism?
3. Do current models approximate faithfulness?
4. Small group technical questions

## Gary Marcus @GaryMarcus • Feb 6, 2018

completely agreed, @tdietterich! lots of cases where variable binding is absolutely necessary. no binding, no AGI.

Thomas G. Dietterich @tdietterich•Feb 6, 2018
Replying to @tdietterich @jahendler and @GaryMarcus
There are lots of cases where binding appears to be necessary. Ex 1: If you put X into your pocket and then walk to work, you will be able to take $X$ out of your pocket at work. Ex 2: If I ask you query $X$ and you know $X$, you will tell me $X$, forall $X$.

Break


## TPR Tutorial

## TPR Tutorial

(1) Symbolic Structure

## TPR Tutorial

(1) Symbolic Structure
(2) Encoding w/ TPRs

## TPR Tutorial

(1) Symbolic Structure
(2) Encoding w/ TPRs
(3) Representation Proofs
(1) Symbolic Structure : Give, the programming language

## (1) Symbolic Structure : Give, the programming language

```
Syntax
p ::= (Give p) |
```


## (1) Symbolic Structure : Give, the programming language

## Syntax <br> p ::= (Give p) |

Examples
(Give $\square$ )
(Give (Give $\square$ ))
(Give (Give (Give $\square$ ))

## (1) Symbolic Structure : Give, the programming language

```
Syntax
p ::= (Give p) |
```


## Examples

(Give $\square$ )
(Give (Give $\square$ ))
(Give (Give (Give $\square$ ))

## Semantics

(Give $p$ ) $\rightarrow p$

$$
\frac{\mathrm{p} \rightarrow \mathrm{p}^{\prime}}{(\text { Give } \mathrm{p}) \rightarrow\left(\text { Give } \mathrm{p}^{\prime}\right)}
$$

## (1) Symbolic Structure : Give, the programming language



## Semantics

$$
\begin{gathered}
(\text { Give } p) \rightarrow p \\
p \rightarrow p^{\prime} \\
(\text { Give } p) \rightarrow\left(\text { Give } p^{\prime}\right)
\end{gathered}
$$

## (2) Encoding w/ TPRs : Give, the programming language

A TPR is a mapping, 【p】: Give ${ }_{4} \mapsto R^{4 \times 3}$, from a set of symbols to a vector space via filler/role decompositions. Here, Give ${ }_{4}$ denotes the set of all Give programs up to length 4.

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\text { Give }_{4}=\{\square \text {, (Give } \square \text { ), (Give (Give } \square \text { )), (Give (Give (Give } \square \text { )))\} }
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R^{4 \times 3}=\left\{\begin{array}{lll}
{\left[\begin{array}{lll}
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\end{array}\right]} & {\left[\begin{array}{lll}
1 & 2 & 0
\end{array}\right]} & {\left[\begin{array}{lll}
2 & 2 & 1
\end{array}\right]} \\
{\left[\begin{array}{lll}
1 & 1 & 1
\end{array}\right],\left[\begin{array}{lll}
0 & 1 & 0
\end{array}\right],\left[\begin{array}{lll}
1 & 1 & 3
\end{array}\right], \ldots} \\
{\left[\begin{array}{lll}
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\end{array}\right]
\end{gathered}
$$

## Fillers

$f=\left\{i_{1}, i_{2}, i_{3}, i_{4}\right\}$

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\end{array}\right]}
\end{array}\right]
\end{gathered}
$$

## Fillers

$f=\left\{i_{1}, i_{2}, i_{3}, i_{4}\right\}$

$$
\begin{aligned}
& \text { Roles } \\
& r=\{\text { Give }, \square, \varepsilon\}
\end{aligned}
$$

## (2) Encoding w/TPRs : Give, the programming language

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## Fillers

$f=\left\{i_{1}, i_{2}, i_{3}, i_{4}\right\} \quad($ Give $($ Give $($ Give $\square)))=\left(i_{1}\right.$ :Give $) \wedge\left(i_{2}:\right.$ Give $) \wedge\left(i_{3}:\right.$ Give $) \wedge\left(i_{4}:\right.$

Roles
$r=\{$ Give, $\square, \varepsilon\}$

## (2) Encoding w/ TPRs : Give, the programming language

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\end{array}\right] \quad\left[\begin{array}{lll}
2 & 2 & 1
\end{array}\right]}
\end{aligned}
$$

## Fillers

$f=\left\{i_{1}, i_{2}, i_{3}, i_{4}\right\} \quad$ (Give (Give (Give $\left.\left.\square\right)\right)=\left(i_{1}:\right.$ Give $) \wedge\left(i_{2}:\right.$ Give $) \wedge\left(i_{3}:\right.$ Give $) \wedge\left(i_{4}:\right.$ ]) $($ Give $($ Give $\square))=\left(\mathrm{i}_{1}: \varepsilon\right) \wedge\left(\mathrm{i}_{2}:\right.$ Give $) \wedge\left(\mathrm{i}_{3}\right.$ :Give $) \wedge\left(\mathrm{i}_{4}\right.$ :

## Roles

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$\square)($ Give $($ Give $\square))=\left(i_{1}: \varepsilon\right) \wedge\left(i_{2}:\right.$ Give $) \wedge\left(i_{3}\right.$ :Give $) \wedge\left(i_{4}:\right.$
Roles

$$
(\text { Give } \square)=\left(\mathrm{i}_{1}: \varepsilon\right) \wedge\left(\mathrm{i}_{2}: \varepsilon\right) \wedge\left(\mathrm{i}_{3}: \text { Give }\right) \wedge\left(\mathrm{i}_{4}:\right.
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1 & 0 & 1
\end{array}\right]}
\end{array}\right]
\end{gathered}
$$

## Fillers

$\begin{array}{ccc}{[1]} & {[0]} & {[0]} \\ {[0]} & {[0]} \\ {[0]} & {[0]} & {[0]} \\ {[0]} & {[0]} & {[1]} \\ {[0]} & {[0]} & \llbracket(\text { Give }(\text { Give }(\text { Give } \square))) \rrbracket=\llbracket\left(i_{1}: \text { Give }\right) \wedge\left(i_{2}: \text { Give }\right) \wedge\left(i_{3}: \text { Give }\right) \wedge\left(i_{4} \text { : }\right.\end{array}$ $\mathrm{f}= \begin{cases}{[0]} \\ {[0]} & {[0],\left[\begin{array}{ll}{[1]} \\ {[0]}\end{array},\left[\begin{array}{ll}{[0]}\end{array}\right\}\right.} \\ {\left[\begin{array}{l}{[0]}\end{array}\right]}\end{cases}$

## Roles

$r=\left\{\begin{array}{lll}{[10} & 0\end{array}\right]\left[\begin{array}{lll}{[01} & 0\end{array}\right]\left[\begin{array}{lll}0 & 0 & 1\end{array}\right]$

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\end{array}\right]}
\end{aligned}
$$

## Fillers

$\begin{array}{llll}{[1]} & {[0]} & {[0]} & {[0]} \\ {[0]} & {[1]} & {[0]} & {[0]}\end{array}$
$\left[\begin{array}{llll}{[0]} & {[0]} & {[1]} & {[0]}\end{array}\right.$
【（Give（Give（Give $\square$ ）））
）］
$\rrbracket=\llbracket\left(i_{1}:\right.$ Give $) \wedge\left(i_{2}:\right.$ Give $) \wedge\left(i_{3}\right.$ ：Give $) \wedge\left(i_{4}:\right.$ $=\left(\mathbf{i}_{1} \otimes\right.$ Give $)+\left(\mathbf{i}_{2} \otimes\right.$ Give $)+\left(\mathbf{i}_{3} \otimes\right.$ Give $)+\left(\mathbf{i}_{4} \otimes \square\right)$

## Roles

$r=\left\{\begin{array}{lll}{[10} & 0\end{array}\right]\left[\begin{array}{lll}{[01} & 0\end{array}\right]\left[\begin{array}{lll}0 & 0 & 1\end{array}\right]$

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\end{aligned}
$$

$$
\begin{aligned}
& \text { Fillers } \\
& \left.\begin{array}{cccc}
{[1]} & {[0]} & {[0]} & {[0]} \\
{[0]} & {[1]} & {[0]} & {[0]} \\
{[0]} & {[0]} & {[1]} \\
{[0]} & {[0],} \\
i_{1} & i_{2} & {[0]} & i_{3} \\
\hline & i_{4}
\end{array}\right\} \\
& \begin{array}{l}
\text { 【(Give (Give (Give } \square \text { ))) } \\
\square) \rrbracket
\end{array} \\
& \rrbracket=\llbracket\left(i_{1} \text { :Give }\right) \wedge\left(i_{2} \text { :Give }\right) \wedge\left(i_{3} \text { :Give }\right) \wedge\left(i_{4}:\right. \\
& =\left(\mathbf{i}_{\mathbf{1}} \otimes \text { Give }\right)+\left(\mathbf{i}_{\mathbf{2}} \otimes \text { Give }\right)+\left(\mathbf{i}_{\mathbf{3}} \otimes \text { Give }\right)+\left(\mathbf{i}_{\mathbf{4}}^{\otimes} \square\right) \\
& =\left[\begin{array}{lll}
1 & 0 & 0
\end{array}\right] \\
& \text { Roles } \\
& r=\left\{\begin{array}{llll}
{[10} & 0
\end{array}\right],\left[\begin{array}{llll}
\text { Give } & 1 & 0
\end{array}\right], \quad\left[\begin{array}{lll}
0 & 0 & 1
\end{array}\right]
\end{aligned}
$$

(3) Representation Proofs: Give, the programming language

Theorem. The following linear transformation is a representation of the instruction Give.

Give : $\sum_{i} \mathbf{f}_{\mathbf{i}} \otimes \mathbf{r}_{\mathbf{i}} \mapsto \sum_{i} \mathbf{f}_{\mathbf{i}} \otimes \mathbf{r}_{\mathbf{i}}$

## （3）Representation Proofs ：Give，the programming language

Theorem．The following linear transformation is a representation of the instruction Give．

$$
\text { Give : } \sum_{i} \mathbf{f}_{i} \otimes r_{i} \mapsto \sum_{i} \mathbf{f}_{i} \otimes r_{i}
$$

Proof．
Recall that w／TPRs we encode Give programs as conjunctions of filler／role decompositions，i．e．$\llbracket p \rrbracket=\sum_{i} \mathbf{f}_{\mathbf{i}} \otimes \mathbf{r}_{\mathbf{i}}$ ．Additionally，recall that：（Give p$) \rightarrow \mathrm{p}$

$$
\begin{aligned}
& \text { 【(Give p)】 = 【p】 } \\
& =\Sigma_{i} \mathbf{f}_{i} \otimes r_{i} \\
& =\text { Give } \sum_{i} \mathbf{f}_{i} \otimes r_{i} \\
& \text { = Give 【p】 }
\end{aligned}
$$

## Discussion

- How does this scale to larger programs in Give?
- What if our programming language was more complicated?
- Other thoughts...


## Benefits \& Shortcomings of Tensor Decomposition

+ No impositions on structure
+ Faithful
+ Variable binding
- Scaling up can be memory and compute demanding
- Using ConceptNet as an example, $\sim 4 \mathrm{M}$ nodes, $\sim 40$ relations might need to play around with pretty large tensors



## Holographic Reduced Representations

- Use Circular Convolutions and Correlations to associate/disassociate vectors that represent structures
- Requires a reconstruction system to sort through the noise
- Circular Conv. and Circular Corr. can be manipulated to query structure


Circular Convolution


## Representations with HRR

Sequences

$$
\begin{aligned}
\mathbf{s}_{a b c} & =\mathbf{a}+\mathbf{a} \odot \mathbf{b}+\mathbf{a} \odot \mathbf{b} \odot \mathbf{c} \\
\mathbf{s}_{d e} & =\mathbf{d}+\mathbf{d} \odot \mathbf{e} \\
\mathbf{s}_{\boldsymbol{f} \boldsymbol{h}} & =\mathbf{f}+\mathbf{f} \odot \mathbf{g}+\mathbf{f} \odot \mathbf{g} \odot \mathbf{h}
\end{aligned}
$$

$\mathrm{S}($ abcdefgh $)=\mathbf{s}_{\boldsymbol{a} b \boldsymbol{c}}+\mathbf{s}_{\boldsymbol{a} b \boldsymbol{c}} \odot \mathbf{s}_{\boldsymbol{d} e}+\mathbf{s}_{a b c} \odot \mathbf{s}_{d e} \odot \mathbf{s}_{f g h}$.

## Representations with HRR

## Sequences

Variable Binding

$$
\tilde{\mathbf{t}}=\tilde{\mathbf{x}} \circledast \tilde{\mathbf{a}}+\tilde{\mathbf{y}} \circledast \tilde{\mathbf{b}} .
$$

Binding a to $X$ and $b$ to $Y$

## Representations with HRR

Running frame: Spot runs

## Sequences

Variable Binding
Frame-Slots

$$
\mathbf{t}_{\text {running }}=\mathbf{l}_{\text {run }}+\mathbf{r}_{\text {agent }} \oplus \mathbf{f}_{\text {spot }}
$$

Seeing frame: Dick saw Spot run

$$
\begin{aligned}
\mathbf{t}_{\text {seeing }}= & \mathbf{l}_{\text {see }}+\mathbf{r}_{\text {agent }} \odot \mathbf{f}_{\text {dick }}+\mathbf{r}_{\text {object }} \odot \mathbf{t}_{\text {running }} \\
= & \mathbf{l}_{\text {see }}+\mathbf{r}_{\text {agent }} \odot \mathbf{f}_{\text {dick }} \\
& +\mathbf{r}_{\text {object }} \odot\left(\mathbf{l}_{r u n}+\mathbf{r}_{\text {agent }} \odot \mathbf{f}_{\text {spot }}\right)
\end{aligned}
$$

## Example: Filling a frame

Frame:
Job Application:

- Name
- Date

Filler

- September 1, 2020


## Example: Filling a frame

| Frame: | Job Application | 0.35 | 0.28 |
| :--- | :--- | :--- | :--- |$\quad 0.11$

## Example: Filling a frame

Binding Date \& Filler

September 1, 2020
$C[0]=0.1 * 0.05+-0.16 * 0.04+-0.22 * 0.06+=-0.0146$
C: $\quad-0.0146$


## Example: Filling a frame

Binding Date \& Filler

September 1, 2020
$C[1]=0.1^{*}-0.16+0.05^{*}-0.22+0.04^{*} 0.06=-0.0246$

C: | -0.0146 | -0.0246 |
| :--- | :--- | :--- |



## Example: Filling a frame

Pairing Job Application + Date Field


September 1, 2020


## Example: Filling a frame

September 1, 2020

C: $\quad$| -0.0146 | -0.0246 | 0.0432 |
| :--- | :--- | :--- |

C: \{Date:September 1,2020\}


## Example: Filling a frame

September 1, 2020


## Example: Filling a frame

C': \{Date: September 1, 2020, Name\}

## Example: Filling a frame



## Example: Filling a frame

C": \{Job Application: Date:September 1,2020, Name\}

Keep in mind representations are stored in a distributed manner We used the "decoder" implicitly to clean the noise Our representations are the result of an FFT

## Example Application for Holographic Representation

## Best role finder:



## Example Application for Holographic Representation

## Best role finder:

Job Application: Name, Date



January, 1, 2021

## Example Application for Holographic Representation

## Best role finder:

Job Application: Name, Date



January, 1, 2021

## Example Application for Holographic Representation

## Best role finder:

Job Application: Name, Date



January, 1, 2021

## Example Application for Holographic Representation

## Best role finder:

Job Application: Name, Date



January, 1, 2021

## Benefits of Holographic Representations

- Format for the two input vectors is not specified, only independently distributed
- Space Efficiency: you just need the 2 vectors rather than the whole Tensor, result is the same size as the input
- Can be calculated in $\mathrm{O}(\mathrm{n} \log \mathrm{n})$ with FFT
- HRRs could retain ambiguity while processing ambiguous input (New York as City and as Name)
- Easy analysis of capacity, scaling and generalization


## Shortcomings of Holographic Representations

- Decoder/cleaner must store all the possible outputs. If it knows everything, then why not find a way to exploit it?
- Is the decoder static? How would you add some new domain?
- Elements of each vector must be independently distributed, but have meaningful features
- Hit until you decode the correct thing?
- Some operations to decode require additional machinery (recursive)


Figure 5: A chunked sequence readout machine.

## Encoding Methods?



## Encoding Methods?


[END]

